

Characterization of the “Size and Shape” of Static RPKI

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Outline

- Goals
- Methodology
- Quantitative analysis of the current resource allocations
- Characterization of static RPKI
- Conclusions

NIST's Goals

- Develop models of the “size and shape” of a potential global RPKI structure from existing RIR/IRR databases.
- Provide quantitative analyses of the scalability and the potential performance impact of global-scale deployed RPKI on routing dynamics.
- Study the potential future changes in routing information infrastructure.
- Evaluate how such issues as IPv4 address exhaustion will impact on the deployed RPKI.
- Assess the potential load and weaknesses of the “moving parts” of the proposed RPKI infrastructure.

Methodology (1 of 2)

- **Use NIST TERRAIN DB data:**
 - Global bulk Whois databases:
 - * 5 RIRs and IRRs from the RADB site.
 - BGP trace data:
 - * RIPE NCC and Route Views.
- **Develop models of the potential global RPKI infrastructure:**
 - Select all distinctively registered objects.
 - For multiple registrations across RIRs:
 - * Select one from a RIR where the resource is allocated to, if exists.
 - * If not, select one arbitrarily among RIRs/IRRs.
 - * *For APNIC, the same resource may be registered in different registries such as RIR and/or NIR. In this case, select one that contains the “status:” attribute.*
 - Build number resources (IPv4 and ASN) structures describing allocation chains.
 - Classify selected objects per region **based on IANA allocation registries:**
 - * ARIN / RIPE / APNIC / AFRINIC / LACNIC / LEGACY / ERX.

Methodology (2 of 2)

- Details of building number resources structures:
 - ASNs:
 - * For SWIP:
 - Distinct ASHandles.
 - Distinct ASNs (aut-nums) registered in RPSL (i.e., aut-num), which are assigned to ARIN but not registered in SWIP (as either a single ASN or AS range).
 - * For RPSL:
 - Unique aut-nums.
 - as-block objects that contain a range of ASNs in RPSL. *Note that some as-blocks contain a single ASN (e.g., ASn – ASn), most of which have corresponding either aut-num or ASHandle objects.*
 - IPv4 addresses:
 - * Globally distinct inetnums in RPSL and NetRanges in SWIP.
 - * For multiple registrations, select one from a RIR where the resource is allocated to, if exists.
 - * If not, select one arbitrarily among RIRs/IRRs.
 - * *Partial registrations from a /8 block may be found in other RIRs but they are considered to belong to the same RIR where the /8 is allocated*
 - * *Exceptions in LEGACY/ERX IP address space:*
 - *The LEGACY/ERX blocks may contain a large number of cross-RIR partial allocations, especially between RPSL and SWIP. These partial allocations are combined before processing.*
 - *Example: If 129.1/16 registered in RIPE (RPSL) and 129.2/16 registered in ARIN (SWIP), then both 129.1/16 and 129.2/16 are considered as LEGACY/ERX.*

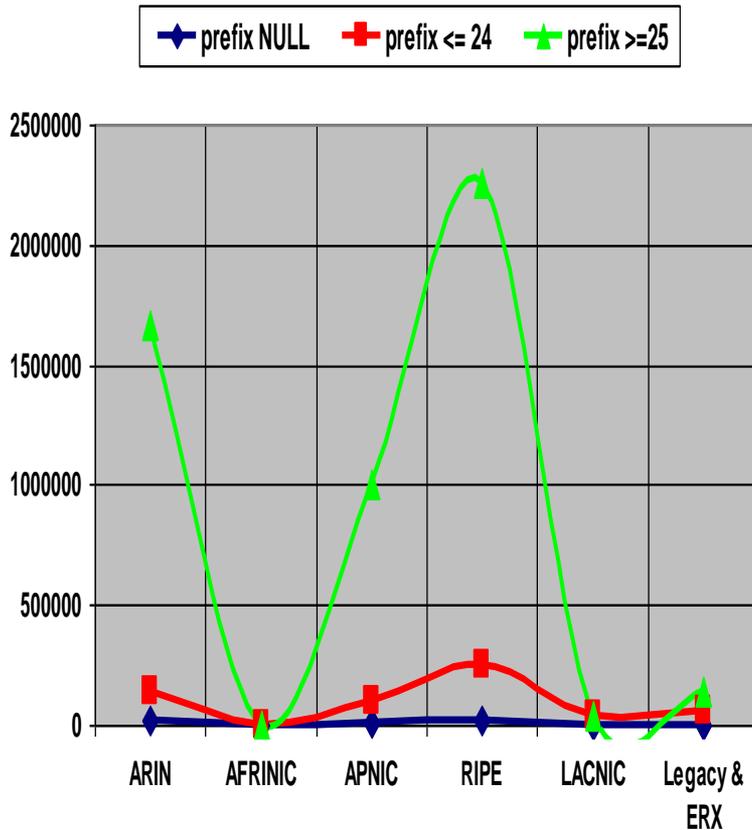
ERX Partial Allocations Examples

- 129/8: currently administered by ARIN:
 - Partial allocations in SWIP: 396
 - Partial allocations in RPSL: 592
 - Multi registrations in both SWIP and RPSL: 30
- 151/8: currently administered by RIPE NCC:
 - Partial allocations in RPSL: 6,999
 - Partial allocations in SWIP: 2,084
 - Multi registrations in both SWIP and RPSL: 15
- 198/8: currently administered by ARIN
 - Partial allocations in RPSL: 320
 - Partial allocations in SWIP: 15,760
 - Multi registration in both SWIP and RPSL: 63

Distribution of Registry IPv4 Address Allocations/Assignments

Registry data date: 2009-02-18

Unit: 1k objects



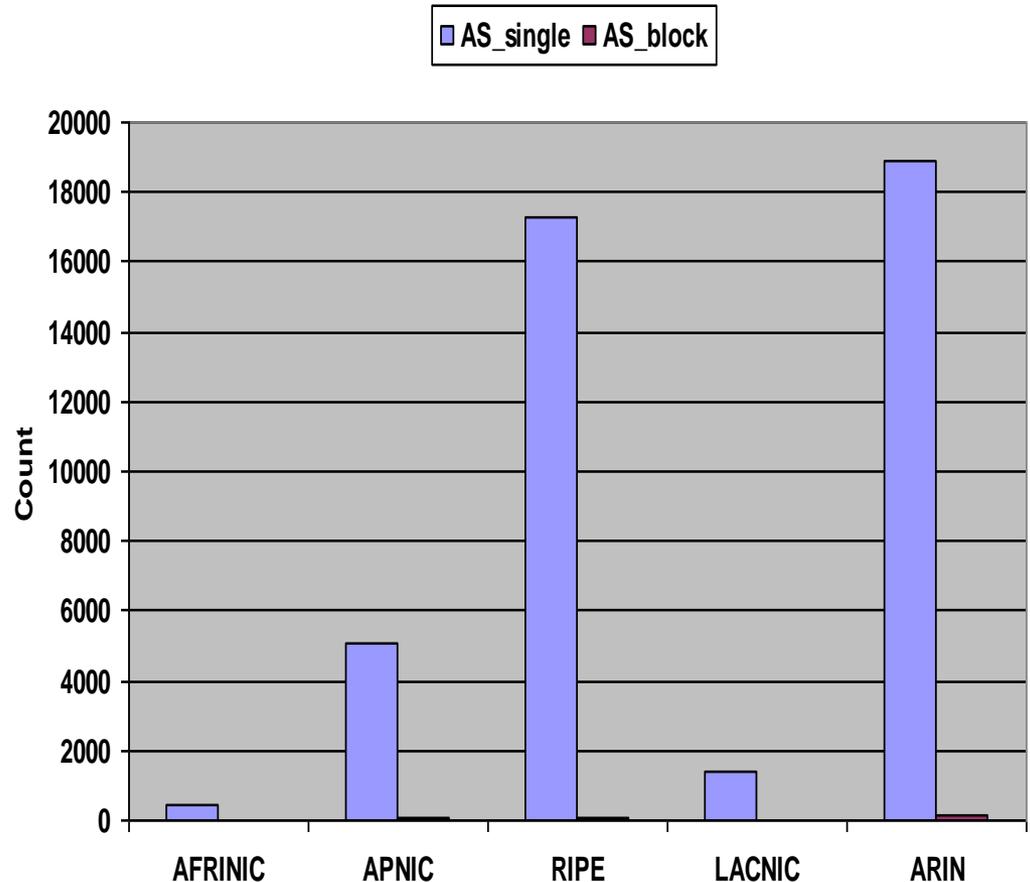
RIR	# of /8 blocks	# objects p_len=NULL*	# objects p_len <= 24	# objects p_len >= 25	Total # objects
ARIN	31	17	145	1,667	1,829
RIPE	28	24	248	2,262	2,534
APNIC	30	13	100	1,004	1,117
AfriNIC	2	0	1	5	6
LACNIC+	6	0	36	40	76
LEGACY/ERX+	92	4	59	144	207
Total	189	58	589	5,122	5,769

* Prefix Length NULL indicates that an address block cannot be represented by a single CIDR.
 + from both RPSL and SWIP except duplicates.
 As of August 2010, 14 /8 blocks are unallocated.

Distribution of Global ASN Assignment Based on IANA and RIR/IRR Datasets

Registry data date: 2009-02-18

RIR	AS single	AS block
ARIN	18,862	137
RIPE	17,280	59
APNIC	5,082	70
AfriNIC	406	4
LACNIC	1,391	2
Total	43,021	272

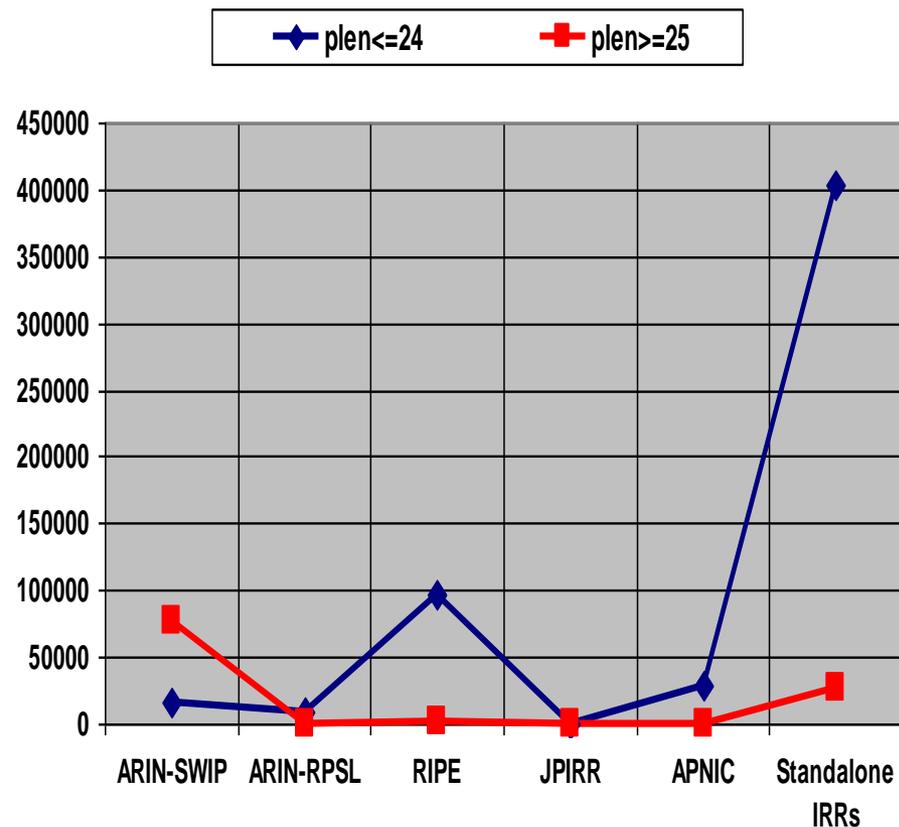


Distribution of Potential ROAs Based on Route Object Registrations

Registry data date: 2009-02-18

Unit: 1k objects

RIR	# objects p_len <= 24	# objects p_len >= 25	Total # objects
ARIN-SWIP	15.5	77	92.6
ARIN-RPSL	8.2	0.2	8.4
RIPE	96.4	1.5	97.9
JPIRR	0.6		0.6
APNIC	28	0	28
Standalone IRRs +	403	27.3	430.3
Total	551.7	106	657.8



+ Standalone IRRs includes all individual IRRs mirrored from the RADB site.

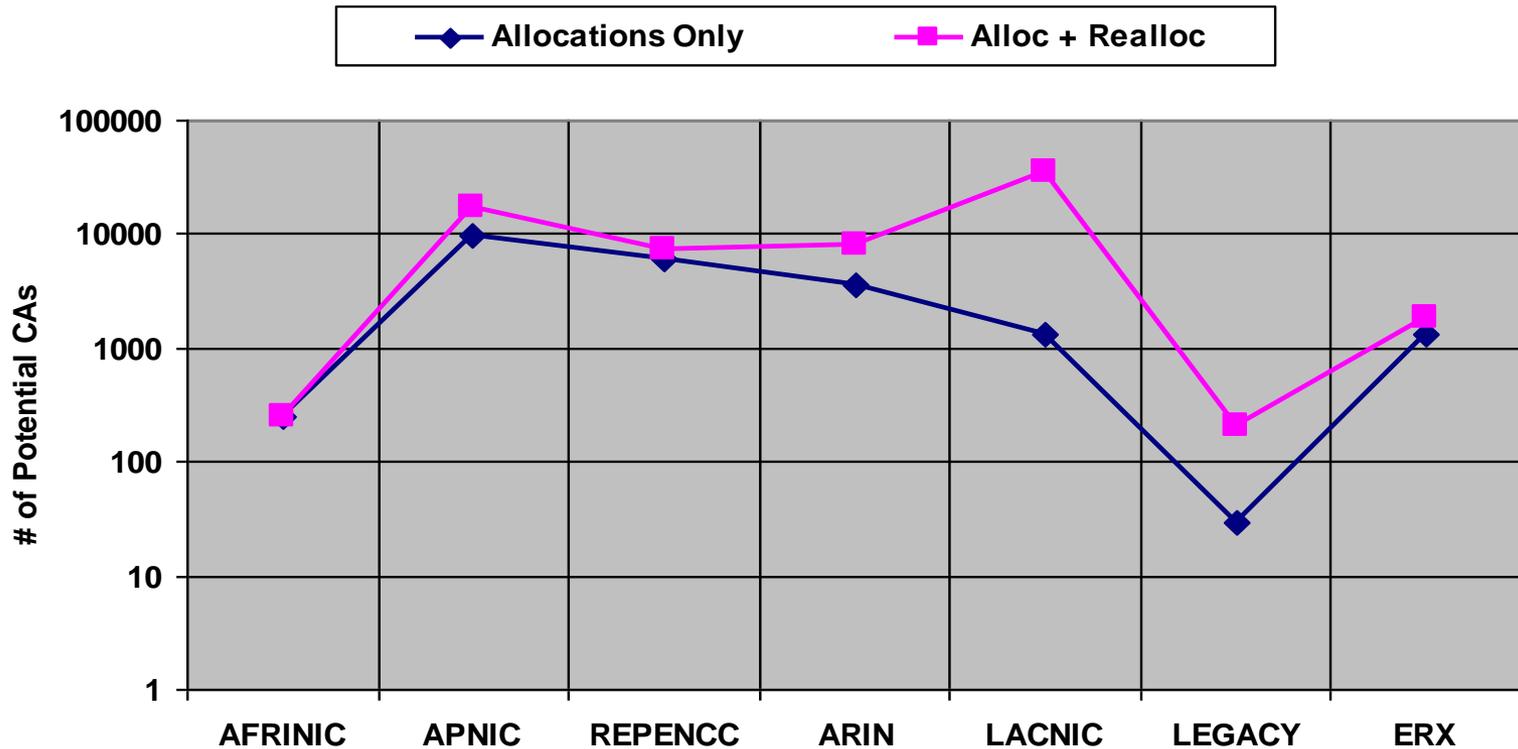
Characterization of the of static RPKI

- Analysis of potential CAs:
 - Distribution of potential CAs per RIR
 - Distribution of CA path depths per RIR
- Analysis of IPv4 certificates:
 - Full deployment vs. optimized deployment
 - IPv4 prefix lengths vs. IPv4 certification path depths
- Analysis of ROAs:
 - The cost estimate of ROA verifications in terms of certification path lengths
 - Distribution of PI address space.
 - Analysis of MOASes of potential ROAs

Potential CAs

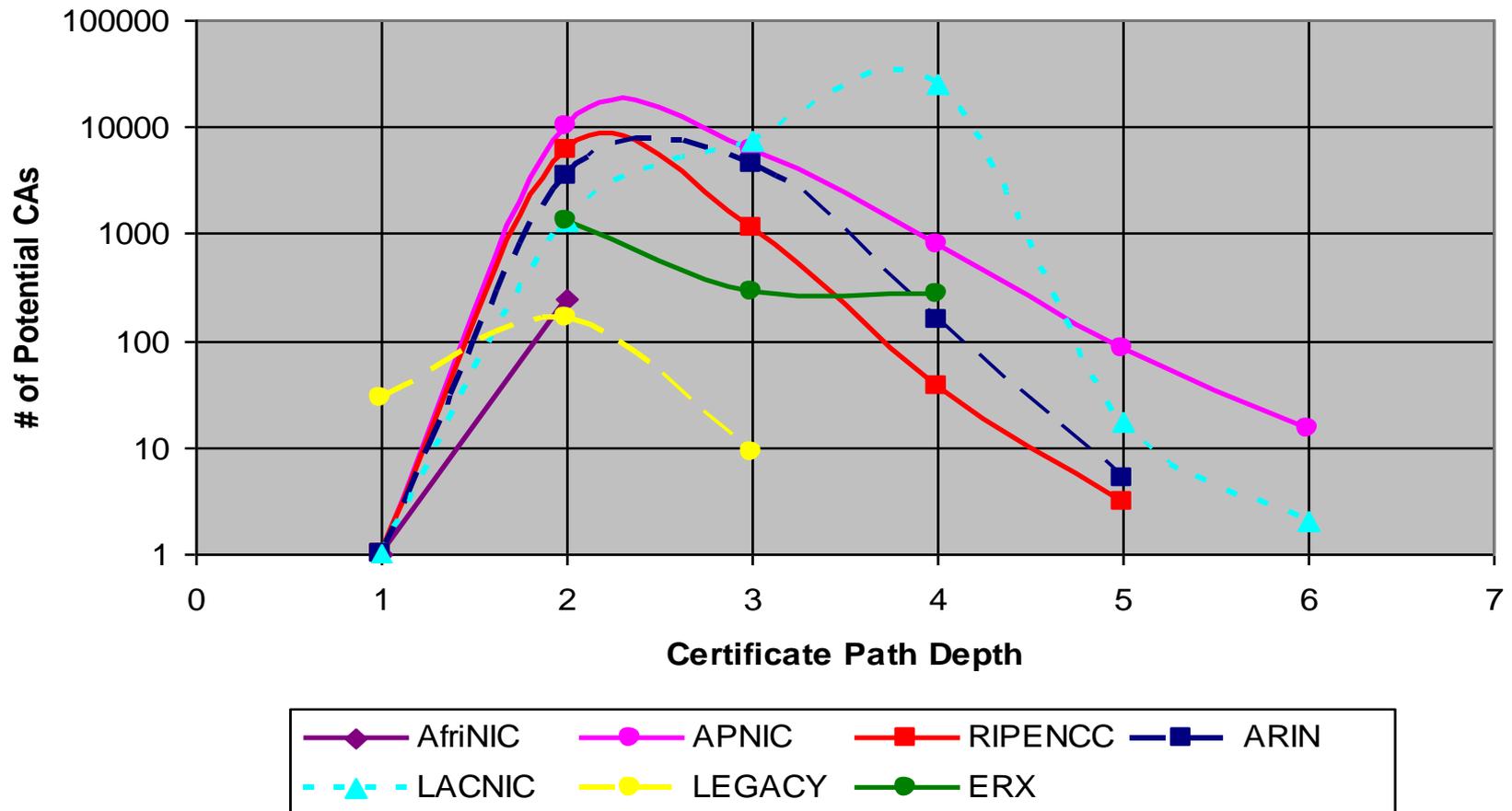
- Selection criteria:
 - Resource allocation objects:
 - * inetnums in RPSL.
 - * NetHandles in SWIP.
 - Attributes contained in an object to identify the allocation type:
 - * “status:” in inetnum.
 - * “NetType:” in NetHandle.
 - Status/NetType Attribute values: Allocation, Re-allocation
 - * First consider “Allocation” ONLY (including both PA and PI)
 - * Then consider “Allocation” and “Re-allocation”
 - Five top level CAs: ARIN, RIPE NCC, APNIC, LACNIC, AfriNIC in addition to IANA
 - * For blocks with prefix length ≤ 8 , the certificates are created by the RIRs
 - * For these blocks, the RIRs are the CAs
 - Eliminate also objects whose size < 255 (i.e, more specific than /24)
- Algorithm for selecting potential CAs:
 - Legacy:
 - * If Org of an object is uniquely defined and the object is either
 - Direct assignment (/8) to an organization; OR
 - Allocation to an ISP under Legacy space (e.g., 4/8 and 8/8 are allocated to Level 3 Comm).
 - Regular allocations and ERX:
 - * If Org of an object is uniquely defined AND the object is allocation (or, reallocation) regardless of the allocation depths

Distribution of Potential CAs per RIR



- # of potential global CAs (allocations only): ~22.4K
- # of potential global CAs (alloc + realloc): ~69.1K
- Note that AfrINIC, APNIC and RIPE NCC do not have the value “re-allocation”. Hence, the first level of direct allocations by these RIR is considered as “Allocation Only”.
- Note also that some objects do not contain “org:” attribute, especially for the regions such as RIPE NCC and APNIC.

Distribution of Certificate Path Depths of potential CAs (Alloc + Realloc)



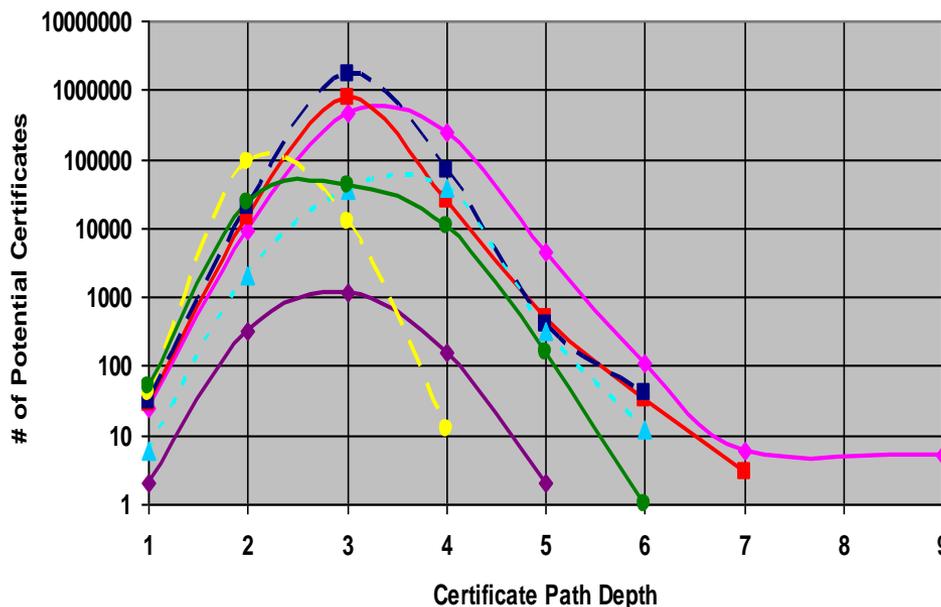
- LACNIC, LEGACY and ERX Data are selected from both RPSL and SWIP excluding duplicates.
- Certification path depth “1” indicates the top-level allocations by IANA to RIRs, i.e., address blocks $\geq /8$.

Analysis of IPv4 certificates

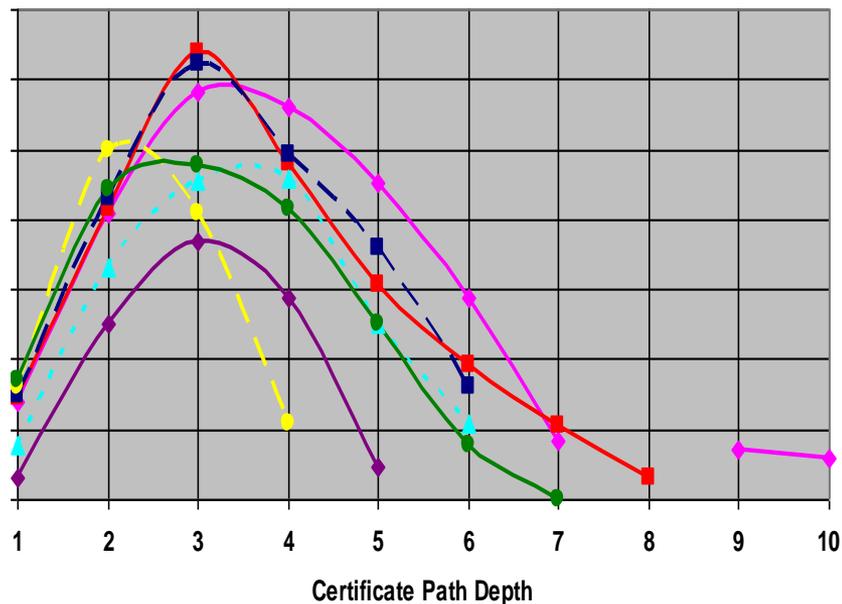
- Full deployment vs. optimized deployment:
 - **Full deployment**: if it was currently deployed based on the registry allocation data.
 - **Optimized deployment** after IPv4 prefix optimization:
 - * Aggregation of adjacent equal length prefixes
- Algorithm for IPv4 prefix optimization:
 - For every possible aggregate (i.e., two adjacent, equal sized, aggregatable prefixes), check the following attributes:
 - If **organizations** in the two objects are defined and the same, aggregate the two.
 - Else if **organizations** in the two objects are defined but different, do not aggregate the two.
 - Else if both or either one of the two contain no organization, then:
 - * If both **country code** and **status (e.g., PI vs. PA and allocation vs. assignment)** between the two are the same:
 - Check **mntner-related attributes** (i.e., mnt-by, mnt-lower, mnt-routes) between the two.
 - If check passes, then aggregate the two.
- Create a new aggregate, if no existing prefix for the aggregate exists, as follows:
 - Aggregated by org:
 - * generate a new aggregate with the **org/status** values of the **first prefix** without mnt values.
 - Aggregated by mnt:
 - * Generate a new aggregate with the **country/status/mnt** values of the first prefix excluding org.

Distribution of IPv4 Certificate Path Depths

Optimized deployment



Full deployment



- LACNIC, LEGACY and ERX data are selected from both RPSL and SWIP excluding duplicates.
- Prefix length "0" indicates that an address block cannot be represented by a single CIDR prefix.
- Certification path depth "1" indicates the top-level allocations to RIRs by IANA, i.e., address blocks $\geq /8$. Each " $\geq /8$ " block is counted separately.

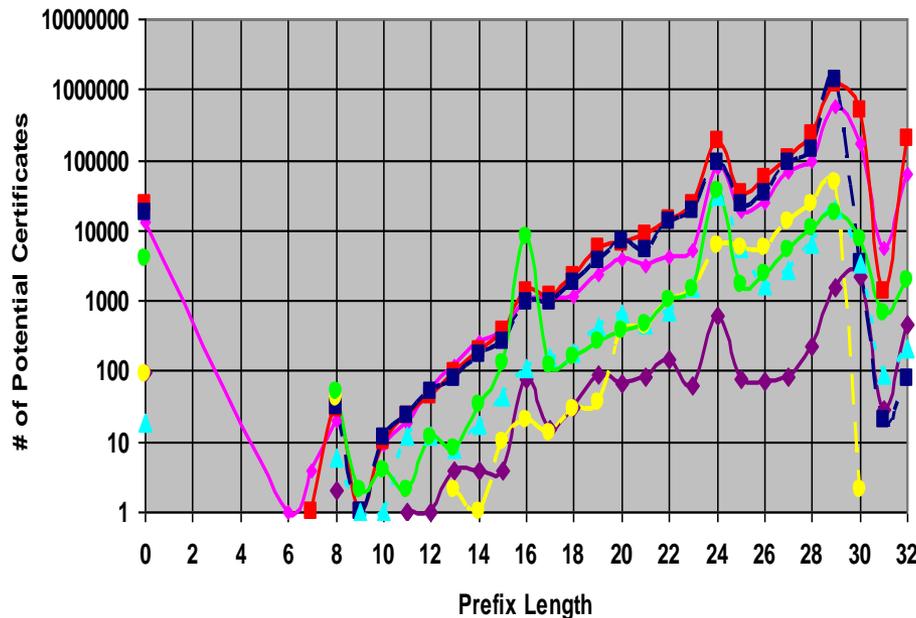
Improvement from optimization for IPv4 Certificates

	All objects			# objects with prefix length <= /24			# objects with prefix length >= /25		
	Full deployment	Optimized deployment	Reduction rate	Full deployment	Optimized deployment	Reduction rate	Full deployment	Optimized deployment	Reduction rate
RPSL	3,733K	1,598K	57%	385K	245K	36%	3,311K	1,316K	60%
SWIP	1,829K	1,816K	0.7%	145K	137K	6%	1,667K	1,662K	0.3%
LEGACY/ERX	207k	178K	14%	59K	48K	19%	144K	126K	13%
Global	5,769K	3,592K	38%	589K	430K	27%	5,122K	3,104K	39%

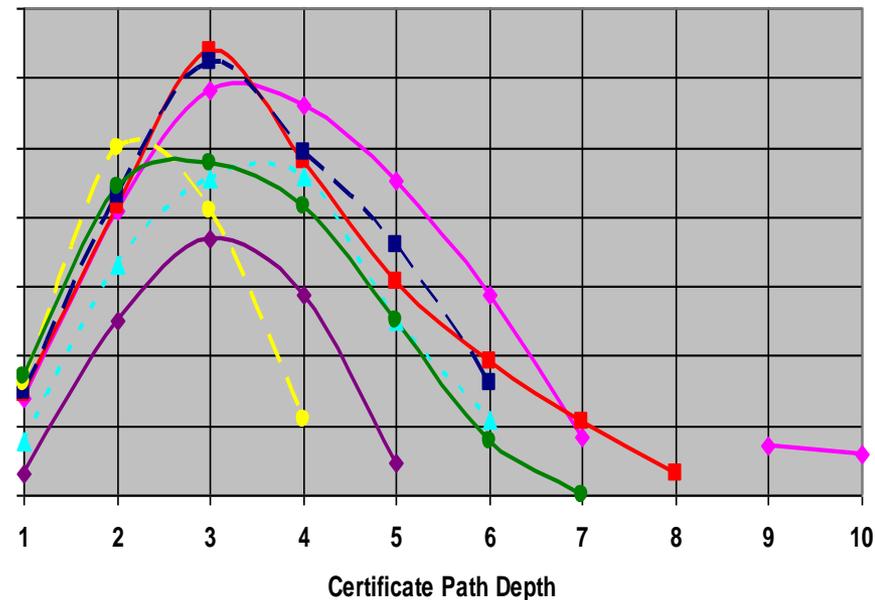
Prefixes with prefix length NULL are not included in this table.

Distribution of Prefix Lengths vs. Certificate Path Depths of IPv4 (full deployment)

Prefix length



Certification path depth



- LACNIC, LEGACY and ERX data are selected from both RPSL and SWIP excluding duplicates.
- Prefix length “0” indicates that an address block cannot be represented by a single CIDR prefix.
- Certification path depth “1” indicates the top-level allocations to RIRs by IANA, i.e., address blocks $\geq /8$. Each “ $\geq /8$ ” block is counted separately.

IPv4 Non-contiguous (Overlapping) Sub-allocations in RPSL (examples)

- RIPE:

- 62.128.192.0 – 62.128.207.255
- 62.128.195.0 – 62.128.223.255

- APNIC:

- 211.100.249.184 – 211.100.250.191
- 211.100.249.192 – 211.100.250.199
- 211.100.249.200 – 211.100.250.207
- 211.100.249.208 – 211.100.250.215
- 211.100.250.216 – 211.100.250.225



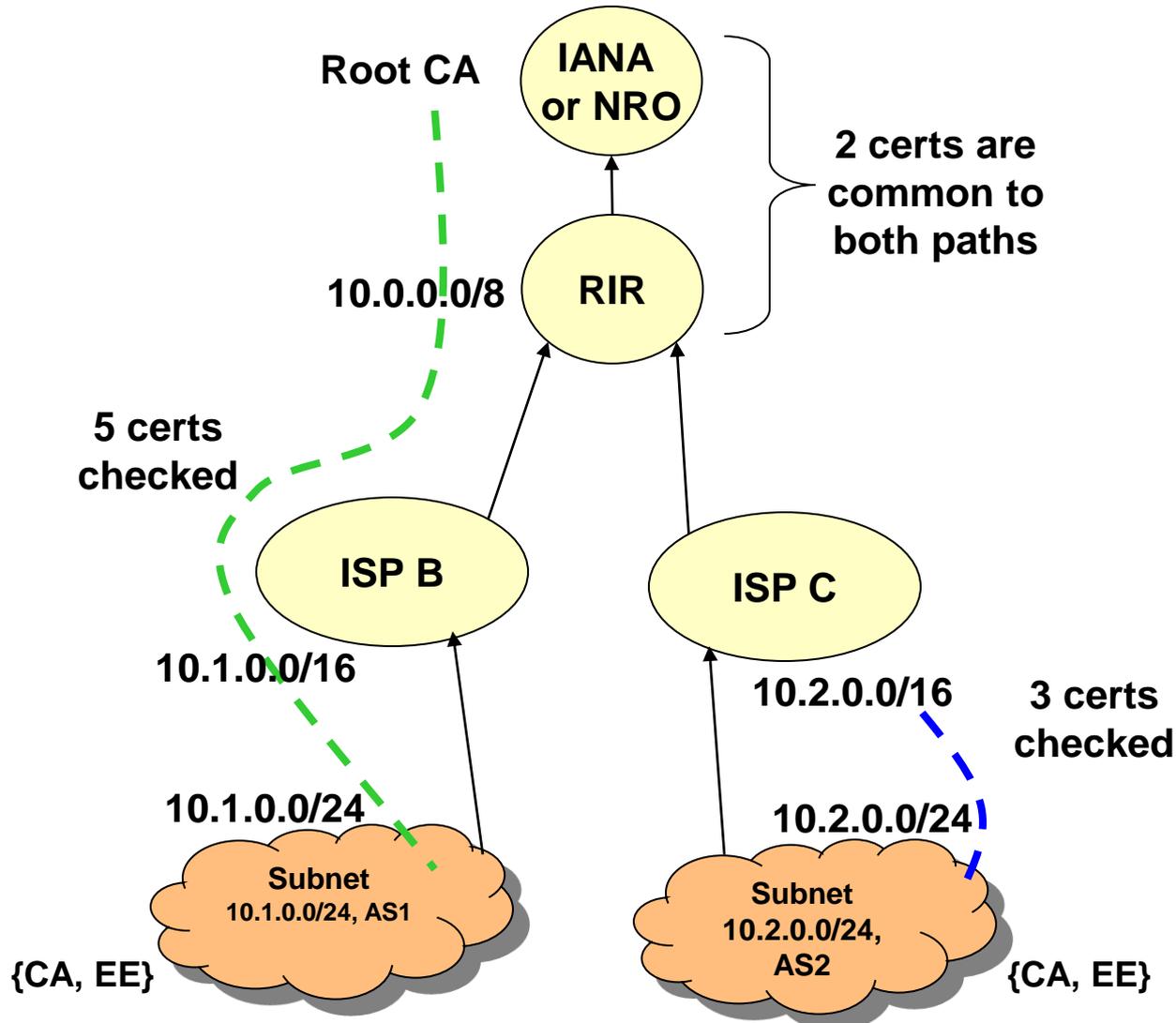
Analysis of ROAs

- The cost estimate of ROA verifications in terms of certification path lengths
- Distribution of PI address space
- Distribution of MOASes of potential ROAs

Analysis of ROAs

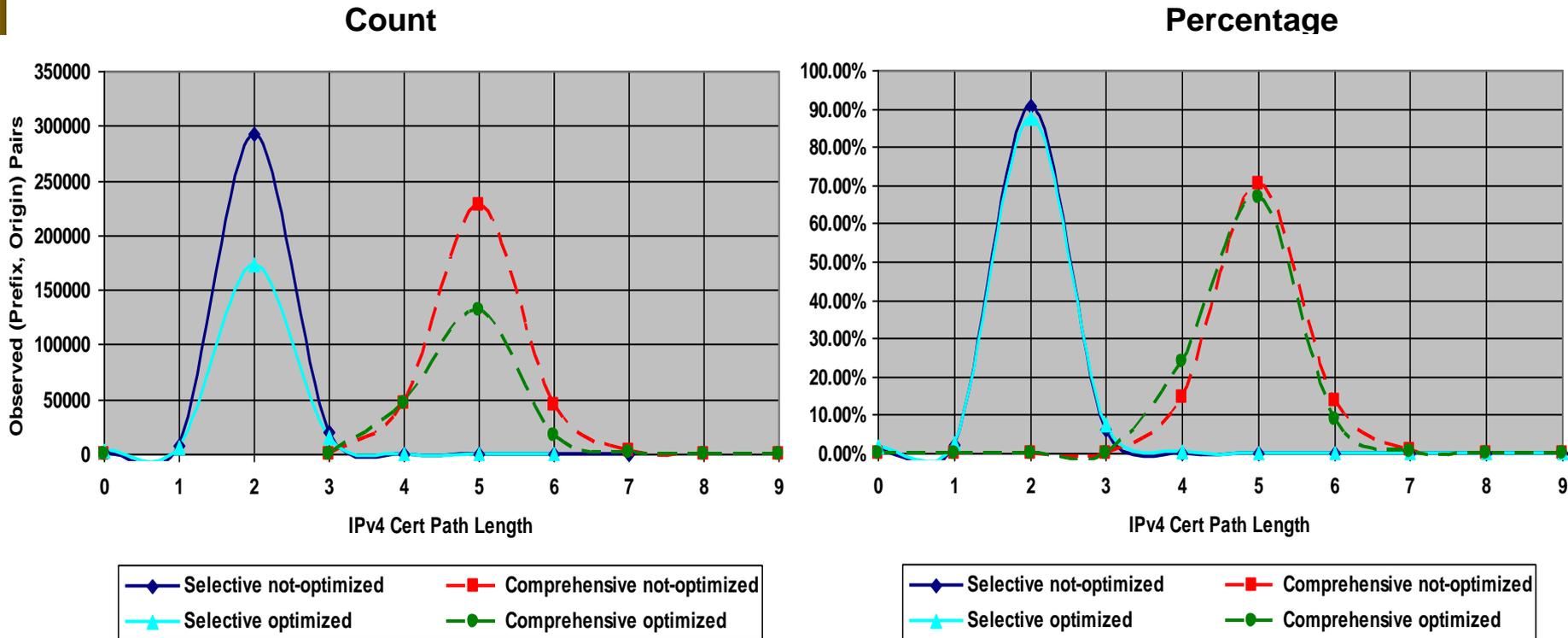
- **ROA analysis techniques:**
 - ROA prefix optimization with the same AS:
 - * *Not optimized:* full-scale
 - * *Optimized:* Aggregation of adjacent equal length prefixes with the same Origin AS
 - ROA prefix verification optimization:
 - * *Comprehensive:*
 - Check every single resource certificate in a certification path including a root.
 - * *Selective:*
 - Use “*validation state*” of a certificate to avoid redundant checks on the certificates that have already been checked.
- **Categorization of the ROA verification:**
 - *Comprehensive and not-optimized*
 - *Comprehensive and optimized*
 - *Selective and not-optimized*
 - *Selective and optimized*
- **Method for computing the length of a certification path:**
 - Does the prefix of a potential ROA have an exact match resource allocation record?
 - * If yes, then that object is considered as a CA and assume an EE for the prefix is created.
 - * If not, then assume both a CA and an EE for the prefix are created.
 - Assume also that routes with prefix length ≥ 25 have only the corresponding EEs, not CAs.
 - Compute the number of certificates included in a particular certification path for the EE including a root certificate and a target EE.
 - *IANA or NRO (the top-level entity) is assumed to be a single trust anchor for this analysis.*

Optimization in ROA Prefix Validation: Selective Method



Distribution of Certification Path Lengths for ROA Prefix Validation

BGP trace data date: February, 2009



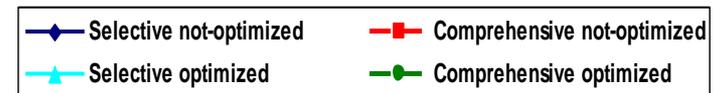
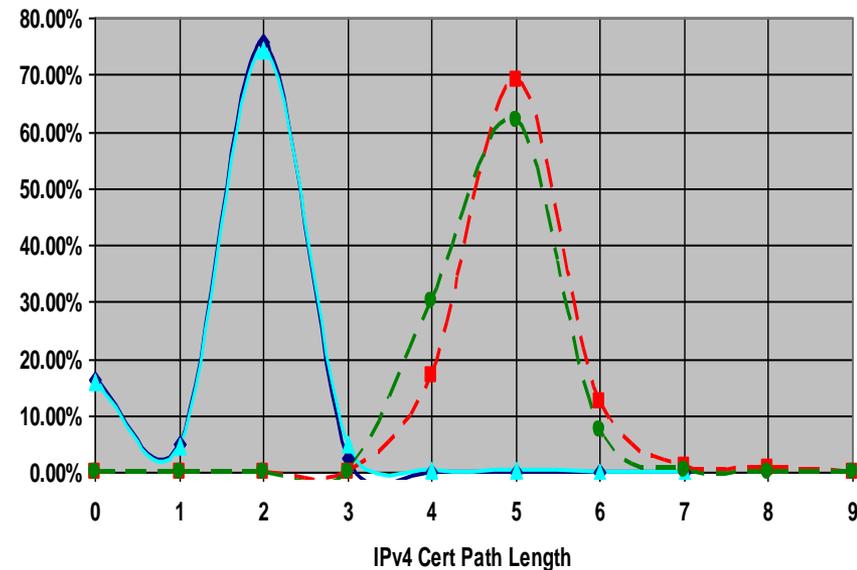
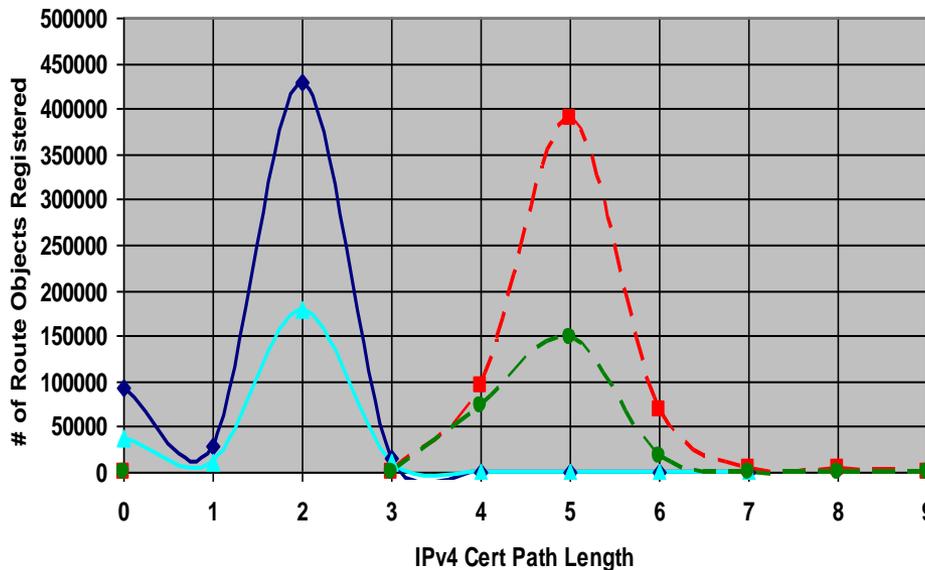
In the case “*Selective and not-optimized*”, a realistic scenario for the global-scale deployed RPKI, the average cert. path length for IPv4 address is ~2.03. About 93.6% of observed (P,O) pairs need to verify about two or less IPv4 address certificates for the prefix of a route.

Distribution of Certification Path Lengths for ROA Prefix Validation

RPSL data date: 2009-02-18

Count

Percentage



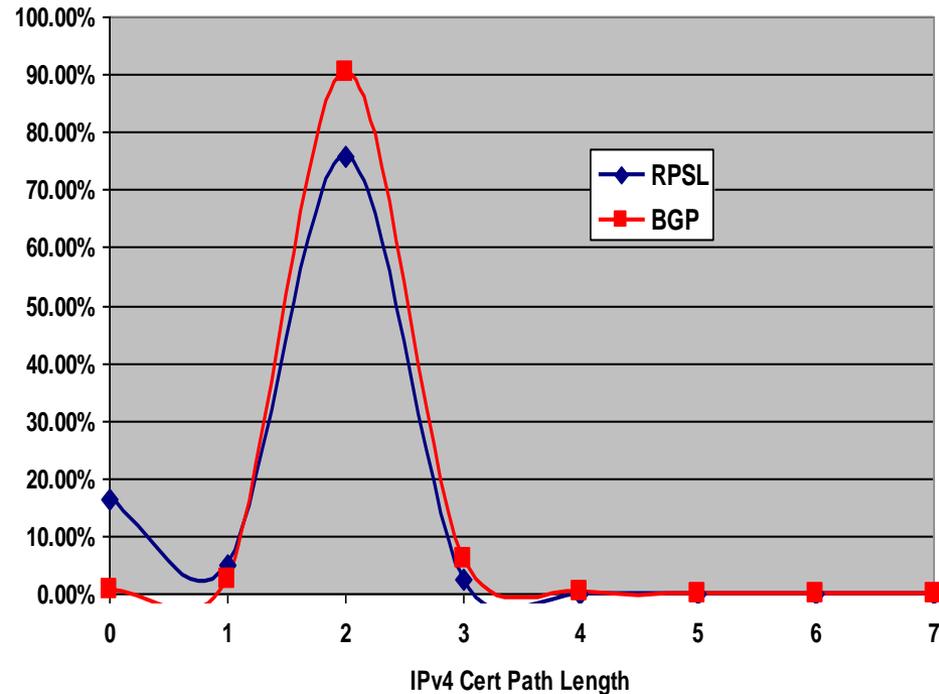
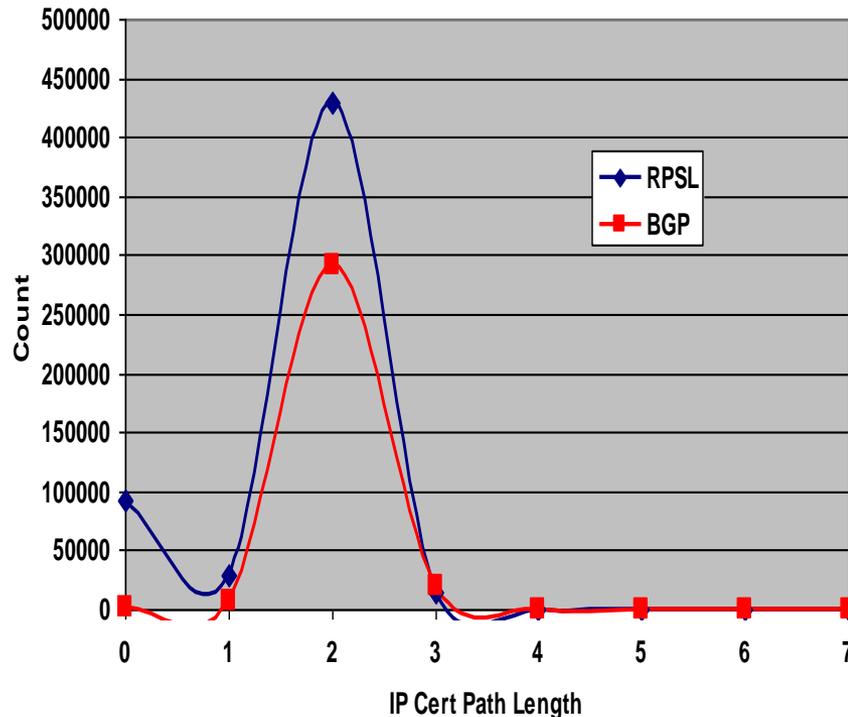
In the case “*Selective and not-optimized*”, the average cert. path length for IPv4 address is ~1.7. About 81% of registered route objects need to verify two or less IPv4 address certificates for the prefix and about 16% need not verify the prefix of a route at all (due to multi-homed prefixes).

Distribution of Certification Path Lengths for ROA Prefix Validation

selective and not-optimized (RPSL vs. BGP Trace)

Count

Percentage



These graphs depict that the two data sources show similar behavior, i.e., the majority of ROAs (94% for BGP and 97% for RPSL) need to check only 2 or less IPv4 address certificates for ROA validation.

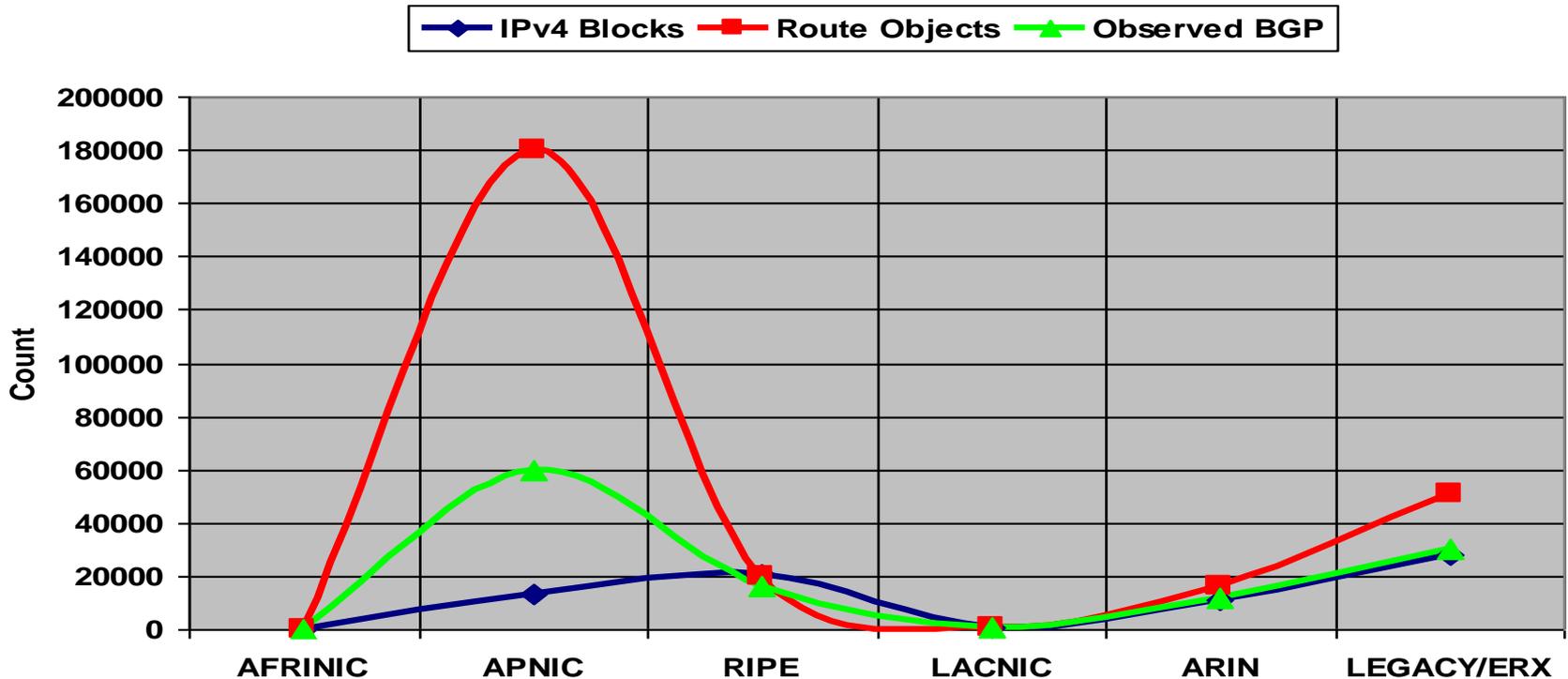
Analysis of PI Space in RPKI Issues

- Attributes “**status:**” in inetnum and “**NetType:**” in NetHandle:
 - Specify the type of address range represented by the address allocation object.
- No globally defined values of these attributes across RIRs. The defined values for PI blocks are as follows:
 - RIPE / AFRINIC:
 - * *ALLOCATED PI / ASSIGNED PI / LIR PARTITIONED PI.*
 - APNIC:
 - * *ALLOCATED PORTABLE / ASSIGNED PORTABLE.*
 - * All /8 blocks are defined as *ALLOCATED PORTABLE.*
 - Some LEGACY blocks are contained in both RPSL and SWIP.
- The LEGACY/ERX blocks are generally assumed to be PI. However, some LEGACY/ERX blocks are specifically defined as PA. These specifically defined PA blocks are excluded for PI analysis.
- Some inetnum objects (in RPSL) do not contain “**status:**” attribute at all:
 - *# of inetnums with no “status:”: 490,661.*
 - *Almost all of these came from JPNIC (one of NIRs under APNIC): 490,559*

Analysis of PI Space in RPKI Methodology

- Select IP resource allocation objects with PI specification.
- Adapt a different approach to each RIR:
 - RIPE / AFRINIC:
 - * *All inetnum objects with the locally defined values for PI (ALLOCATED PI, ASSIGNED PI, LIR PARTITIONED PI).*
 - * */8 blocks are defined as ALLOCATED UNSPECIFIED.*
 - APNIC:
 - * *All inetnum objects with the locally defined values for PI (ALLOCATED PORTABLE, ASSIGNED PORTABLE).*
 - * */8 blocks are defined as ALLOCATED PORTABLE, which are excluded.*
 - ARIN / LACNIC:
 - * All objects that are directly “ASSIGNED” to an organization by the RIR.
 - LEGACY/ERX:
 - * First, select all NetHandle objects with PI from SWIP, which belong to LEGACY/ERX.
 - * Then, select all the LEGACY/ERX inetnum objects with PI from RPSL, which are not included in SWIP.
- Classify these PI blocks based on IANA allocation registry.

Distribution of PI Address Blocks based on Allocation Source



The graph depicts that APNIC-allocated PI address blocks are heavily sub-allocated to both route objects and advertised BGP updates.

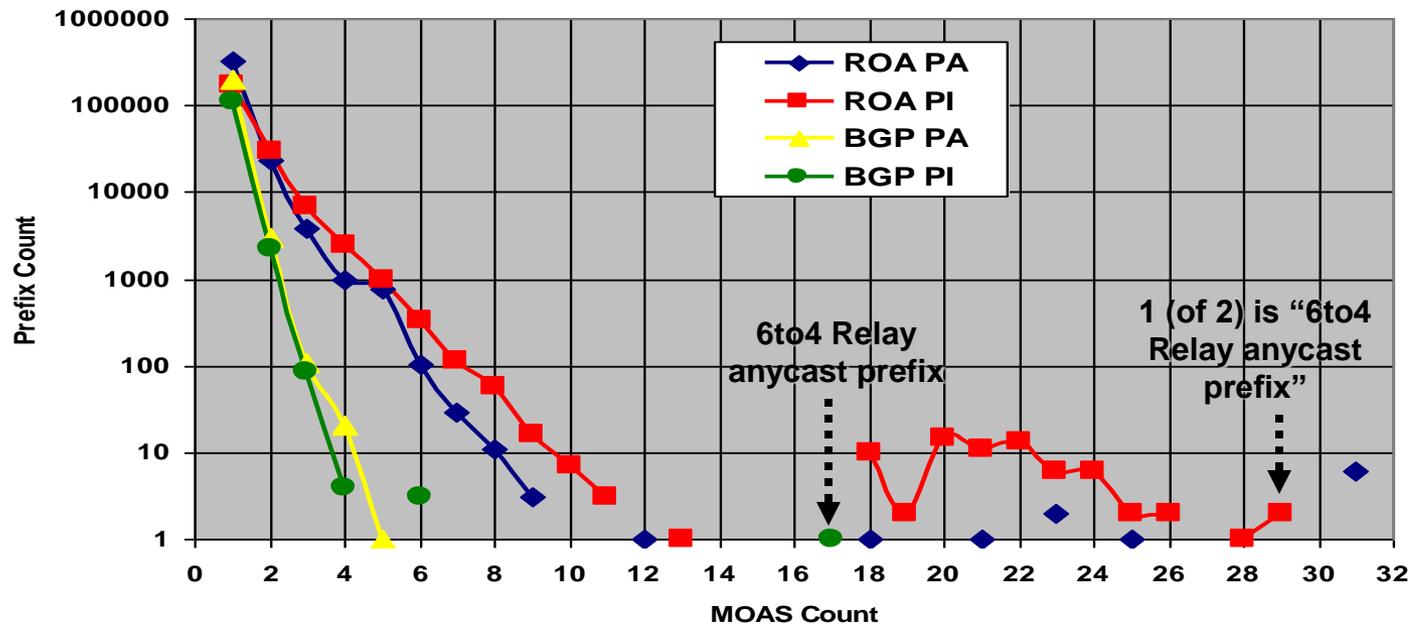
- # IPv4 blocks with the valid "status": 5,281K
- # IPv4 blocks with NULL "status": 491K
- # IPv4 blocks with PI: 74K (~1.4%)

- # route objects (RPSL + SWIP): 654K
- # route objects with PI: 268K (~41%)
- # objects in both RPSL and SWIP: 4K
- There may be many proxy-registered route objects.

- # observed (P, O) pairs: 322K
- # observed (P, O) pairs with PI: 118K (~37%)

Distribution of MOASes of Route Objects (ROA) and Observed BGP Updates (BGP) with PI vs. PA

Registry data date: 2009-2-18



- Here PA means the rest of address blocks other than PI space in the registry.
- PI address blocks tend to have more MOASes, especially in route objects. Does this indicate that many of them could be proxy-registered route objects or stale objects?

- # globally unique route objects (RPSL + SWIP): 654K
- # globally unique route objects with PI: 268K (~41%)
- # multi registrations between RPSL and SWIP: 4K
- There may be many proxy-registered route objects.

- # of observed unique (P, O) pairs: 322K
- # of observed unique (P, O) pairs with PI: 118K (~37%)

Conclusions

- We performed quantitative analysis of potential deployed RPKI and compared two possible deployment scenarios: full vs. optimized deployments
 - The total number of IPv4 certificates can be significantly reduced with prefix aggregation.
 - The global reduction rate of the total number of IPv4 certificates is ~38%, and ~26% on the certificates with prefix length $\leq /24$.
- ROA validation in RPKI may not be a big performance issue:
 - About 89% of the total number of IPv4 address certificates (as of 2/18/2010) are address blocks with prefix length $\geq /25$, which may not call for ROA creation.
 - The performance of ROA verifications can be significantly improved by the use of the cached “validation state” of certificates being verified.
 - * About 933K IPv4 certificates among total of more than 5.8M need to be verified for ROA verification when used with existing route objects.
- Handling of partial allocations across multiple RIRs?
 - Who would be responsible for creating resCerts for LEGACY/ERX address blocks?
- **Future tasks:**
 - Analysis of RPKI growth over time
 - Potential impact of RPKI on global BGP dynamics:
 - * The effect of creation, expiration or revocation of resource certificates and ROAs
 - The models can help generate synthetic RPKI workload models for routers for origin / path validation