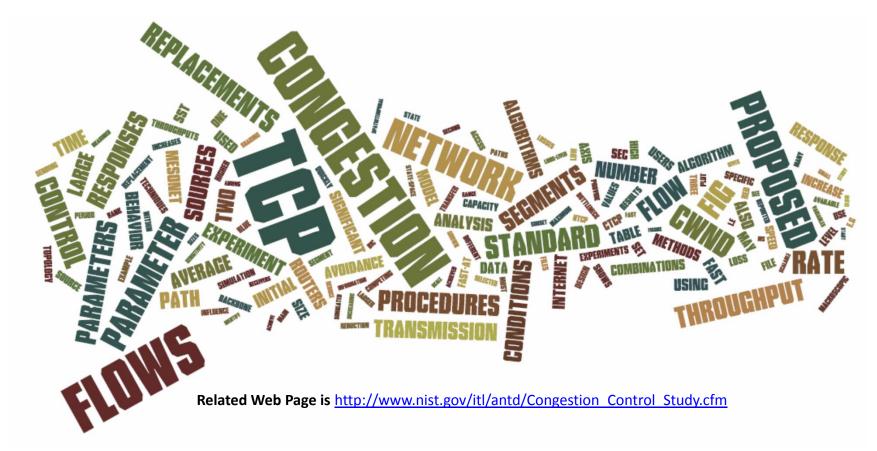
#### Predicting Macroscopic Dynamics in Large Distributed Systems



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ASME PVP 2011 July 21, 2011

We are investigating methods to predict macroscopic dynamics in large distributed systems.

- Modern networking technology allows components in large distributed systems to be updated easily.
- Will any given update radically change macroscopic system dynamics?
- Can we predict how such updates will influence systemwide behavior?

Here, we demonstrate some of our methods applied to predict how various proposed replacements for the transmission control protocol (TCP) will affect macroscopic dynamics in the Internet.

# Outline

- Introduction to the Problem
- Methods to Reduce Stimulus and Response State Spaces
- Methods to Analyze Multidimensional Datasets
- Application to Predict Macroscopic Dynamics of Proposed TCP Replacements

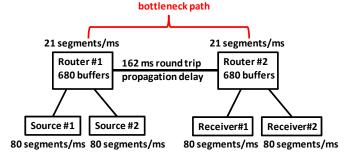
## Introduction to the Problem

#### **Internet Congestion Control & TCP**

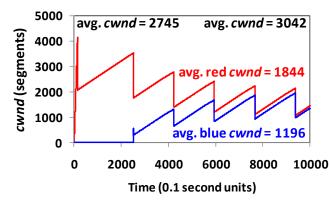
#### Phases in a TCP Flow



#### 252 Mbps Bottleneck Path



#### TCP Flows Achieve 89 % Utilization



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#### Utilizations Much Lower on High-Speed, Long-Delay Paths

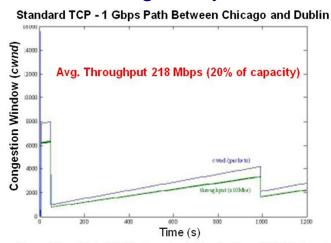
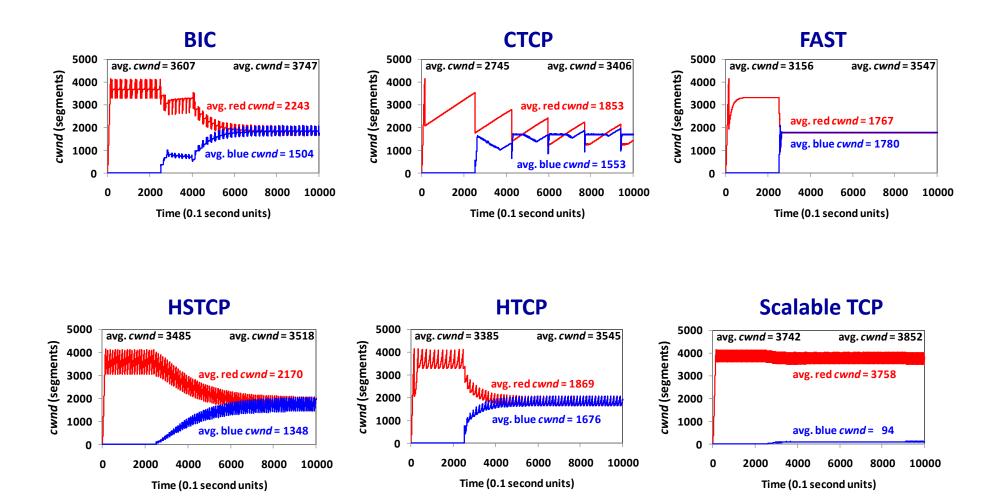


Figure 1 from Li et al. 2007. Experimental Evaluation of TCP Protocols for High-Speed Networks. *Transactions on Networking*. 15:5, 1109-1122.

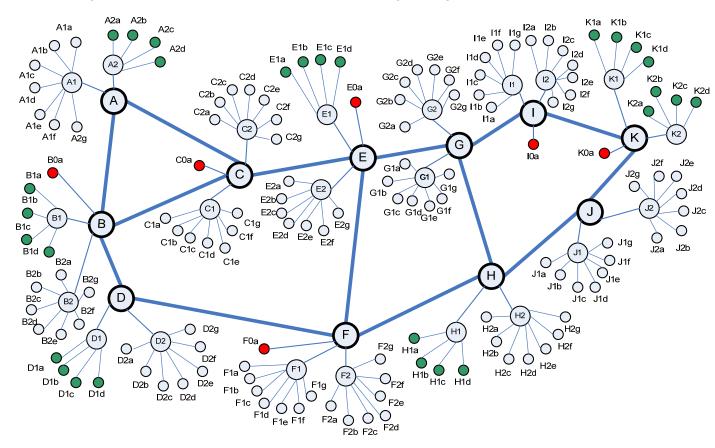
# So Researchers Proposed TCP Replacements

Identifier	Acronym	Name of Algorithm
1	віс	Binary Increase Congestion Control
2	CTCP	Compound TCP
3	FAST	Fast Active Queue Management Scalable TCP
4	HSTCP	High Speed TCP
5	HTCP	Hamilton TCP
6	STCP	Scalable TCP
7	ТСР	Transmission Control Protocol

#### Proposed TCP Replacements Well-Studied on Bottleneck Path



# What would be the macroscopic dynamics if these proposed TCP replacements were deployed in the Internet?



Four-Tier Topology (based on the Abilene Network) with 11 Backbone Routers (A-K), 22 Point of Presence Routers (A1-K2) and 139 Access Routers (A1a-K2d) – 6 red and 28 green Access Routers may operate at different speeds from the 105 others – not shown is a fourth tier of hundreds of thousands of attached computers

## Stimulus and Response State Spaces Are Large for Network Simulation Models

$$y_1, ..., y_z = f(x_{1|[1,...,\ell]} ..., x_{p|[1,...,\ell]})$$
Response State-Space Stimulus State-Space
e.g.,  $z = 500$  e.g.,  $p = 1000$  and  $\ell = 2^{32}$ 

# Methods to Reduce the Scale of the Stimulus and Response Spaces

# We Employ Stimulus State-Space Reduction Techniques

**Sensitivity** 

**Analysis** 

#### X4 Topology Web Object Size for Browsing X5 **Proportion & Size of Larger Files** X6 $(2^{32})^{1000} \longrightarrow O(10^{9633})$ [10<sup>80</sup> = atoms in visible universe] **X7** Think Time User **Behavior X8 Patience Selected Spatiotemporal Congestion** X9 Discard parameters not germane to study-reduce by 944 parameters X10 Long-lived Flows X11 Source & Receiver Interface Speeds $(2^{32})^{56} \longrightarrow O(10^{539})$ X12 Number of Sources & Receivers Sources & Receivers X13 Distribution of Sources Group related remaining parameters reduce by 36 parameters X14 **Distribution of Receivers Parameter** X15 **Congestion Control Procedures** $2^{32}$ ) $^{20} \longrightarrow O(10^{192})^{11}$ **Protocols** X16 Initial Congestion Window Reduction X17 **Initial Slow Start Threshold** Select only 2 values for each parameter Simulation X18 **Measurement Interval Size** Control & Level X19 Simulation Duration $2^{20} \longrightarrow O(10^6)$ Measurement X20 Startup Pattern Reduction Use experiment design theory to reduce parameter combinations to 256

Design July 18-19, 2011

**Orthogonal** 

Design

**Fractional-Factorial** 

(OFF) Experiment

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**OFF** 

**Experiment** 

parameter combinations to 32

Use sensitivity analysis

significant parameters

to identity six most

Use experiment design theory again to reduce

#### **NIST MesoNet Simulator**

**Network Speed** 

**Propagation Delay** 

**Buffer Provisioning** 

Identifier Name
X1 Netwo

**X2** 

**X3** 

Category

Network Configuration

# We Employ Response State-Space Reduction Techniques (here we demonstrate correlation analysis & clustering)

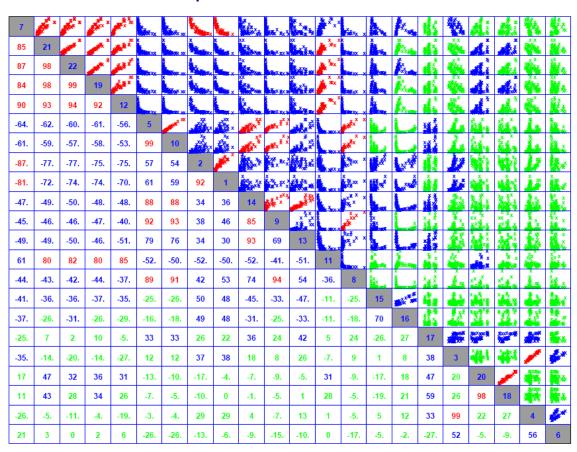
Response	Definition
v1	Number o

y1	Number of sources transmitting
y2	Proportion transmitting [y1 / all sources]
у3	Number segments entering network
y4	Number of segments leaving network
у5	Loss Rate [y4/ (y3 + y4)]
у6	Flows completing per interval
у7	Flow-completion rate [y6/(y6+y1)]
у8	Connection failures per interval
у9	Connection-failure rate [y8/(y8+y1)]
y10	Retransmission rate
y11	Congestion window (cwnd)
y12	cwnd increases per interval
y13	Duplicate ACKs per interval
y14	Timeouts per interval
y15	Round-trip time (rtt)
y16	Queuing delay

#### Response Definition

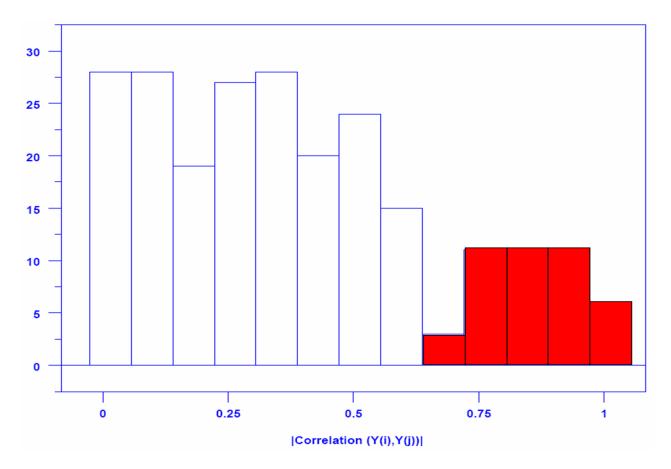
y17	Average Segments/Second on DD flows
y18	Average Segments/Second on DF flows
y19	Average Segments/Second on DN flows
y20	Average Segments/Second on FF flows
y21	Average Segments/Second on FN flows
y22	Average Segments/Second on NN flows

#### **Response Correlation Matrix**



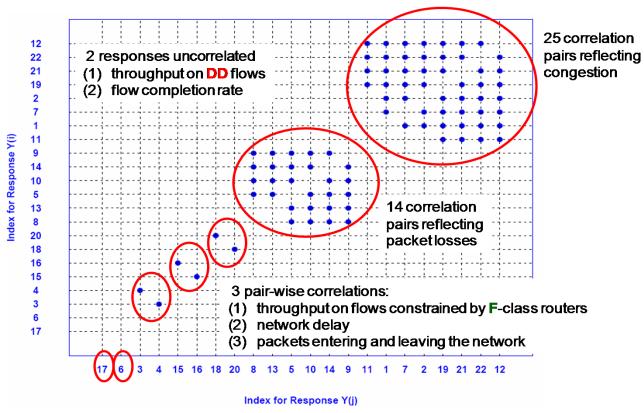
Red  $80 \ge |r| \times 100 \le 100$ Blue  $30 \ge |r|x100 < 80$  Green |r|x100 < 30

#### Decide Which Pairs of Correlated Responses to Analyze



Frequency distribution of |r| for pair-wise correlations – bins highlighted for |r| > 0.65

#### **Cluster Retained Correlated Responses**



Index-Index Plot Sorted by Increasing Count of Correlated Pairs to Indentify Clusters of Mutual Correlations that Represent Seven Behavioral Dimensions

# Methods to Analyze Multidimensional Datasets

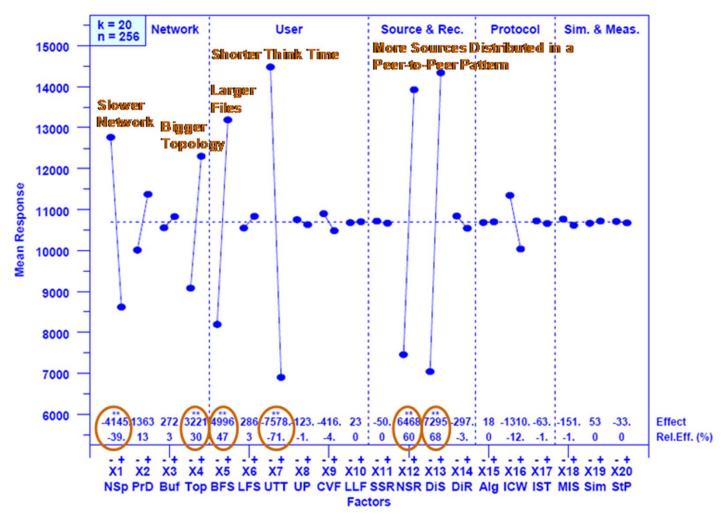
# Even with Response State-Space Reduction Combinatorial Comparisons Can Yield Large, Multidimensional Datasets

For example, comparing 7 congestion control algorithms with respect to 48 responses under 32 conditions over three time periods with a mixture of competing TCP flows yields a multidimensional dataset with > 64,000 numbers

We used a range of multidimensional data analysis techniques, including:

- 1. Main-Effects Analysis [a, b, c]
- 2. Cluster Analysis [a, b]
- 3. Primary Principal Components Plots
- 4. Condition-Response Summaries [a, b, c]
- 5. Biplots
- 6. Colored Bar Graphs
- 7. Rank Analysis

## 1(a) Sample Main-Effects Analysis on One Response



Main effects plot showing absolute and relative influence of each *MesoNet* parameter (x axis) on mean (y axis) number of sources transmitting (y1) and identifying six statistically significant parameters – network speed (X1), topology (X4), file size (X5), think time (X7) and number (X12) and distribution (X13) of sources

#### 1 (b) Summary of Main-Effects Analyses Across 18 Responses

Significance of influence of 20 MesoNet parameters (columns) on 18 macroscopic responses (rows) during  $2^{nd}$  time period: blue cells indicate significance p < 0.01 and orange cells indicate significance p < 0.05, where a - or + in highlighted cells indicates the parameter setting that causes an increase in the corresponding response

Metric		Network					User Behavior					Sou	urce/R	eceive	er	Protocol			Sim. Control & Meas.		
Class	Y#	X1	X2	Х3	Х4	X5	Х6	X7	Х8	Х9	X10	X11	X12	X13	X14	X15	X16	X17	X18	X19	X20
	Y1	-**		+**	+*			-**					+**	+**			+**				
	Y2	-**			+**	+**		-**					+**	+**							
Flows	Y3	+**	+**	+**		-**		+**					-**	-**				+**			
	Y4	-**	-**	-**		+**		-**					+**	+**		-*					
	Y5	+**		+**				+**					-**	-**		+**	+*	-**			
	Y6	-**	-**	-**		+**		-**					+**	+**			+**				
Congestion	Y7	+*																			
	Y8	-**	-**	-**		+**							+**	+**			+**				
Dalass	Y9	-**	+**	+**									+*	+*							
Delay	Y10	-**	+**	+**				-*					+**	+**							
Aggregate	Y11	+**	_*		+**	_**		-**					+**	+**							
TP	Y12	+**		+**	+**	+**		-**		-**			+**								
	Y13	+**		+*	+*					+**	-**										
Long-Lived	Y14	+*			+**					+**	_**										
Flow TP	Y15	+**								+**	-*				-*						
Other Ele	Y16	+**	_**		-**			+*		+**			-**		-*			+**	-*		
Other Flow	Y17	+**	-**		+*			+**		+**			-**	+**	-**		+**		-*		
ТР	Y18	+**	-**		-**			+**					_**	_**			+**				

1 (c) Summary of Parameter Influence on Responses for Three Time Periods and Two Congestion Control Algorithms

$$\Psi = (|\{y \mid p < 0.01\}| + \frac{1}{2} |\{y \mid p < 0.05\}|) / |\{y\}|$$

_	Network				User Behavior						Source/Receiver					Protocol			Control & Meas.		
Ψ x 100	X1	X2	Х3	Х4	<b>X</b> 5	Х6	Х7	X8	Х9	X10	X11	X12	X13	X14	X15	X16	X17	X18	X19	X20	
Period #1	95	53	64	50	50	0	67	0	0	17	0	75	67	14	6	44	14	8	0	0	
Period #2	94	53	53	42	39	0	61	0	33	14	0	75	64	11	8	31	17	6	0	0	
Period #3	97	56	61	42	53	0	69	0	25	17	0	75	67	11	6	28	17	11	0	0	
TCP	79	71	56	40	56	0	50	0	4	0	0	21	67	0	0	33	67	4	0	0	
CTCP	79	75	46	44	65	0	60	0	4	0	0	46	67	8	0	31	50	0	0	0	
Weighted Average	87	64	55	43	54	0	60	0	12	8	0	54	67	8	3	33	37	5	0	0	

Parameters with largest influence: network speed (x1), distribution of sources (x13), propagation delay (x2), user think time (x7), buffer size (x3) and file size (x5) and number of sources (x12)

#### 2 (a) Cluster Analysis

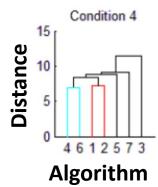
Compute normalized Euclidian distance between a pair of congestion-control algorithms in 45 dimensional space (45 responses used in this case)

$$Dist(Y_i, Y_j) = \sqrt{\sum_{m=1}^{45} \frac{(Y_{im} - Y_{jm})^2}{(\sigma_m)^2}}$$

Compute linkage between clusters of congestion-control algorithms

$$D(r,s) = \frac{1}{n_r n_s} \sum_{k=1}^{n_r} \sum_{l=1}^{n_s} Dist(Y_{k,r}, Y_{l,s})$$

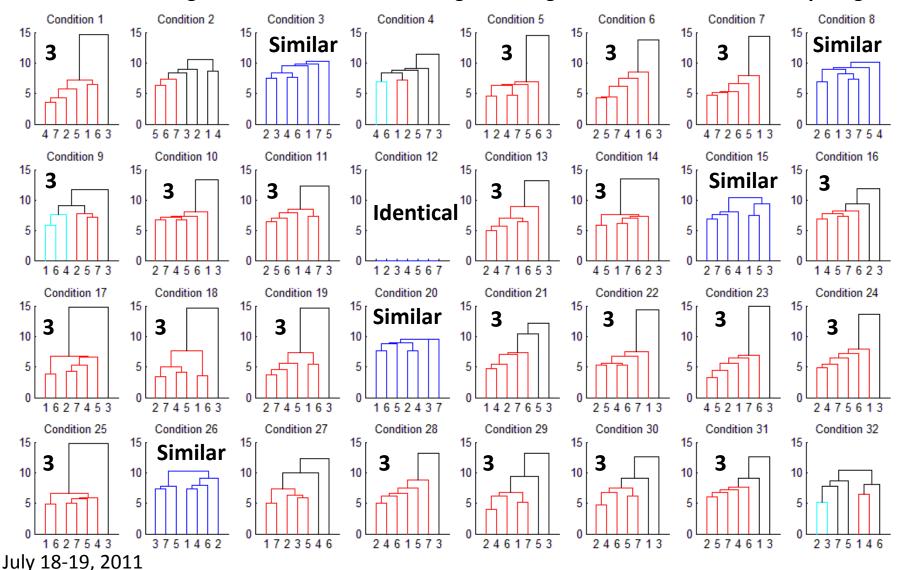
Plot as dendrogram and color when the linkage value for a pair of clusters falls below 70 % of the maximum linkage value



Note that congestion-control algorithms may only be compared under the same conditions (i.e., parameter combinations) because the conditions are very different from each other – this leads to the need to plot one dendrogram for each condition (32 here)

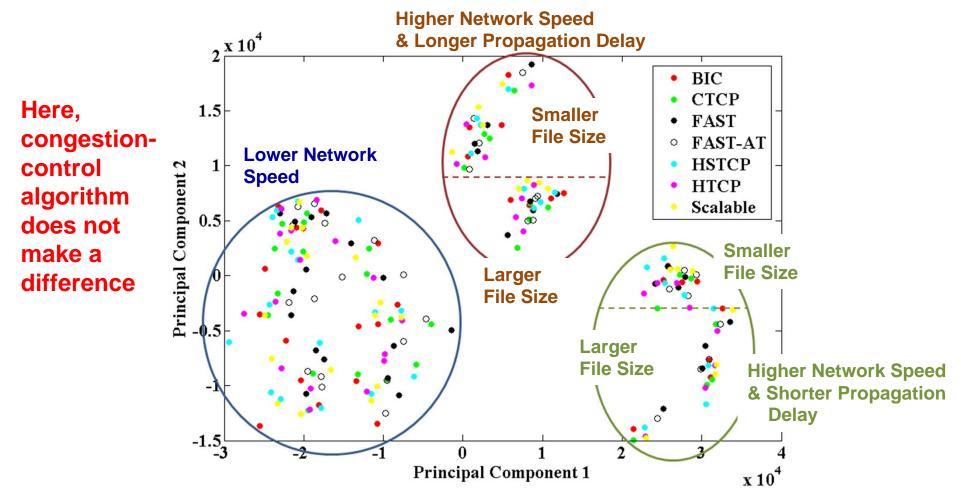
#### 2 (b) Dendrograms for Each of 32 Conditions

Identical under no congestion; Similar under low congestion; algorithm 3 distinct under heavy congestion



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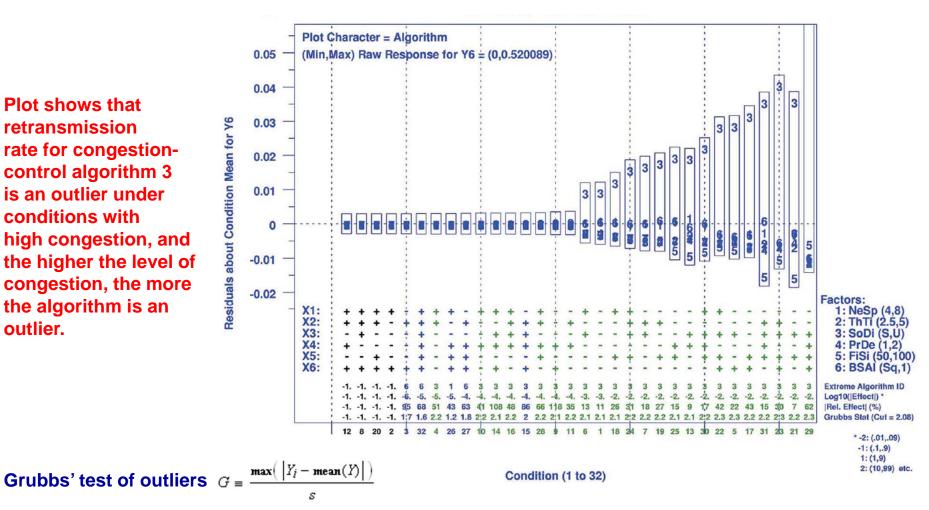
#### **3** Primary Principal Components Plots



PC1 (x axis) vs. PC2 (y axis) from a multivariate dataset containing average throughput achieved on network connections using seven different congestion control algorithms under 32 different parameter combinations. Circles identify groupings of points and lines within circles distinguish additional subgroups.

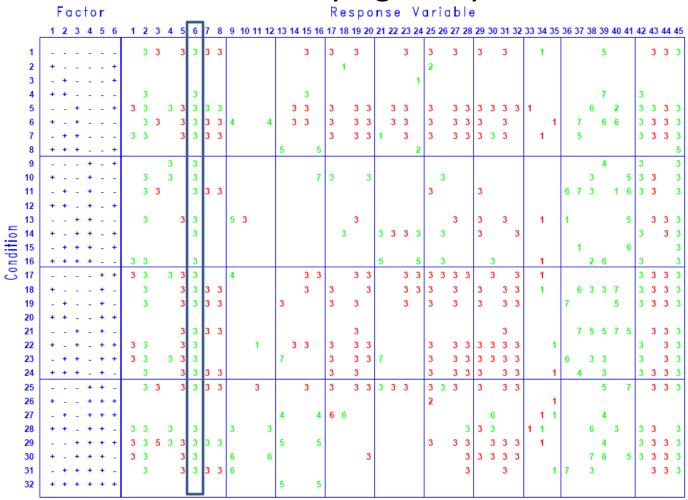
## 4 (a) Condition-Response Summaries: Underlying Analyses

Plot shows that retransmission rate for congestioncontrol algorithm 3 is an outlier under conditions with high congestion, and the higher the level of congestion, the more the algorithm is an outlier.



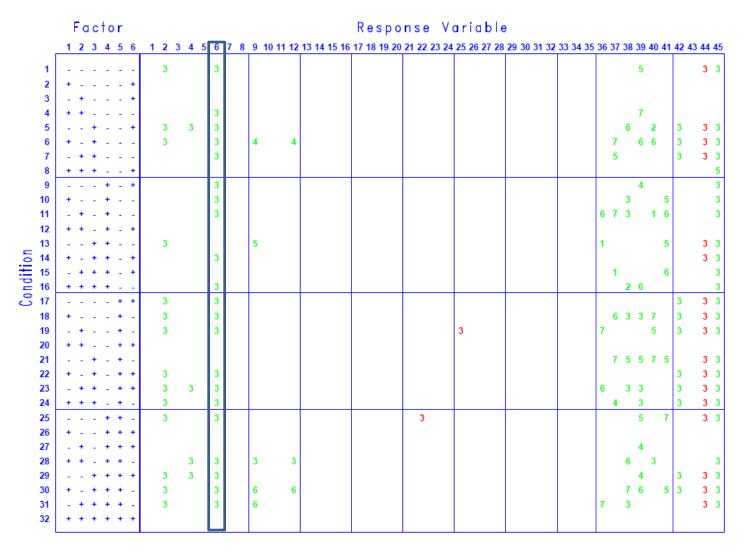
Sample plot analyzing the influence of condition and congestion control algorithm on the segment retransmission rate (in this experiment retransmission rate was designated as y6) – y axis gives residuals around the mean value for each condition and x axis gives conditions ordered by increasing range of residuals

# 4 (b) Condition-Response Summaries: Rollup Outliers from Underlying Analyses



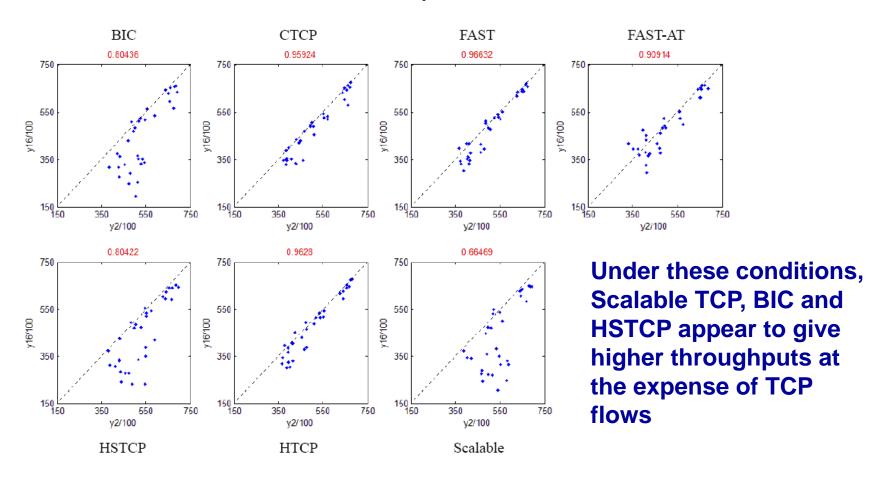
Condition-response summary identifying any statistically significant outliers among congestion control algorithms for each of 45 responses measured over 32 different conditions in this particular experiment

#### 4 (c) Condition-Response Summaries: Apply Filters



Filtered condition-response summary identifying any congestion control algorithms that are statistically significant outliers and that exhibit at least a 10 % relative effect for each of 45 responses measures over 32 different conditions

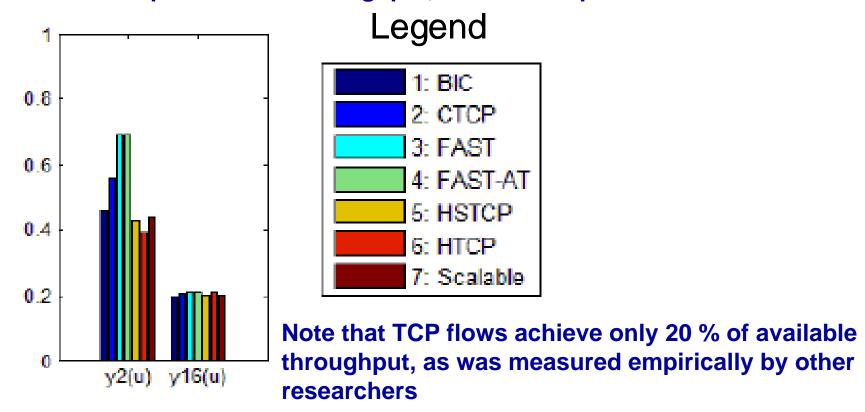
#### **5** Biplots



Seven biplots, where each biplot compares throughput (in units of 100 segments per second) for large flows transferred over very fast paths with fast interface speeds using a proposed TCP replacement (x axis) with throughput on identical competing flows using standard TCP – here a high initial slow-start threshold is used

#### **6** Colored Bar Graphs

#### FAST & FAST-AT provide best throughput, while CTCP provides second best

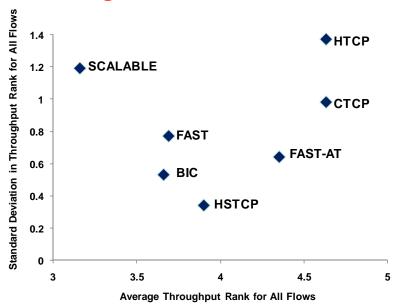


Fraction (y axis) of maximum available throughput achieved when transferring movies over uncongested paths with a maximum rate of 1 Gbps – each of the leftmost seven bars – y2(u) – represents flows using one of the proposed replacements for TCP (see legend), while each of the rightmost seven bars – y16(u) – represents flows using standard TCP and competing with flows using one of the proposed TCP replacements (see legend) – here a low initial slow-start threshold is used.

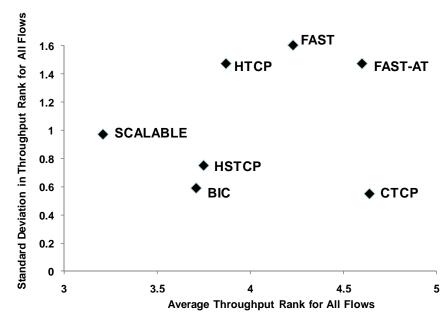
## 7 Rank Analyses

#### CTCP leads high ranking throughputs under a wide range of conditions

#### high initial slow-start threshold



#### low initial slow-start threshold



Average throughput rank (x axis) vs. standard deviation in throughput rank (y axis) for flows using each proposed TCP replacement and for competing flows using standard TCP

# Application to Predict Macroscopic Dynamics of Proposed TCP Replacements

## **Experiments Conducted**

How do the proposed TCP replacements react to and recover from spatiotemporal congestion?

	spatiotemporal congestion:
	Compared seven proposed TCP replacements (excluded
	FAST-AT) in a large (up to
	278,000 sources), fast (up to
	192 Gbps backbone) network;
Experiment #1a	Web browsing users; 25
	minutes simulated; three 5-
	minute time periods; large (> 1
	billion segments) initial sst; all
	sources use the same
	congestion control algorithm
	Same as #1a, except; smaller
	(up to 27,800 sources) and
Experiment #1b	slower (up to 28.8 Gpbs
	backbone) network and low
	(100 segments) initial sst.
	(Added FAST-AT)

How do the proposed TCP replacements improve flow throughputs and affect competing TCP flows?

tin ougnputo una	allect competing for nows:
	Compared eight proposed TCP
	replacements in a small (up to
	26,085 sources), slow (up to 38.4
	Gpbs backbone) network; Web
	browsing users and interspersed
Experiment #2a	users who download software and
	movies; 60 minutes simulated; large
	(> 1 billion segments) initial sst;
	some sources use standard TCP
	congestion control procedures and
	some sources use one of the
	proposed TCP replacements
Experiment #2b	Same as #2a except for low (100
Experiment #20	segments) initial sst.
	Same as #2a except for a larger (up
Experiment #2c	to 261,792 sources) and faster (up to
	384 Gbps backbone) network.

# What distinguishes among proposed TCP replacements? Will they be useful? Will they be safe?

- 1. Increase rate: How quickly can the maximum transmission rate be achieved?
- 2. Loss/Recovery processing:
  - a. How much does the protocol reduce transmission rate upon a loss?
  - b. How quickly does the protocol increase transmission rate after a reduction?
- **3.** TCP Fairness: How well do standard TCP flows do when competing with proposed TCP replacements?
- 4. Utility bounds: Under what circumstances can proposed TCP replacements provide improved user throughputs?
- **5.** Safety: Will widespread deployment of proposed TCP replacements induce undesirable macroscopic dynamics into the Internet?

#### Increase Rate

- Assuming low congestion, setting of initial sst is a key factor
  - High initial sst all algorithms (standard TCP included) achieved maximum transmission rate with the same (exponential) quickness
  - Low initial sst proposed TCP replacements achieved maximum transmission rate more quickly than (linear) increase of standard TCP
- Under heavy congestion, setting of initial sst matters little because initial slow start terminates upon first packet loss and a flow enters congestion avoidance, where loss/recovery processing determines throughput
- On real TCP flows receivers may convey a window (rwnd) that can restrict throughput because sources pace transmission based on min(cwnd, rwnd). Typically, rwnd < cwnd. In our studies, we assume an infinite rwnd in order to compare effects of congestion control algorithms. Throughput on many TCP flows in a real network might be constrained by rwnd, so that proposed TCP replacements would provide little advantage over standard TCP. In fact, even TCP congestion control does not have much influence when rwnd < cwnd.</p>

## Loss/Recovery Processing

- ➤ One group of algorithms (Scalable TCP, BIC¹ and HSTCP) reduce transmission rate less than standard TCP after a packet loss
  - Unfair to TCP flows and to new flows using alternate algorithms
- Another group of algorithms (CTCP, FAST and FAST-AT) reduce transmission rate by ½ following a loss (HTCP is a hybrid with reduction between 20 and 50 %)
  - These algorithms seek to obtain higher throughput by increasing transmission rate more quickly than standard TCP (the rate of increase varies with the algorithm)
  - HTCP reverts to TCP congestion avoidance for 1 s after each loss, which can lead to lower throughputs than other alternate algorithms
- Under extreme spatiotemporal congestion, most proposed TCP replacements have a low-window threshold and revert to standard TCP congestion avoidance procedures (giving no advantage to alternate procedures)
  - FAST and FAST-AT do not use TCP congestion avoidance under any conditions, which can lead to oscillatory behavior and increased loss rates

<sup>&</sup>lt;sup>1</sup>Note that on repeated losses occurring close in time, BIC can reduce *cwnd* substantially more than standard TCP – thus, on paths with very severe congestion TCP can provide higher throughput than BIC

#### **TCP Fairness**

- All proposed TCP replacements take steps to provide improved throughput over TCP thus comparing fairness must consider relative performance of TCP flows when competing with flows using each proposed replacement
- We found CTCP, HTCP and FAST-AT to be most fair to TCP flows
  - Under low initial sst FAST-AT is more unfair because of its quick increase in rate
  - Injecting more FAST-AT packets induced more losses in TCP flows, which could recover only linearly
- We found Scalable TCP, BIC and FAST to be most unfair to TCP flows
  - Established Scalable and BIC flows (on large files) tended to maintain higher transmission rates than TCP flows after losses, while FAST recovered more quickly, and these proposed TCP replacements induced more losses in TCP flows
- ➤ HSTCP appeared moderately fair to TCP flows, especially under conditions of lower congestion and under low initial *sst* − HSTCP appeared unfair under conditions of heavy congestion
- We found that Scalable TCP, BIC and HSTCP are also unfair to competing flows that are newly arriving

#### **Utility Bounds**

- ➤ We found that proposed TCP replacements could provide increased utility (throughput) for users however, this utility would arise only under a specific combination of circumstances
  - Flow's rwnd must not be constraining flow transmission rate
  - Flow's initial sst must be relatively low
  - Flow must be transferring a large file
  - Flow's packets must be transiting a relatively uncongested path (i.e., experiencing only sporadic losses) or else users must be willing to tolerate marked unfairness in trade for increased throughput
- How likely is this combination of circumstances on a given Internet flow?
  - Certainly possible to engineer a network, or segments of a network, to provide specific users with improved throughput compared with TCP
  - We suspect a rather low probability for such circumstances to arise generally in the Internet
- We conclude that proposed TCP replacements can provide improved user throughput – however, most users seem unlikely to benefit very often

## Safety

- We can answer this only in part additional cautionary findings may be possible
  - We simulated either homogeneous networks where all flows used one congestion control algorithm or mixes of TCP flows competing with flows using one proposed TCP replacement at a time
  - The real Internet could contain a mix of many different types of congestion algorithm
- For most algorithms we studied, under most conditions, we found little significant change in macroscopic network characteristics
- FAST and FAST-AT are exceptions to this general finding
  - Under high spatiotemporal congestion, where there were insufficient buffers to support flows transiting specific routers, FAST and FAST-AT entered an oscillatory behavior where the flow *cwnd* increased and decreased rapidly with large amplitude
  - Under such conditions the network showed increased loss and retransmission rates, a higher number of flows pending in the connecting state and a lower number of flows completed over time
- We recommend the need for additional study of FAST and FAST-AT prior to widespread deployment and use on the Internet