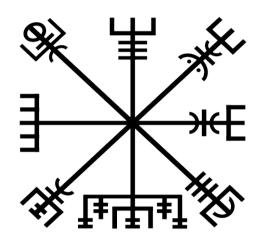
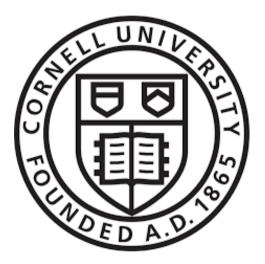
Emergency Edge Supercloud



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DISCLAIMER

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Motivation



- The 2017 Atlantic Hurricane Season was deadly and destructive
- Hundreds of lives were lost
- Over a quarter trillion dollars in estimated damages

Emergency First Responders



Loss of lives, limbs, and property would have been a lot higher if not for the efforts of thousands of first responders

Problem: communication infrastructure might not be available



- Cell towers rendered inoperable
- First responders can use only equipment they bring with them

Opportunity: every first responder is carrying a computer and network router



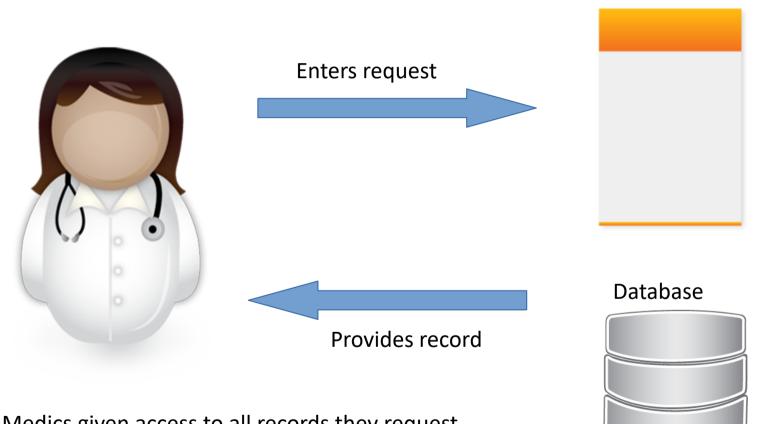
- Smartphones come with a variety of communication modalities
- Can form ad-hoc networks

Prompt and privacy aware access to medical records



Problem: loss of communication with central server

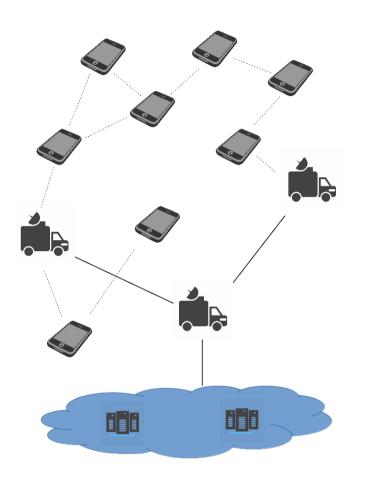
Accountability over access control



Tamperproof log

- Medics given access to all records they request
- Provided request is entered in a tamperproof log
- After emergency is over, logs are reviewed

Tamperproof log, but how?



- In an ad hoc network limited access to public cloud
- Not all nodes can be trusted

Requirements of a tamperproof log

Availability

- Must be readable at all times
- Appendable "under normal circumstances"

Scrutinizable

- Protocol and implementation must be easy to understand
- Ideally formally verified

Distributed Trust, or "trustless"

- Not under a single administrative domain
- Yet should be impersonationresistant

Performance

- Append-only

records

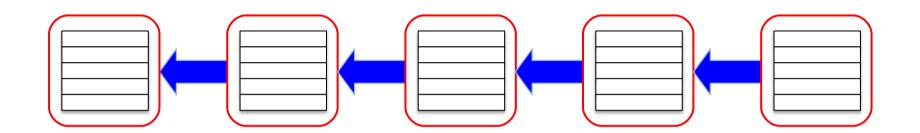
- No modification or

deletion of existing

Integrity

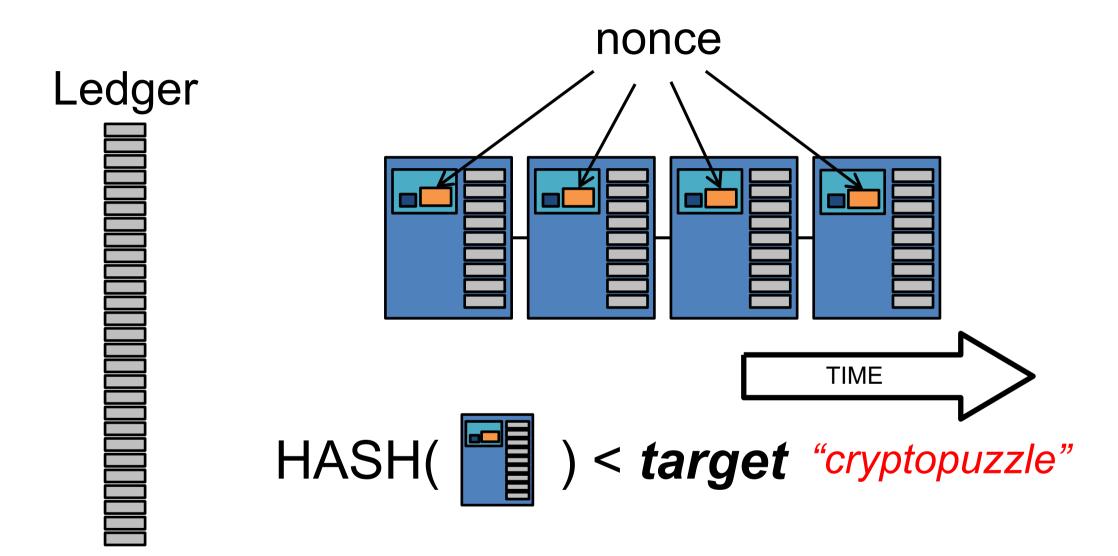
 Scale, high throughput, low latency, efficiency

Blockchain as tamperproof log

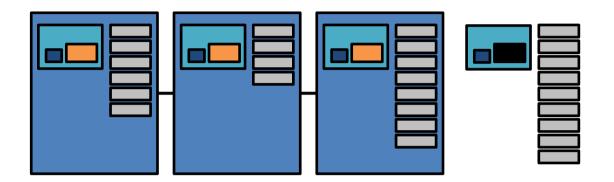




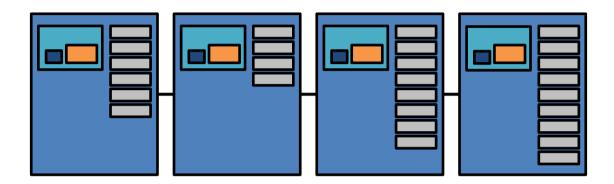
The Bitcoin Blockchain



The Bitcoin Blockchain



The Bitcoin Blockchain

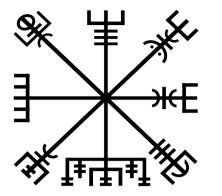


Bitcoin-style blockchains not an option

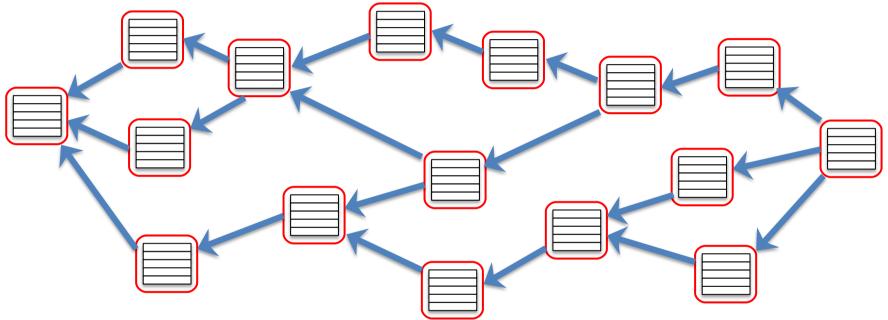
- Are computationally expensive and thus batterydraining
- Require high network connectivity
 - Miners typically want to broadcast new blocks asap
 - first miner wins the prize
 - Protocol can recover from temporary network partitions, but leads to blocks being discarded and work wasted, as well as security issues
- Lack of decentralization harms security

"Permissioned" blockchains can dispense with proof of work

- Blockchain doubles as a PKI
- Owner's self-signed certificate in genesis block
- Additional users added/removed by placing certificates/revocations on blockchain
- But system-wide consensus is not an option either



Vegvisir: tolerate branches

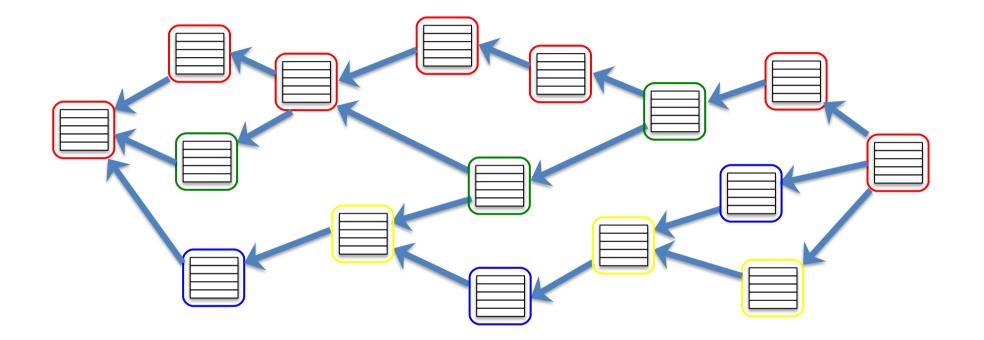


- Leads to DAG structure instead of linear blockchain
- Not good for cryptocurrencies...
- Still maintains full causal history of events

Properties

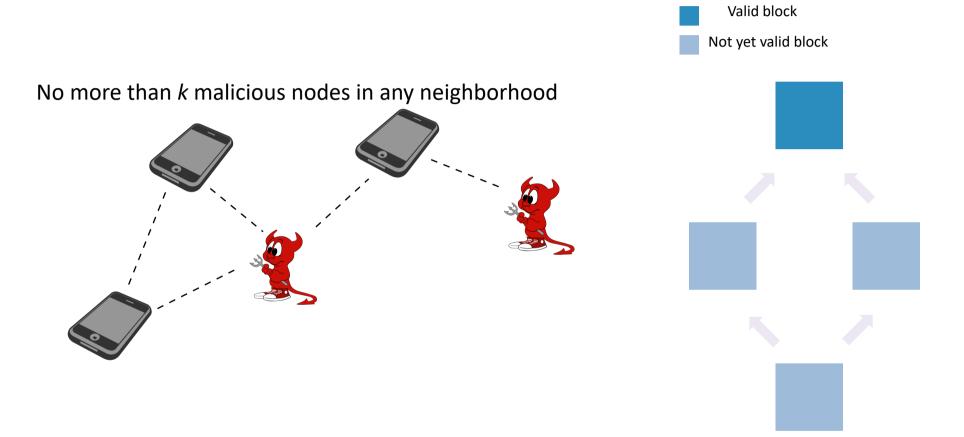
- Availability:
 - blocks, once added, cannot be removed
- Integrity:
 - unique genesis block (sink)
 - each DAG has a unique "leader block" (source)
 - each DAG is connected (and loop-free...)
 - each block signed by authorized principal
 - blocks signed by same (honest) principal on unique path
- Confidentiality
 - end-to-end encryption of transactions (optional)

Integrity



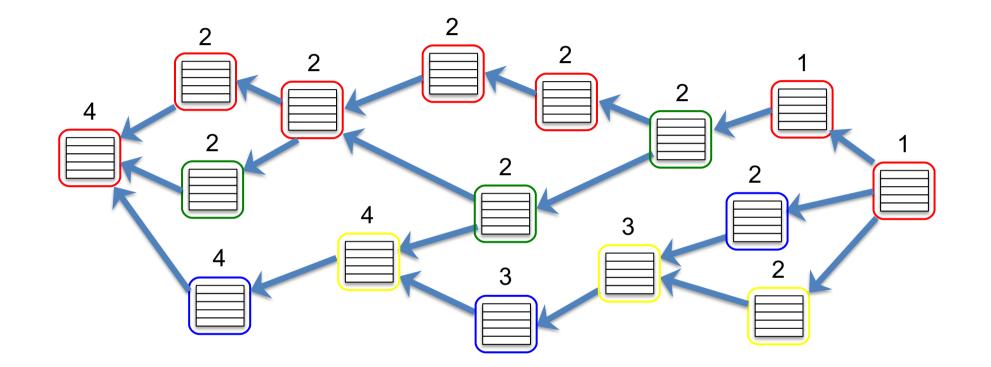
- Maximum "width" is determined by the number of principals (Byzantine principals can temporarily create higher width)
- Signed blocks on different paths are "proof-of-misbehavior"

Availability: Proof-of-Witness

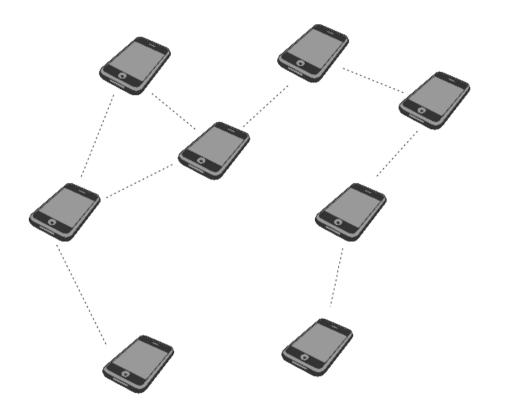


- Block will survive if >k witnesses
- If one block has PoW, so has all its ancestor blocks
- Research question: who makes good witnesses?

Counting Witnesses

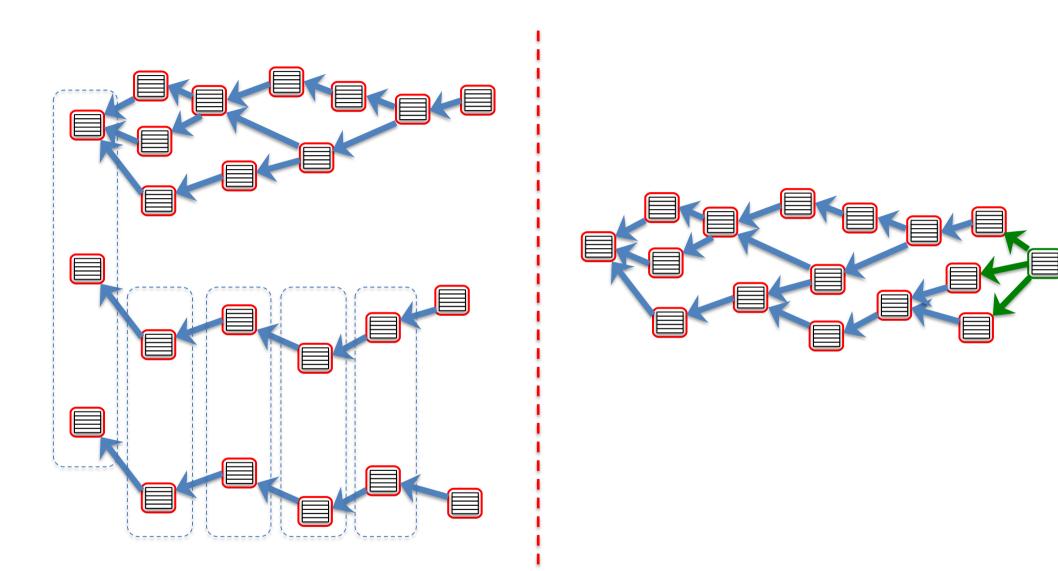


Blocks gossiped over ad hoc network

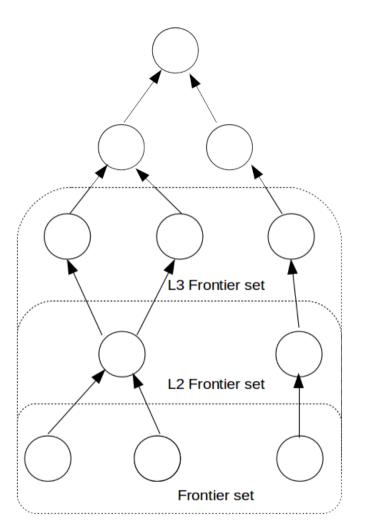


Heterogeneous, opportunistic networking

Reconciliation



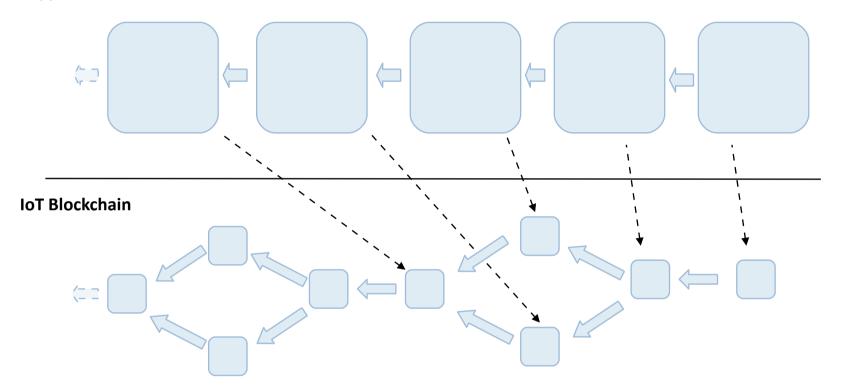
DAG Reconciliation



Peers exchange frontier sets incrementally

Offloading to "support blockchain"

Support Blockchain



- Allows regular peers to discard old blocks when storage space is low
- Design invariant: availability of a block is monotonically increasing

Interpreting Vegvisir

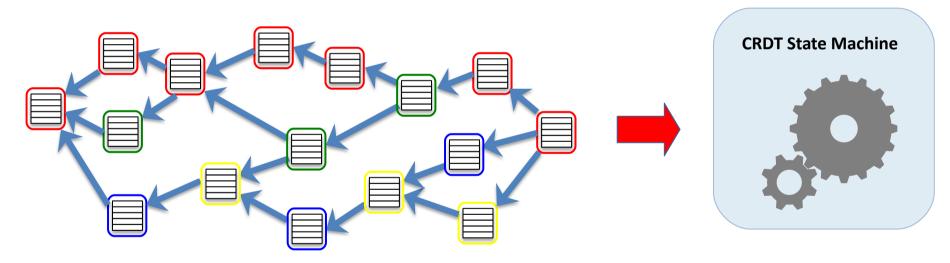
- Vegvisir provides a shared, tamperproof data repository that keeps track of data provenance and distributes trust over peers
- Only requirement is that updates on shared data structure in some sense commute

Conflict-Free Replicated Datatypes (CRDTs)

- Updates must be associative, commutative, idempotent
- Nodes can be updated independently
- Basic CRDTs form registers, counters, sets
- Can be combined and composed

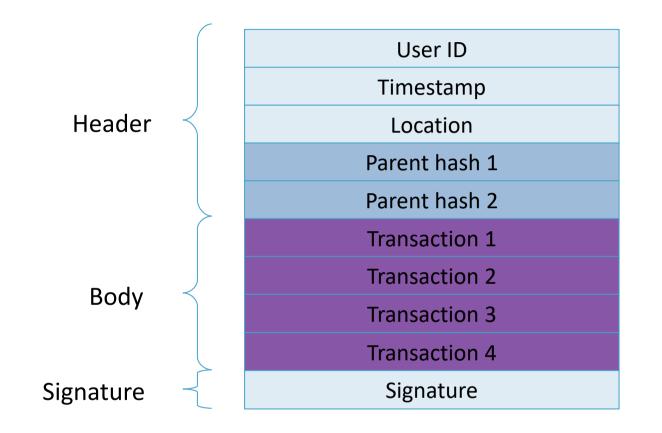
Vegvisir has two main components

- Blockchain itself
- CRDT state machine



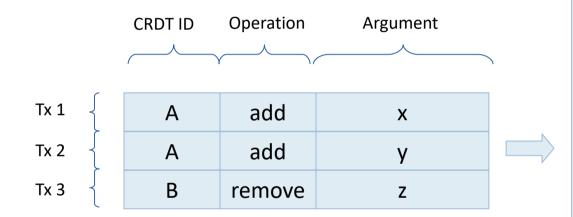
Operations only applied if PoW available

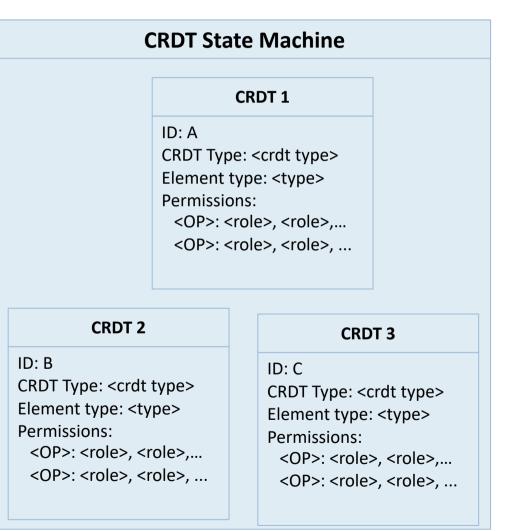
Vegvisir Block Structure

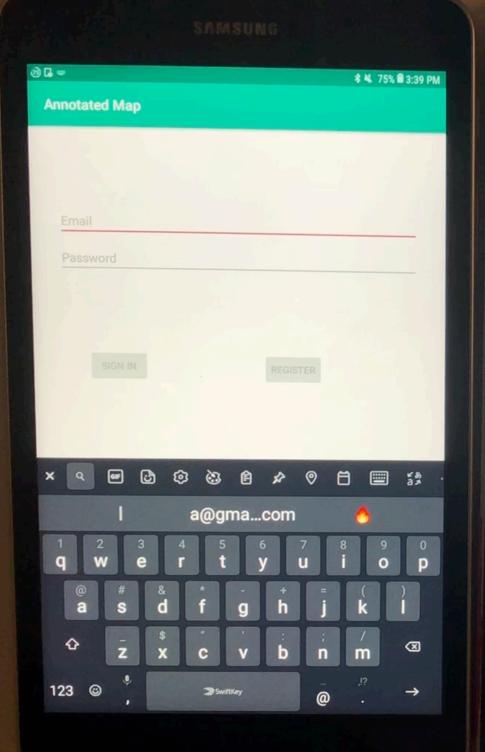


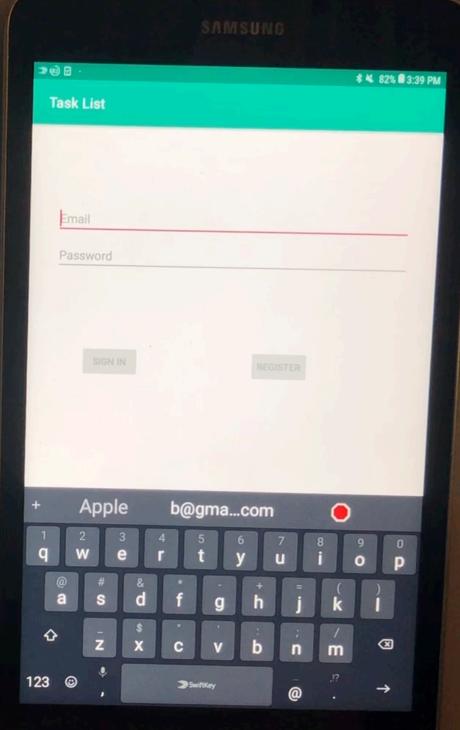
Blocks are certificates

Transactions manipulate CRDTs







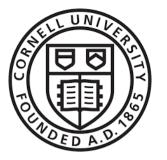


Membership

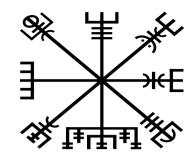
- One special CRDT (2P) maintains membership
 - add-membership
 - revoke-membership
- Proof-of-misbehavior also implicitly revokes membership
- Only members can add new blocks

ARM TrustZone

- ARM TrustZone "secure worlds" can help:
 - Who is a good witness?
 - secure access to device location and time
 - Check PoW and provide access to secured data
 - Secure sensor values
 - secure retrieval of sensor values







- Vegvisir is a DAG-based blockchain to allow for partitioned operations
 - not for higher throughput per se
- Replaces Proof-of-Work with *"Proof-of-Witness"*
- CRDTs enable consistently evolving views
- Prototype available for Android devices

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