



NFORMATION ECHNOLOGY ABORATORY

## Routing Data Quality and Its Impact on BGP Anomaly Detection Algorithms

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http://www.antd.nist.gov/bgp\_security/

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## **Talk Based on Previous Publications**

- DHS CATCH Conference, Washington DC, March 2009 <u>http://www.nist.gov/itl/antd/upload/NIST\_BGP\_Robustness-</u> <u>2.pdf</u>
- NANOG-45, Santo Domingo, January 2009

http://www.nanog.org/meetings/nanog45/abstracts.php?pt= MTE5NSZuYW5vZzQ1&nm=nanog45

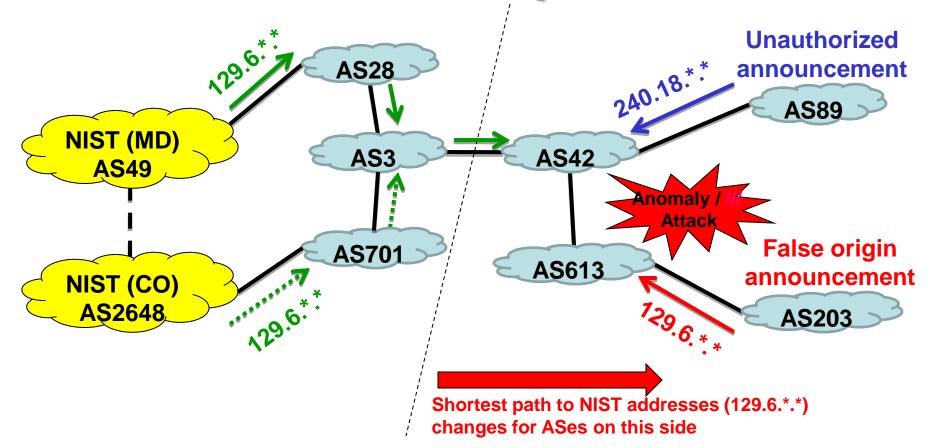
• ARIN-23, San Antonio, TX, April 2009

https://www.arin.net/participate/meetings/reports/ARIN\_XXI II/pdf/monday/nethandles.pdf

# Theme of the Talk

- Registered data (RIR, IRR, RADB, RPKI, etc.) as well as historical BGP trace data are and/or will likely be the basis for implementing routing robustness and security
- Characterization of correctness and completeness of the data
- Data pruning to improve its reliability
- What implications does the data quality have on BGP robustness/security algorithms?
  - Focus: Reduce probability of false alarms & false negatives

# One Aspect of BGP Robustness Problem Space



Other aspects: Route leaks, Path modification

# **Data Driven BGP Robustness**

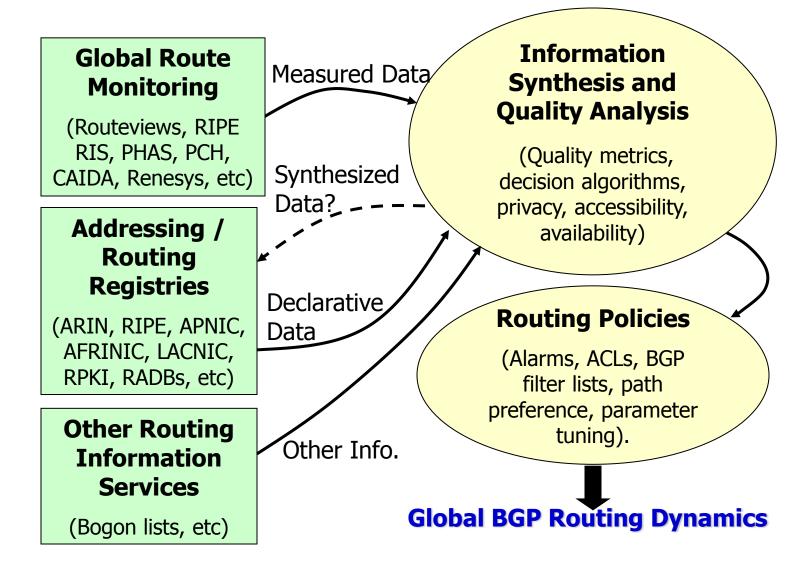
#### What are the Data Sources?

- Addressing Registries
  - global databases of address block and autonomous system number assignments.
- Routing Registries
  - loosely maintained global databases of contractual relationships for routing services.
- Monitoring Data
  - public BGP monitoring and measurement projects that collect BGP protocol exchanges at various spots around the Internet.

#### Why is this hard?

- Registries
  - known to be incomplete and inaccurate, and are maintained in differing formats, by differing processes in different regions of the world.
- Robustness Algorithms
  - to be effective, must make precise policy decisions from imperfect data.
- Needle in a Hay Stack
  - millions of BGP update messages per day; millions of registry entries; rare but potent threats.

# **Solution Components / Players**



# **Registry Data and Analysis of Its Completeness and Correctness**

## **Registry Data Object Counts by Source**

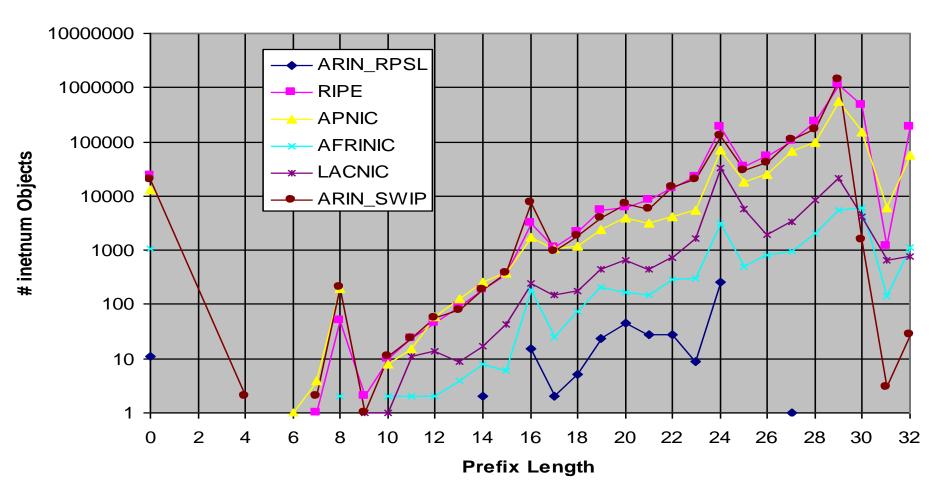
		route		inetnum (ARIN NetHandle)			aut-num (ARIN ASHandle)			
RIR/IRR	06/18/2007	10/18/2008	Incr	06/18/2007	10/18/2008	Incr	06/18/2007	10/18/2008	Incr	
ARIN	7,330	8,201	12%	338 (1,618,197)	434 (1,924,454)	28% 19%	758 (18,050)	890 (19,678)	17% 9%	
RIPENCC	71,569	89,957	26%	2,044,536	2,458,119	20%	14,106	16,969	20%	
APNIC*	23,616	35,515	50%	822,891	1,080,999	31%	4,559	5,347	17%	
AFRINIC	0	0		13,948	22,706	63%	342	445	30%	
LACNIC**	0	0		45,346	83,036	83%	1,219	1,339	10%	
Standalone IRRs+	345,129	497,124	44%	1	1		3,785	4,643	23%	
Total:	447,644	630,797	41%	2,927,060 (1,618,197)	3,645,295 (1,924,454)	25% 19%	24,769 (18,050)	29,633 (19,678)	20% 9%	

\* Includes TWNIC, JPIRR, JPNIC and APNIC

\*\* RIR only

+ Independent IRR databases that are mirrored via the RADB website including RADB, but EXCLUDING ARIN, APNIC, JPIRR and RIPE Note that route objects can be registered at any IRR regardless of where the address spaces are allocated.

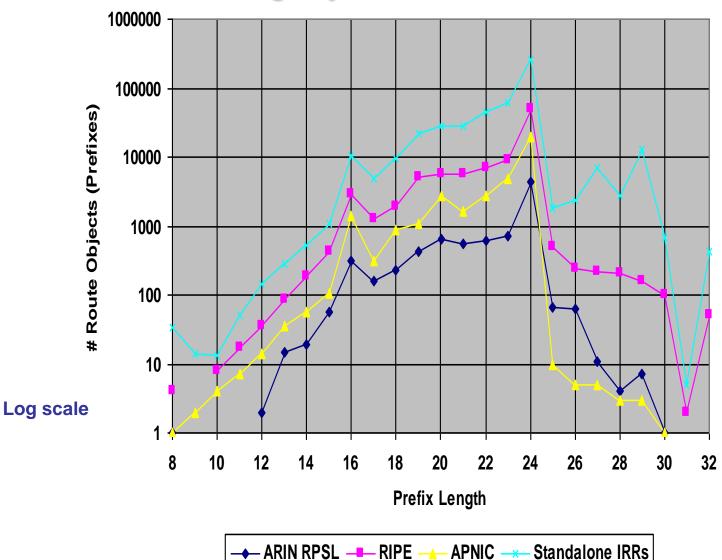
# Distribution of Prefix Length of inetnum (RPSL) and NetHandle (SWIP)



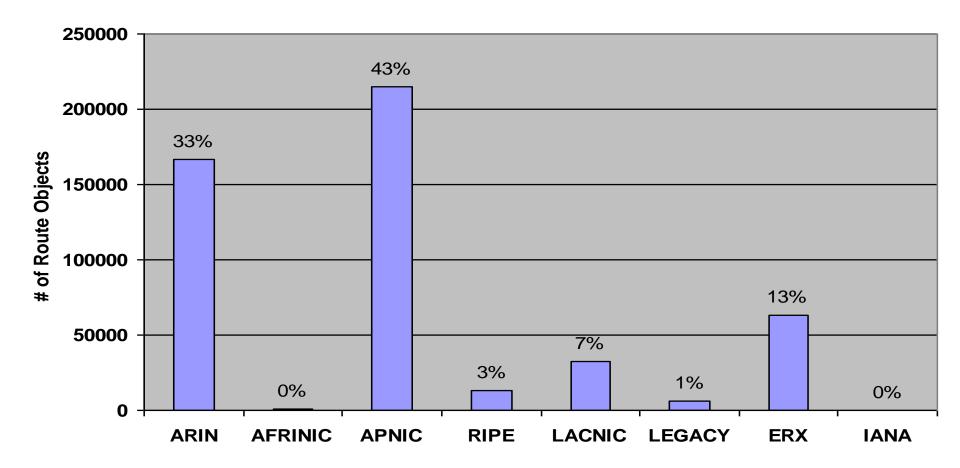
- Length 0 indicates that an address block cannot be represented by a single CIDR
- Length 4 specifies Multicast and Reserved Future Use blocks
- Some Legacy and ERX blocks may be included in one or more RIRs

#### Distribution of Prefix Length of Route Objects in IRR

Registry Data Date: 2008-10-18

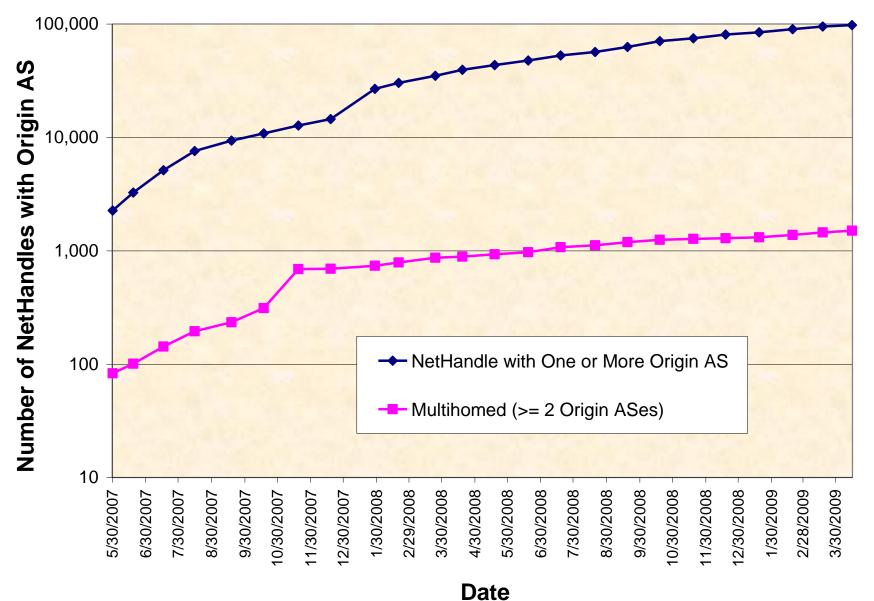


#### Distribution of Sources of Prefix Allocations of Route Objects Registered to Standalone IRRs

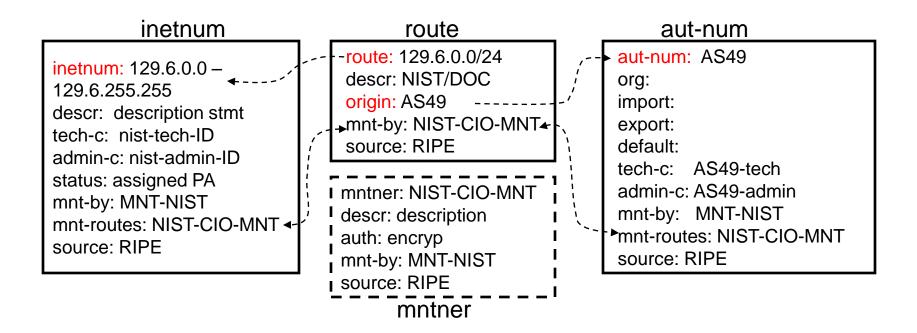


All route objects registered in standalone IRRs on 2008-10-18: 497,124

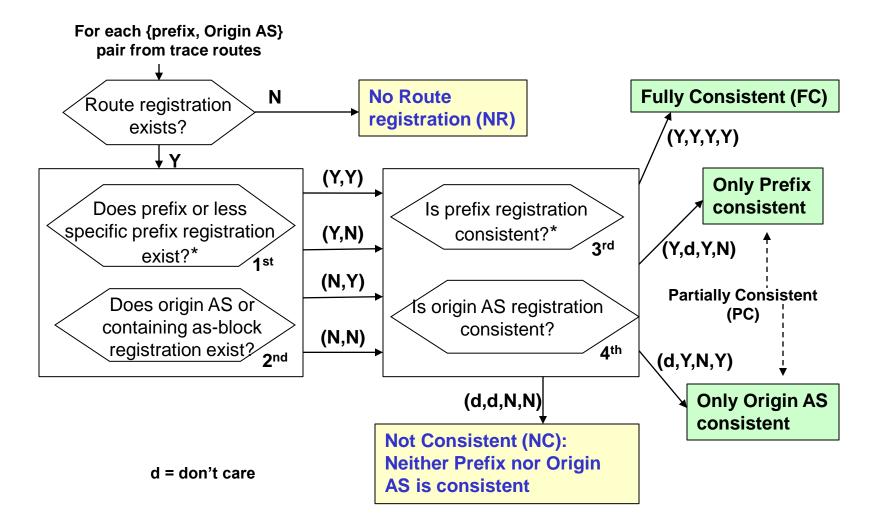
## **Growth of NetHandles with OriginAS**



#### Checking Consistency of a Registered Route with Corresponding Inetnum and Aut-Num

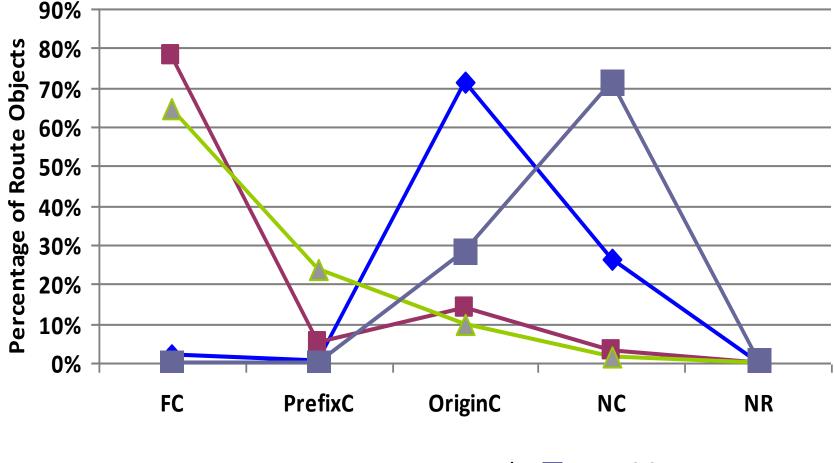


#### Registry-Based Algorithm for Scoring Routes Observed in Trace Data



#### Characterization of IRR Consistency Based on Route Object Registrations

Registry Data Date: 2008-10-18



ARIN RPSL — RIPE — APNIC\* — Standalone IRRs

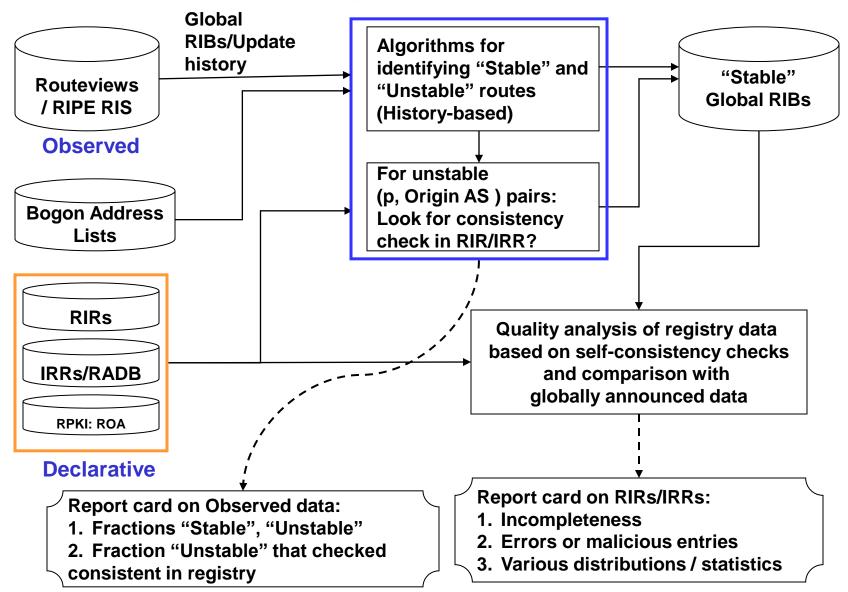
 Note: This does not reflect the total number of routes registered at each IRR. For example ARIN has only 8K whereas RIPE has 90K as of 2008-10-18.

## **BGP Robustness Algorithms**

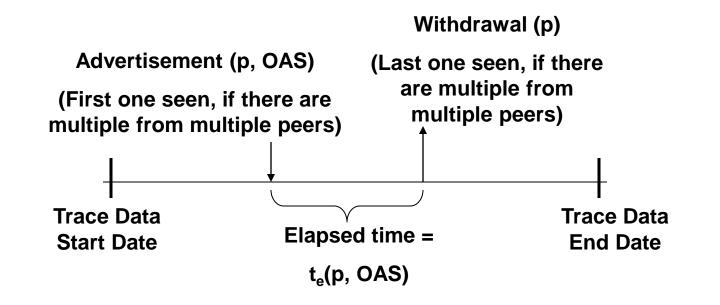
## **Known BGP Robustness Algorithms**

- General goal: Validate an observed (p, Origin AS) pair
- Nemecis: Compare with registered objects (route, inetnum, autnum)
- PHAS: Compare with historically observed (p, Origin AS) pairs, AS-paths:
  - Identify origin changes, subprefix announcements; generate alerts
- Pretty Good BGP (PGBGP): Compare with historically observed (p, Origin AS) pairs
  - Influence forwarding or holding back of updates in real-time in BGP processing

### **New Integrated Approach**

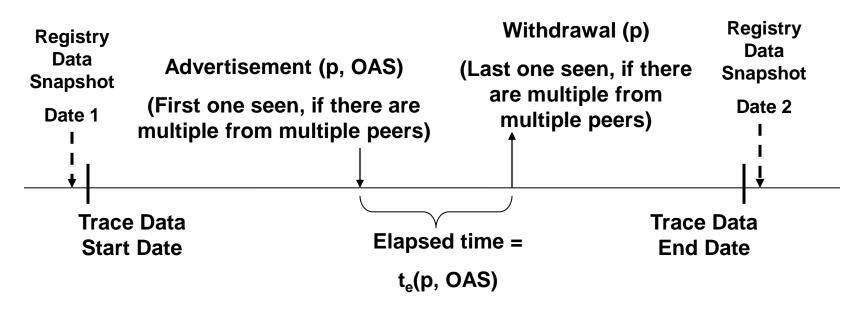


#### Enhanced History-Based Algorithm for Determining Stability of (p, OAS) in the Trace Data



- If  $t_e(p, OAS) \ge 48$  hours, then (p, OAS) is a stable (prefix, Origin AS) pair
- If  $t_e(p, OAS) < 48$  hours, then (p, OAS) is an unstable (prefix, Origin AS) pair
- Update data is initialized with stable (i.e., persistent for 
  <u>></u> 48 hours) RIB entries

# Enhanced Hybrid Algorithm for Validating (p, OAS) in the Trace Data



- Use enhanced history-based (i.e., trace-data-based) algorithm as in previous slide
- Complement it with combined results of the registry-based algorithm with data from two dates (close to start and end dates of the history algorithm)
- Result: Better performance of anomaly detection algorithms

## Comparative Analysis of Existing and Enhanced Algorithms

- We have encoded Registry-based, Enhanced Trace-data-based and Enhanced Hybrid algorithms for evaluation
- Algorithms are run on top of the NIST TERRAIN\* framework
  - Unified database of Registry / Trace data (RIRs, IRRs, RIPE-RIS, Routeviews)
- Tested and compared the algorithms

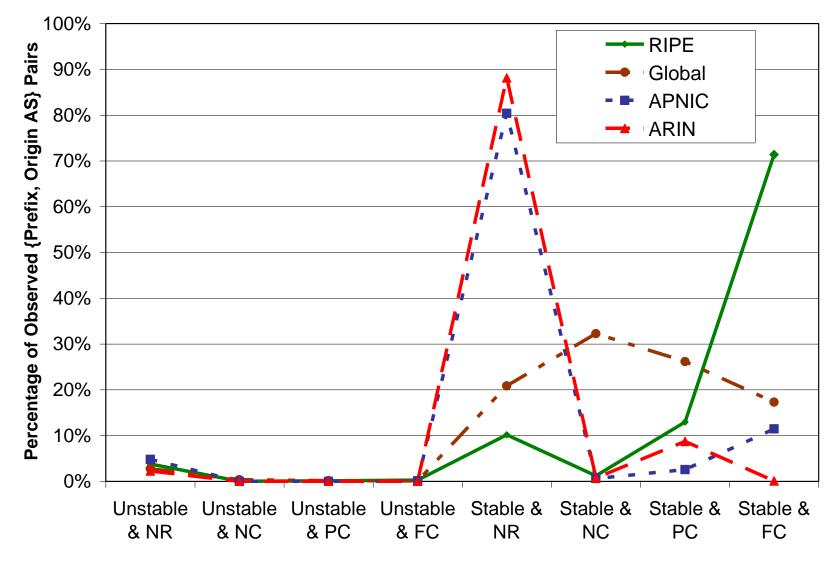
\* TERRAIN: Testing and Evaluation of Routing Robustness in Assurable Inter-domain Networking

## Comparative Analysis of Existing and Enhanced Algorithms (Contd.)

For the purpose of this presentation:

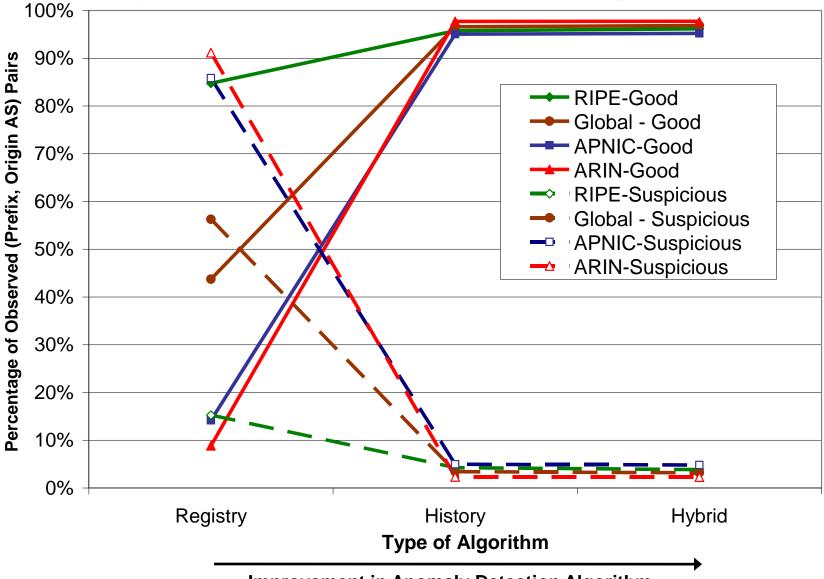
- Results focus on Origin AS validation
- Results are reported globally for all prefixes as well as selectively for regional (RIPE, ARIN, ...) prefixes
- Six-month trace-data window (January through June 2007); initialized with stable RIB entries
- Registry data two dates prior to and towards the end of the six-month window (December 12, 2006 and June 18, 2007)

#### Classification of Observed (p, OAS) Pairs According to Stability / Consistency Scores



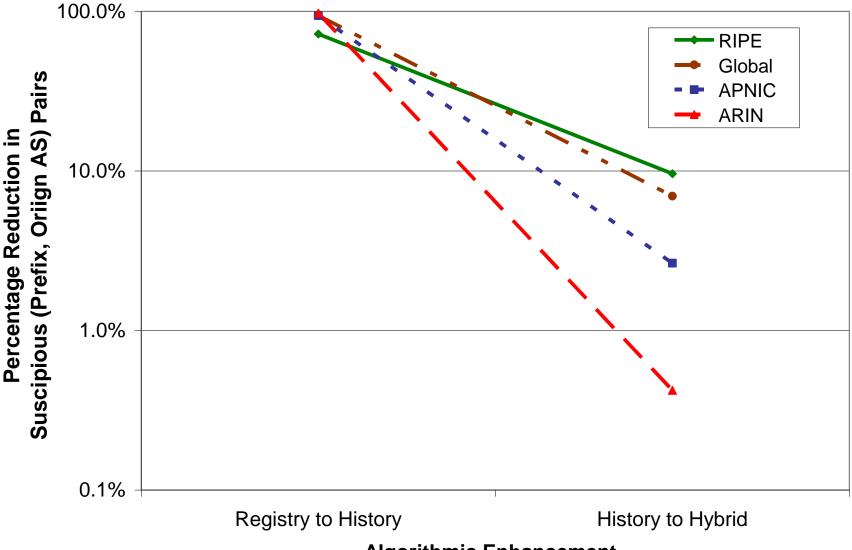
p = prefix; OAS = Origin AS; FC = Fully Consistent; PC = Partially Consistent; NC = Not Consistent; NR = Not Registered

#### **Comparative Performance of Algorithms**

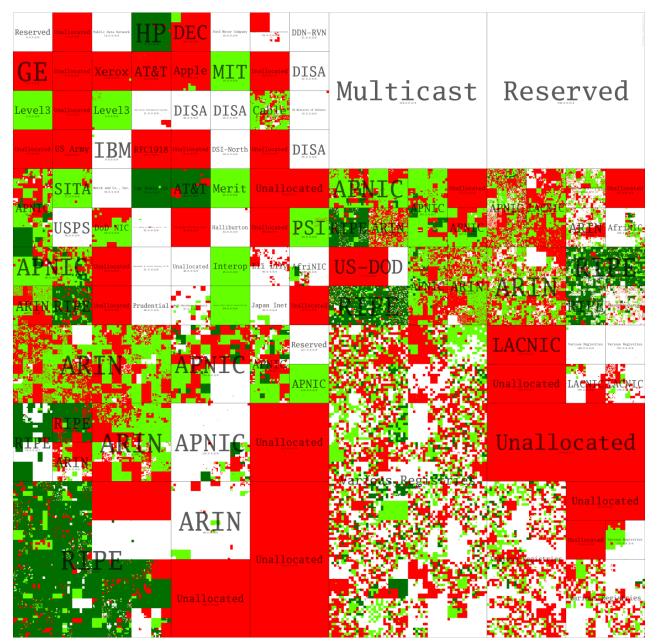


Improvement in Anomaly Detection Algorithm

#### **Comparative Performance of Algorithms**



#### **Checking Origin AS : Comparison of Algorithms**



#### Registry-based Algorithm

Green: Good / FC Light Green: Good / PC Red: Suspicious White: Not found in trace data

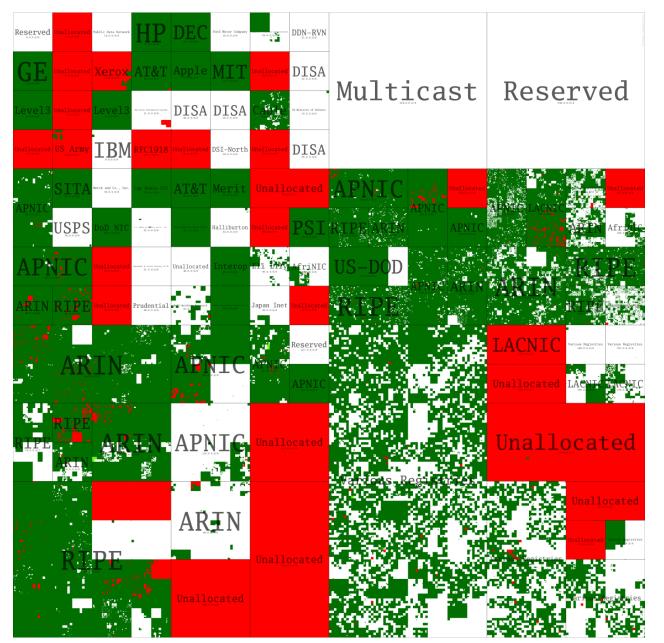
#### **Checking Origin AS : Comparison of Algorithms**



#### Enhanced tracedata-based Algorithm

Green: Good Red: Suspicious White: Not found in trace data

#### **Checking Origin AS : Comparison of Algorithms**



#### Enhanced Hybrid Algorithm

Green: Good / FC Light Green: Good / PC Red: Suspicious White: Not found in trace data

## Summary

- Examined and quantified the quality (completeness, correctness) of registry data
- Enhanced hybrid algorithm history and registry data have complementary influence on improvement in origin validation
- Further testing for robustness of the algorithms needs to be performed with extensive real and synthetic trace data
- NIST has begun to monitor and quantify the growth and quality of the RPKI data

# **Backup slides**

## **Prefixes with Multiple Origin ASes**

# Origin ASes	# Prefixes		
1	476243		
2	55673		
3	10419		
4	2683		
5	965		

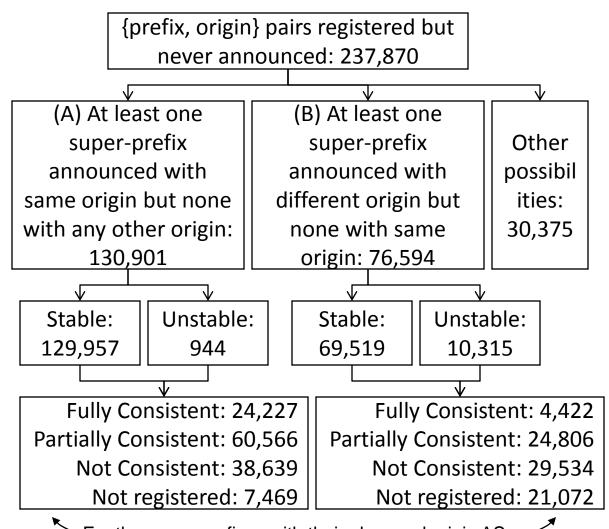
#### For prefixes with two Origin ASes:

OAS1	OAS2	# Prefixes
FC + Stable	FC/PC + Unstable	23
PC + Stable	FC/PC + Unstable	41
NC + Stable	FC/PC + Unstable	104
NR + Stable	FC/PC + Unstable	0
Total		168

 Statistics of prefixes with two Origin ASes where the primary path is stable (with or without consistency in the registry), while the secondary (failover) path is <u>transient (unstable) but consistent</u> in the registry

#### **Analysis of Registered But Unobserved Routes**

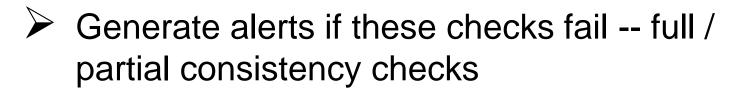
- Large number of {prefix, origin} pairs registered but never announced
- In most cases, superprefixes are announced with the same origin AS (as in registered route) or a different origin AS
- Is it due to aggregation by a higher tier ISP?



 $^{\sim}$  For the super-prefixes with their observed origin ASes  $^{\prime\prime}$ 

## **Nemecis: Registry Based Algorithm**

- For (p, Origin AS) pair from an update:
  - Check for existence of prefix, autnum, and route objects in RIR/IRR
  - Check for consistency between these declared objects by matching Organization, maintainer, email, etc.



G. Siganos and M. Faloutsos, "A Blueprint for Improving the Robustness of Internet Routing," 2005. <u>http://www.cs.ucr.edu/%7Esiganos/papers/security06.pdf</u>

G. Siganos and M. Faloutsos, "Analyzing BGP policies: methodology and tool," IEEE Infocom, 2004. <u>http://www.cs.ucr.edu/~siganos/papers/Nemecis.pdf</u>

## **PHAS: Prefix Hijack Alert System**

- Make use of BGP trace data
- Provide alert messages if:
  - Origin AS set changes
  - New subprefix is added to observed set of subprefixes

#### Last-hop AS set changes

Mohit Lad, Dan Massey, Yiguo Wu, Beichuan Zhang and Lixia Zhang, *PHAS: A prefix hijack alert system*, North American Network Operators Group Meeting (NANOG-38), October, 2006. <u>http://www.nanog.org/mtg-0610/presenter-pdfs/massey.pdf</u>

Mohit Lad, Dan Massey, Dan Pei, Yiguo Wu, Beichuan Zhang and Lixia Zhang, *PHAS: A prefix hijack alert system*, in Proceedings of 15th USENIX Security Symposium (USENIX Security 2006). <u>http://www.cs.ucla.edu/~mohit/cameraReady/ladSecurity06.pdf</u>

#### PGBGP: Pretty Good BGP Old Version of the Algorithm

- Observed {prefix, Origin AS} pairs based on update history and RIB entries over the last h days (h = 10 days) are recorded
- An update for a prefix is considered suspicious if the origin AS is new relative to the history record; the update is propagated with lower local pref
- A new subprefix (of a prefix in history record) is always considered suspicious and quarantined
- The quarantine lasts for suspicious period of s hours (s = 24 hours); if the subprefix is not withdrawn during that time, then the update is propagated

#### **One Weakness of Old PGBGP** From NANOG discussions back in 2006

- Q: Panix's first, obvious countermeasure aimed at restoring their connectivity announcing subprefixes of their own address space would <u>also</u> have been considered suspicious, since it gave two "sub-prefixes" of what ConEd was hijacking?
- A: [Here] things get a little more subtle. We have considered allowing the trusted originator of a prefix to split the space among itself and those downstream of it without considering that suspicious behavior.

Note: This was part of the Q&A after the paper on PGBGP was presented by J. Karlin at NANOG-37. <u>http://www.nanog.org/mtg-0606/pdf/josh-karlin.pdf</u>

# **New Version of PGBGP**

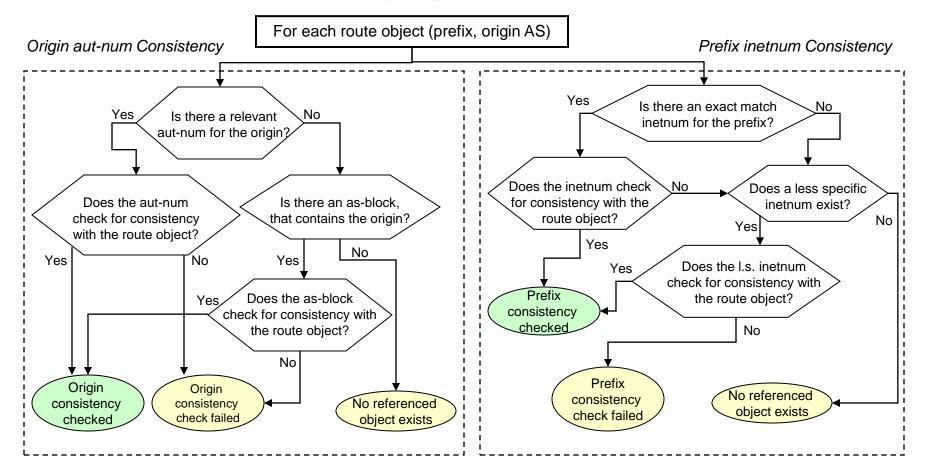
- From an updated new version of PGBGP paper:
  - "PGBGP would not interfere if an AS announces sub-prefixes of its own prefixes in order to gain traffic back during a prefix hijack."

Josh Karlin, Stephanie Forrest, and Jennifer Rexford, "Pretty Good BGP: Improving BGP by Cautiously Adopting Routes," The 14th IEEE International Conference on Network Protocols, November 2006. <u>http://www.cs.unm.edu/~treport/tr/06-06/pgbgp3.pdf</u>

## **Potential Weaknesses of (New) PGBGP**

- The <u>short-span historical view</u> (last ten days) has the following negative implications:
  - PGBGP will typically unnecessarily lower local-pref on path announcements due to multi-homing related AS origin change.
  - If a malicious user observes a prefix withdrawal by genuine origin AS and announces the prefix at that time, the malicious path propagates with a lower localpref value and will be used (Effectively - False Negative).
  - If the prefix owner sometimes announces sub-prefixes in conjunction with multi-homing related AS origin change, PGBGP will quarantine the announcements.

#### Checking Registry Consistency of Registered Routes (Algorithm)



#### **Origin AS Approval Check List: Comparison**

		Which checks are included in each approach?				
	Checks/Questions	Registry based (e.g., Nemecis)	Trace- data based (PGBGP)	Enhanced Trace- data based	Enhanced Hybrid	
Q1.	Is prefix registered (same or less specific)?					
Q2.	Is there a route registered (with same or less specific prefix and origin AS)?	$\checkmark$				
Q3.	Is announced (p, origin AS) fully consistent with corresponding registry objects in RIR/IRR?	$\checkmark$				
Q4.	Is announced (p, origin AS) partially consistent with corresponding registry objects in RIR/IRR?	$\checkmark$			$\checkmark$	
Q5.	Was (p, origin AS) seen in RIB in the last $h (= 10)$ days? (Also, if it was suspicious, did it remain in RIB beyond the suspicious period of $s (= 24)$ hours?)		$\checkmark$			
Q6.	Would a less specific prefix with the same origin AS pass the test in Q5?		$\checkmark$			
Q7.	Was prefix previously announced by the same origin AS and remained stably (48 hrs or more) in the RIB over the observation period ( <i>d</i> months)?			$\checkmark$		
Q8.	Would a less specific prefix with the same origin AS pass the test in Q7?					

#### **Algorithm Robustness Checklist**

	Algorithmic Features	Registry based (e.g., Nemecis)	Trace-data based (PGBGP)	Enhanced Trace-data based	Enhanced Hybrid
Sets	Utilization of self-consistent registry objects	Yes	No	No	Yes
a Se	Utilization of update history	No	Yes	Yes	Yes
Data	Utilization of historical RIB entries	No	Yes	Yes	Yes
	Pass a subprefix announcement if a less specific prefix with same origin AS could be passed	Yes	Yes	Yes	Yes
Handled	False Positives: Alert raised when genuine prefix owner announces multi-homing related AS origin change	Moderate probability	High probability	Moderate probability	Low probability
Situations	Alert raised when attacker announces a prefix after sensing it has just been withdrawn	Yes	NO (Path propagates with lower pref)	Yes	Yes
Si	Pass a subprefix announcement in conjunction with multi-homing related AS origin change	Moderate probability	Low probability	Moderate probability	High probability

\* This is a ballpark qualitative assessment; subject to corroboration using extensive quantitative studies.

## Some Caveats Apply

- This presentation is mainly to demonstrate the capability and to solicit feedback on approach
- Quantitative results are subject to change when the following enhancements to the study are made (ongoing / future work)
  - Consideration of new NetHandle format in ARIN which includes origin AS information
  - Consideration of multiple trace-data collectors (here we considered trace-data from RRC00 only)
  - Use of ROAs based on RPKI efforts (in future)

#### Heatmap Depicting Origin Validation for Announced Prefixes

Reserved Unallocated Party	HP	DEC	institutor Communi 1933 - Vit		DDN-RVN					a	
GE Inallocated Xe	erox AT&T	Apple	MIT	Unallocated	DISA	Multicast					
Level3 thallocated Ley	vel3	DISA	DISA	Cable	Statistics of primer (0.4 + 1+						
Unallocated US Army I	BM RFC1918	final located	DSI-North	Unallacated	DISA						
APNIC SITA-	end in , line Case Baltice OCS	AT&T	Merit	Unalļ	ocated	APNIC APNIC		Unallocated	APNIC LACNIC	Group Segment Thallocated	
	D.NIC		Halliburton	Unallocated	PSI	RIPE ARI		APNIC	AFNIC LICHIC	ARIN Afrinic	
APNIC	llocated	thallocated	Interop	Eli Lily	AfriNIC	US-DO	APNIC			RIPE	
ARIN RIPE	llocated Predential	hili darapa insara E 1 + 10		Japan Inet	Unallocated	RIPE				RIPE THE ADDRESS	
ART	N	APNIC		APNIC Reserved					LACNIC	far laan Bayleardan 1800 - Arigan Bayleardan 1800 - Arigan	
AKL	IN								Unallocated LACNIC LACNI		
$\frac{\text{RIPE}}{\text{ARIN}} \text{ARIN}$		AP <u>N</u> IC		Unallocated		Various Registries			Unallocated		
RIPE		ARIN Unallocated		Unallocated		various registites			Unallocated		
									Various Registries	Unallocated Tering Auginotics	
								and a second sec	Various Registries		



- a. Allocations
- b. Registry-based Algorithm
- c. Enhanced Tracedata-based Algorithm
- d. Enhanced Hybrid Algorithm

For (b), (c), (d) : Green: Good / FC Light Green: Good / PC Red: Suspicious White: Not found in trace data

Reference: http://maps.measurementfactory.com/software/ipv4heatmap.1.html

