Area Laws for Entanglement

Fernando G.S.L. Brandão

University College London

joint work with

Michal Horodecki

arXiv:1206.2947

arXiv:1406.XXXX

Stanford University, April 2014

Goal: Lay down the theory for future quantum-based technology (quantum computers, quantum cryptography, ...)

Quant. Comm.
Entanglement theory
Q. error correc. + FT
Quantum comp.
Quantum complex. theo.

Goal: Lay down the theory for future quantum-based technology (quantum computers, quantum cryptography, ...)

Ultimate limits to information transmission

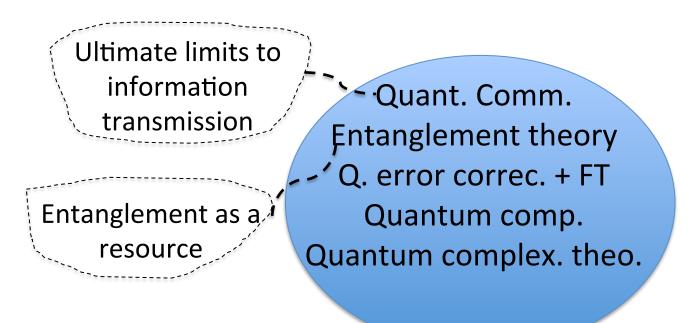
• Quant. Comm.

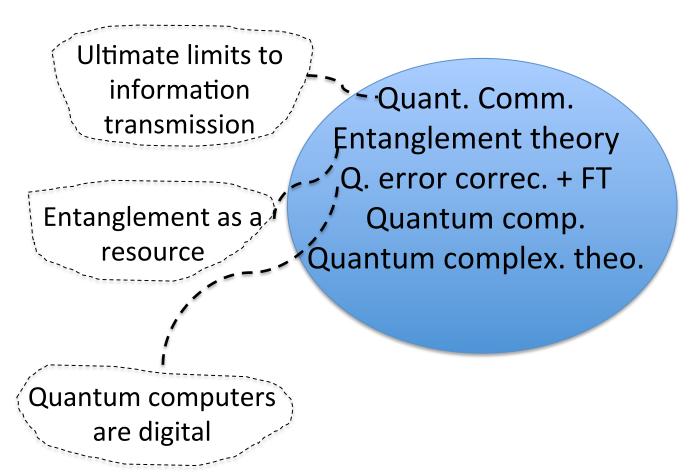
Entanglement theory

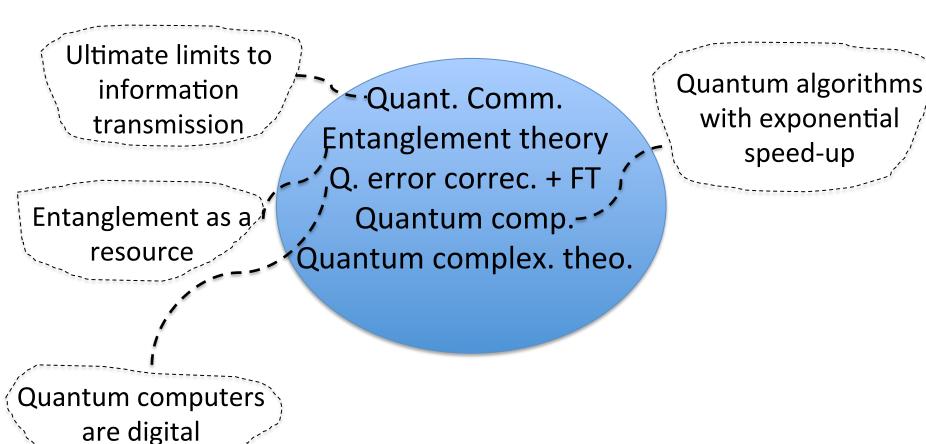
Q. error correc. + FT

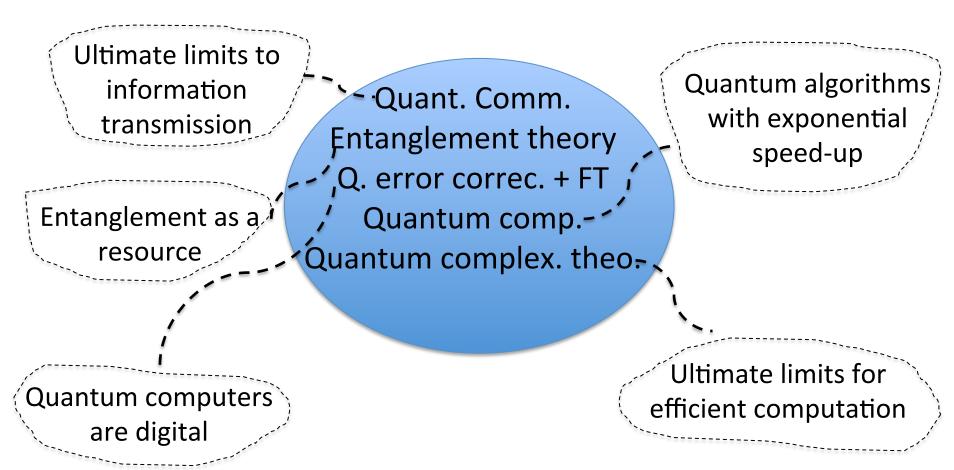
Quantum comp.

Quantum complex. theo.









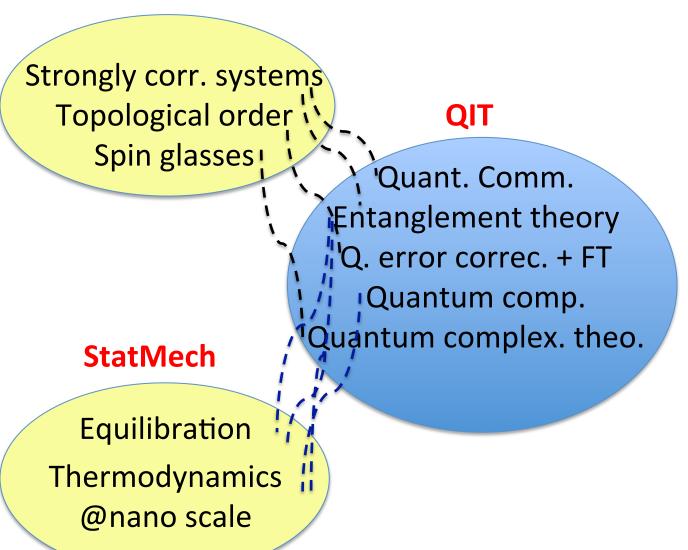
QIT

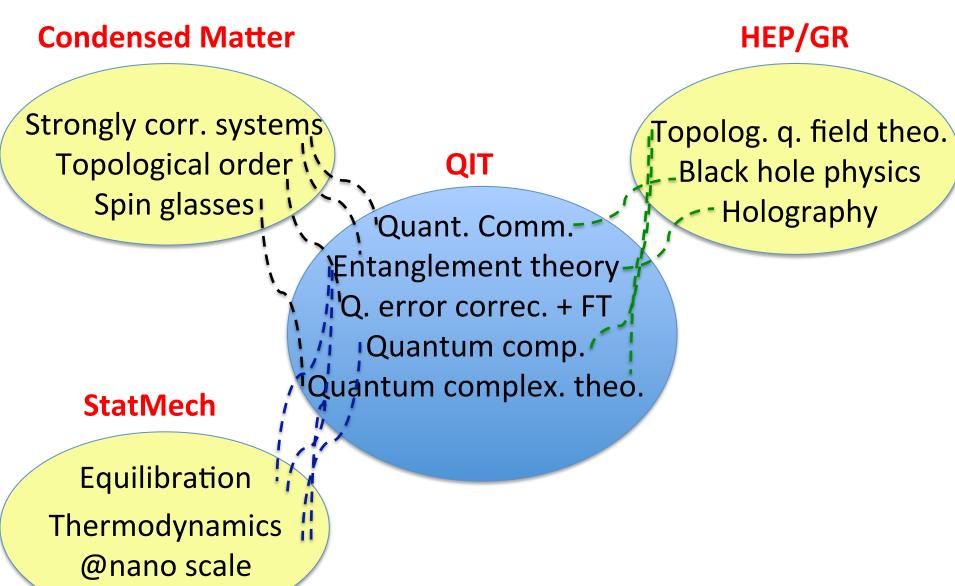
Quant. Comm.
Entanglement theory
Q. error correc. + FT
Quantum comp.
Quantum complex. theo.

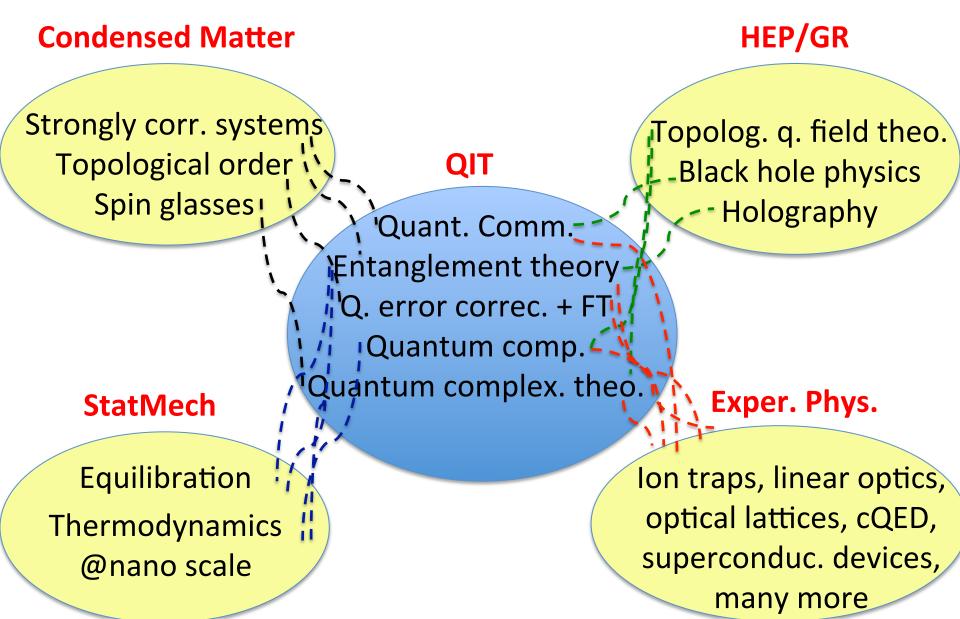
Condensed Matter

```
Strongly corr. systems
Topological order
Spin glasses
QIT
Spin glasses
Quant. Comm.
Entanglement theory
Q. error correc. + FT
Quantum comp.
Quantum complex. theo.
```

Condensed Matter







This Talk

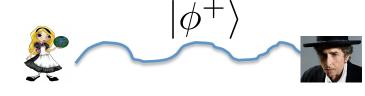
Goal: give an example of these emerging connections:

Connect behavior of correlation functions to entanglement

Entanglement

Entanglement in quantum information science is a **resource** (teleportation, quantum key distribution, metrology, ...)

Ex. EPR pair
$$|\phi^+\rangle=(|0,0\rangle+|1,1\rangle)/\sqrt{2}$$



How to quantify it?

Bipartite Pure State Entanglement

Given $|\psi\rangle_{AB}$, its entropy of entanglement is

$$E(|\psi\rangle_{AB}) := S(\rho_A) = S(\rho_B)$$

Reduced State: $\rho_A := \operatorname{tr}_B(|\psi\rangle\langle\psi|_{AB})$

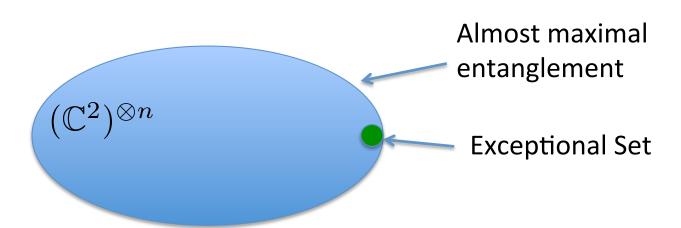
Entropy: $S(\rho) = -\operatorname{tr}(\rho \log \rho)$ (Renyi Entropies: $S_{\alpha}(\rho) := \frac{1}{1-\alpha} \log \operatorname{tr}(\rho^{\alpha})$)

Entanglement in Many-Body Systems

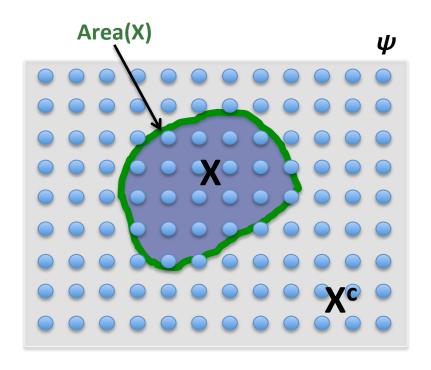
A quantum state ψ of n qubits is a vector in $(\mathbb{C}^2)^{\otimes n} \cong \mathbb{C}^{2^n}$

$$|\psi\rangle = \sum_{i_1,\dots,i_n} c_{i_1,\dots,i_n} |i_1,\dots,i_n\rangle$$

For almost every state ψ , $S(X)_{\psi} \approx |X|$ (for any X with |X| < n/2) |X| := # qubits in X



Area Law

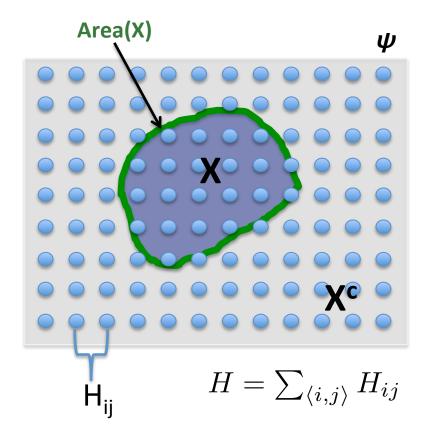


Def: ψ satisfies an area law if there is c > 0 s.t. for every region X,

 $S(X) \le c Area(X)$

Entanglement is Holographic

Area Law



Def: ψ satisfies an area law if there is c > 0 s.t. for every region X,

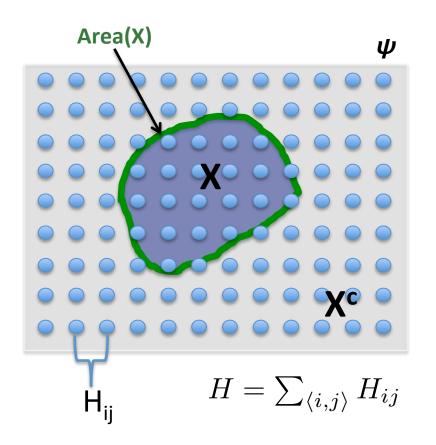
 $S(X) \le c Area(X)$

Entanglement is Holographic

When do we expect an area law?

Low-energy states of many-body local models: $H|\psi_0\rangle=E_0|\psi_0\rangle$

Area Law



Def: ψ satisfies an area law if there is c > 0 s.t. for every region X,

 $S(X) \le c Area(X)$

Entanglement is Holographic

When do we expect an area law?

Low-energy states of many-body local models: $H|\psi_0\rangle=E_0|\psi_0\rangle$

(Bombeli *et al* '86) massless free scalar field (connection to Bekenstein-Hawking entropy) (Vidal *et al* '03; Plenio *et al* '05, ...) XY model, quasi-free bosonic and fermionic models, ... (Holzhey *et al* '94; Calabrese, Cardy '04) critical systems described by CFT (log correction)

(Aharonov et al '09; Irani '10) 1D model with volume scaling of entanglement entropy!

Why is Area Law Interesting?

- Connection to Holography.
- Interesting to study entanglement in physical states with an eye on quantum information processing.
- Area law appears to be connected to our ability to writedown simple Ansatzes for the quantum state.
 (e.g. tensor-network states: PEPS, MERA)

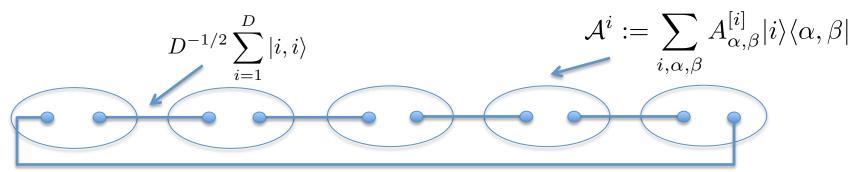
This is known rigorously in 1D:

Matrix Product States

(Fannes, Nachtergaele, Werner '92; Affleck, Kennedy, Lieb, Tasaki '87)

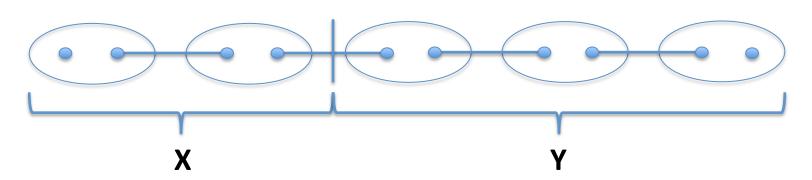
$$|\psi\rangle_{1,...,n} = \sum_{i_1=1}^{2} ... \sum_{i_n=1}^{2} tr(A_{i_1}^{[1]}...A_{i_n}^{[n]})|i_1,...,i_n\rangle, A_j^{[l]} \in Mat(D,D)$$

D: bond dimension



- Only nD² parameters.
- Local expectation values computed in nD³ time
- Variational class of states for powerful DMRG (White ')
- Generalization of product states (MPS with D=1)

MPS Area Law



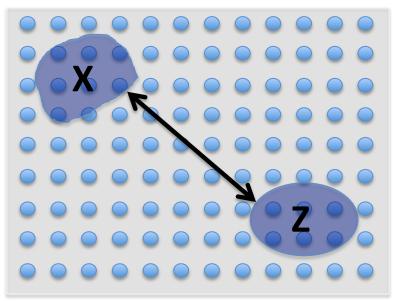
- For MPS, $S(\rho_x) \leq log(D)$
- (Vidal '03; Verstraete, Cirac '05)

If ψ satisfies $S(\rho_X) \le \log(D)$ for all X, then it has a MPS description of bond dim. D

(obs: must use Renyi entropies)

Correlation Length





Correlation Function:

$$cor(X:Z)_{\psi} := \max_{M \in X, N \in Y} \langle \psi | M \otimes N | \psi \rangle - \langle \psi | M | \psi \rangle \langle \psi | N | \psi \rangle$$

Correlation Length: $|\psi|$ has correlation length ξ if for every regions X, Z:

$$cor(X:Z)_{\psi} \leq 2^{-dist(X,Z)/\xi}$$

$$dist(X, Z) := \min_{x \in X, z \in Z} dist(x, z)$$

When there is a finite correlation length?

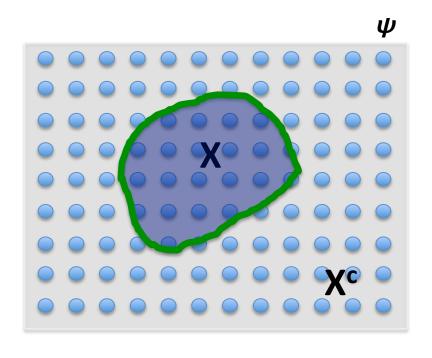
```
(Hastings '04) In any dim at zero temperature for gapped models (for groundstates; \xi = O(1/gap))
```

(Hastings '11; Hamza et al '12; ...) In any dim for models with mobility gap (many-body localization)

```
(Araki '69) In 1D at any finite temperature T (for \rho = e^{-H/T}/Z; \xi = O(1/T))
```

(Kliesch et al '13) In any dim at large enough T

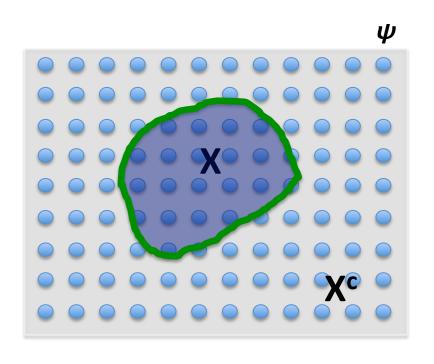
(Kastoryano et al '12) Steady-state of fast converging dissipative processes (e.g. gapped Liovillians)



$$\operatorname{Ent}(X:X^c) \approx \sum_{i \in X^c} \operatorname{Ent}(X:X^c_i)$$

$$\approx \sum_{i \in X^c} 2^{-\operatorname{dist}(X,i)/\xi}$$

$$= O(\operatorname{Area}(X)\xi)$$



$$\operatorname{Ent}(X:X^c) \approx \sum_{i \in X^c} \operatorname{Ent}(X:X^c_i)$$

$$\approx \sum_{i \in X^c} 2^{-\operatorname{dist}(X,i)/\xi}$$

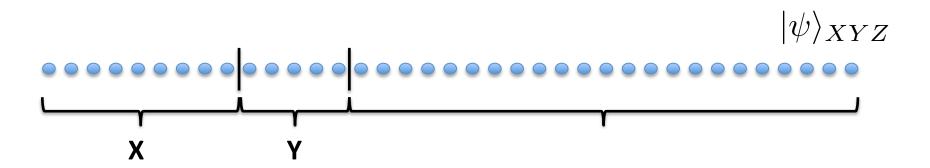
$$= O(\operatorname{Area}(X)\xi)$$

That's incorrect!

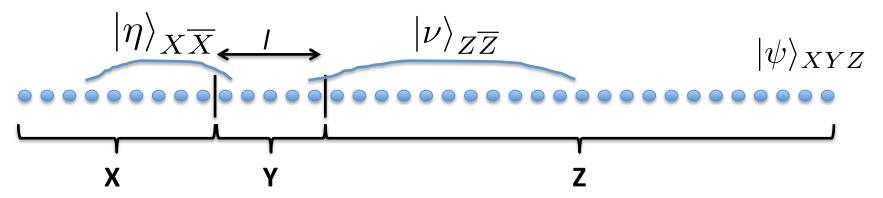
Ex. For almost every n qubit state: $S(X) \approx \operatorname{vol}(X)$

and for all i in \mathbf{X}^{c} , $\|\rho_{XX_{i}^{c}}-\rho_{X}\otimes\rho_{X_{i}}^{c}\|_{1}\leq2^{-cn}$

Entanglement can be *scrambled*, non-locally encoded (e.g. QECC, Topological Order)



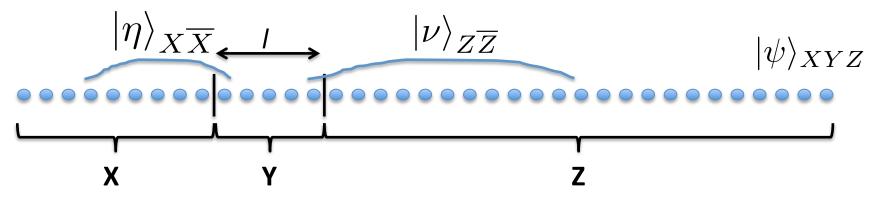
Suppose $\rho_{XZ} = \rho_X \otimes \rho_Z$.



Suppose
$$\rho_{XZ} = \rho_X \otimes \rho_Z$$
.

Then
$$|\psi\rangle_{XYZ}=(U_{\overline{XZ}\to Y}\otimes I_{XZ})|\eta\rangle_{X\overline{X}}\otimes |\nu\rangle_{Z\overline{Z}}$$

X is only entangled with Y

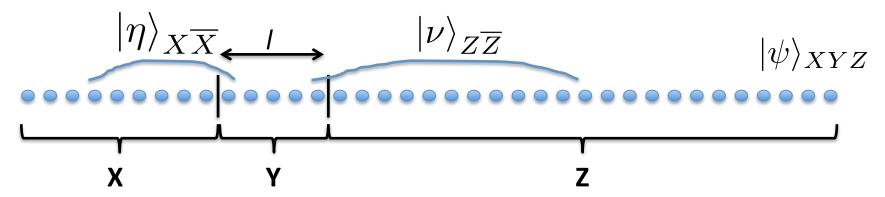


Suppose
$$ho_{XZ}=
ho_X\otimes
ho_Z$$

Then
$$|\psi\rangle_{XYZ}=(U_{\overline{XZ}\to Y}\otimes I_{XZ})|\eta\rangle_{X\overline{X}}\otimes |\nu\rangle_{Z\overline{Z}}$$

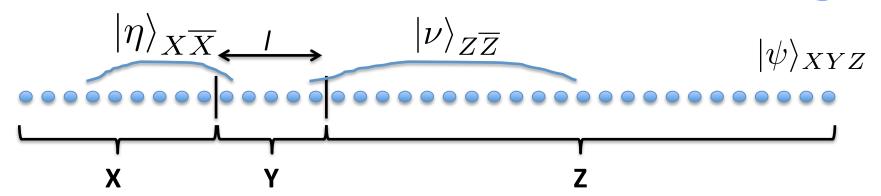
X is only entangled with Y

What if merely $cor(X:Z) \leq 2^{-l/\xi}$?



Suppose
$$ho_{XZ} pprox
ho_X \otimes
ho_Z$$

Then
$$|\psi\rangle_{XYZ} \approx (U_{\overline{XZ}\to Y} \otimes I_{XZ})|\eta\rangle_{X\overline{X}} \otimes |\nu\rangle_{Z\overline{Z}}$$



Suppose
$$ho_{XZ} pprox
ho_X \otimes
ho_Z$$

Then
$$|\psi\rangle_{XYZ} \approx (U_{\overline{XZ}\to Y} \otimes I_{XZ})|\eta\rangle_{X\overline{X}} \otimes |\nu\rangle_{Z\overline{Z}}$$

True (Uhlmann's thm). But we need 1-norm (trace-distance):

$$\rho_{XZ} \approx \rho_X \otimes \rho_Z \iff \|\rho_{XZ} - \rho_X \otimes \rho_Z\|_1 \le \varepsilon$$
$$\|\rho_{XZ} - \rho_X \otimes \rho_Z\|_1 := 2 \max_{0 \le M \le I} \operatorname{tr}(M(\rho_{XZ} - \rho_X \otimes \rho_Z))$$

In contrast
$$\operatorname{cor}(X,Z) := \max_{A,B} \operatorname{tr}((A \otimes B)(\rho_{XZ} - \rho_X \otimes \rho_Z))$$

Data Hiding States

Well distinguishable globally, but poorly distinguishable locally

(DiVincenzo, Hayden, Leung, Terhal '02)

Ex. 1 Antisymmetric Werner state $\omega_{AB} = (I - F)/(d^2-d)$

$$\operatorname{cor}(A:B) \le 1/d \qquad \|\omega_{AB} - \omega_A \otimes \omega_B\|_1 \approx 1/2$$

Data Hiding States

Well distinguishable globally, but poorly distinguishable locally

(DiVincenzo, Hayden, Leung, Terhal '02)

Ex. 1 Antisymmetric Werner state $\omega_{AB} = (I - F)/(d^2-d)$

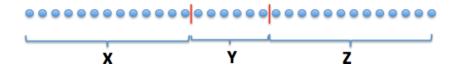
$$cor(A:B) \le 1/d$$

$$\operatorname{cor}(A:B) \le 1/d \qquad \|\omega_{AB} - \omega_A \otimes \omega_B\|_1 \approx 1/2$$

Ex. 2 Random state $|\psi\rangle_{XYZ}$ with |X|=|Z| and |Y|=|Z|

$$cor(X:Z) \le 2^{-cl}$$

$$S(X) \approx (n-l)/2$$



Data Hiding States

Well distinguishable globally, but poorly distinguishable locally

(DiVincenzo, Hayden, Leung, Terhal '02)

Ex. 1 Antisymmetric Werner state $\omega_{AB} = (I - F)/(d^2-d)$

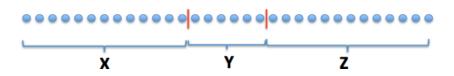
$$cor(A:B) \le 1/d$$

$$\operatorname{cor}(A:B) \le 1/d \qquad \|\omega_{AB} - \omega_A \otimes \omega_B\|_1 \approx 1/2$$

Ex. 2 Random state $|\psi\rangle_{_{XYZ}}$ with |X|=|Z| and |Y|=|Z|

$$cor(X:Z) \le 2^{-cl}$$

$$S(X) \approx (n-l)/2$$

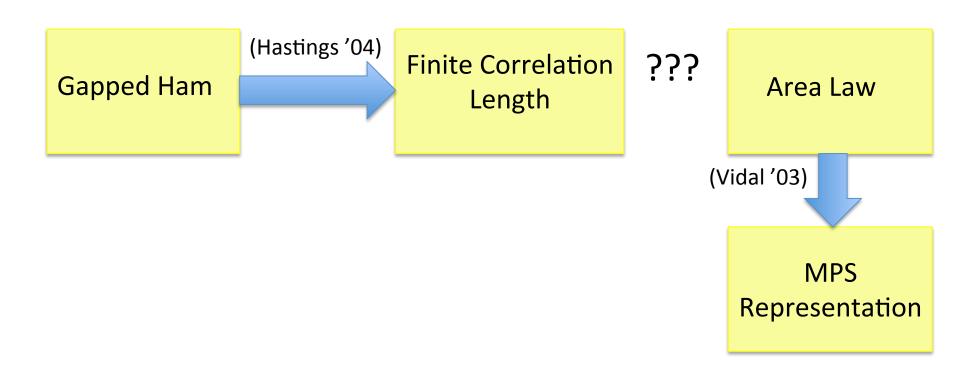


Ex. 3 (Hastings '07) Quantum Expanders States: States with big entropy but s.t. for every regions X, Z far away from edge

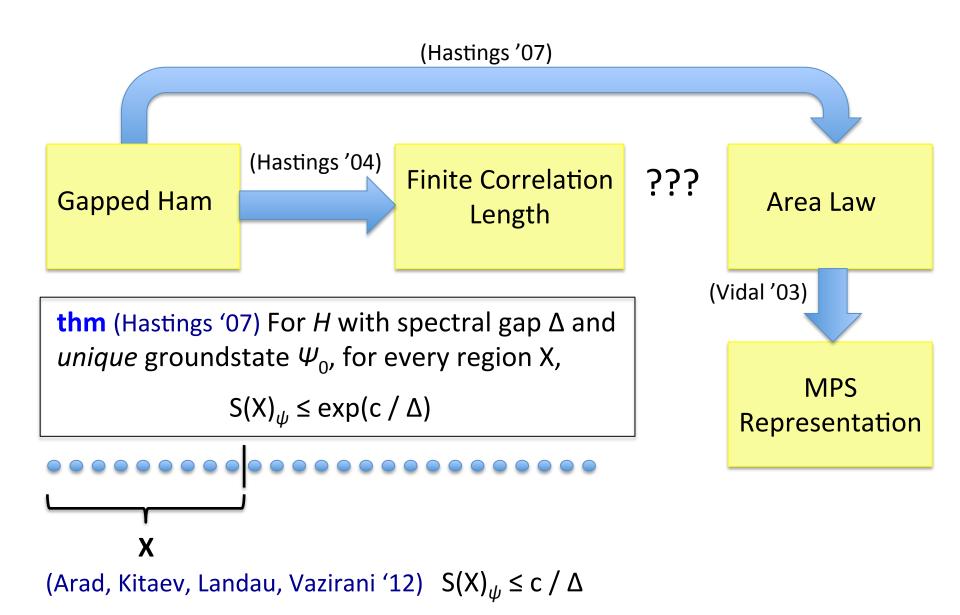
$$cor(X:Z) \le 2^{-c \operatorname{dist}(X,Z)}$$



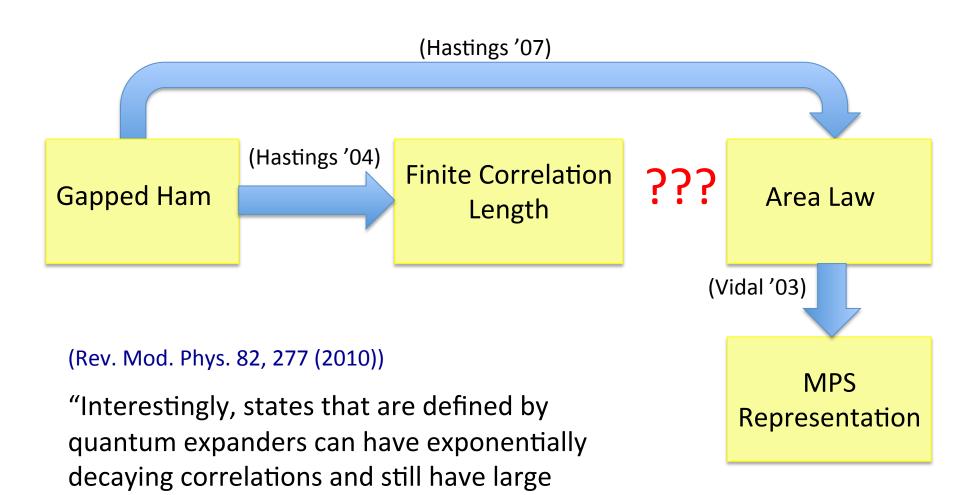
Area Law in 1D?



Area Law in 1D



Area Law in 1D

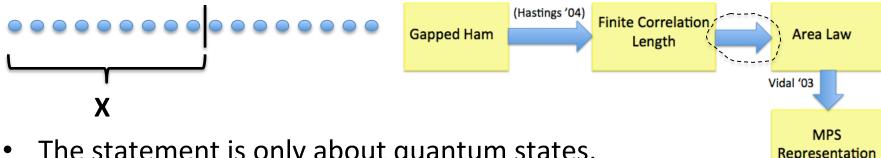


entanglement, as has been proven in (...)"

Correlation Length vs Entanglement

thm 1 (B., Horodecki '12) Let $|\psi\rangle_{1,...,n}$ be a quantum state in 1D with correlation length ξ . Then for every X,

$$S(X)_{\psi} \leq 2^{c\xi}$$



- The statement is only about quantum states, no Hamiltonian involved.
- Applies to degenerate groundstates, and gapless models with finite correlation length (e.g. systems with mobility gap; many-body localization)

Summing Up

Area law always holds in 1D whenever there is a finite correlation length:

- Groundstates (unique or degenerate) of gapped models
- Eigenstates of models with mobility gap (many-body localization)
- Thermal states at any non-zero temperature
- Steady-state of gapped dissipative dynamics

Implies that in all such cases the state has an efficient classical parametrization as a MPS

(Useful for numerics – e.g. DMRG. Limitations for quantum information processing e.g. no-go for adiabatic quantum computing in 1D)

Proof Idea

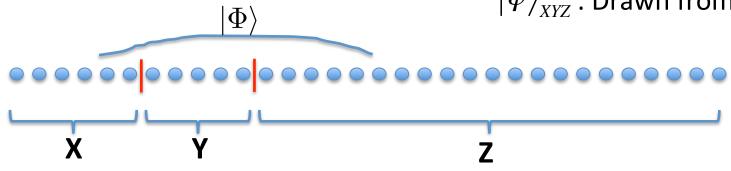


We want to bound the entropy of X using the fact the correlation length of the state is finite.

Need to relate entropy to correlations.

Random States Have Big Correl.

 $|\psi
angle_{\scriptscriptstyle
m VVZ}$: Drawn from Haar measure



Let size(XY) < size(Z). W.h.p.
$$\|\rho_{XY} - \tau_X \otimes \tau_Y\|_1 \le 2^{-\Omega(n)}$$
, $\tau_X := \mathrm{id}/d_X$

X is **decoupled** from Y.

Extensive entropy, but also large correlations:

$$U_{Z \to Z_1 Z_2} |\psi\rangle_{XYZ} \approx |\Phi\rangle_{XZ_1} \otimes |\Phi\rangle_{YZ_2}$$

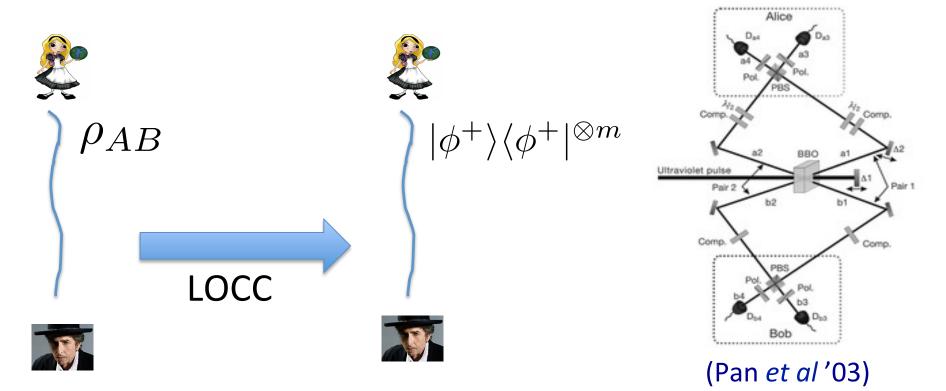
 $|\Phi
angle_{\scriptscriptstyle XZ_{\scriptscriptstyle 1}}$: Maximally entangled state between XZ₁.

 $Cor(X:Z) \ge Cor(X:Z_1) = 1/4 >> 2^{-\Omega(n)} : long-range correlations$

Entanglement Distillation

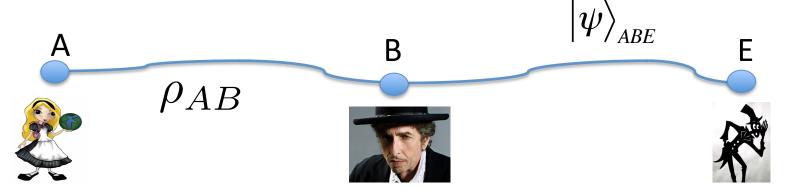
Consists of extracting EPR pairs from bipartite entangled states by Local Operations and Classical Communication (LOCC)

Central task in quantum information processing for distributing entanglement over large distances (e.g. entanglement repeater)



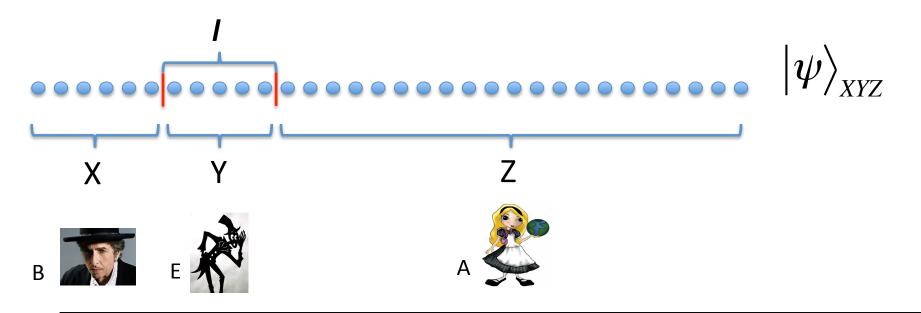
Entanglement Distillation Protocol

We apply entanglement distillation to show large entropy implies large correlations



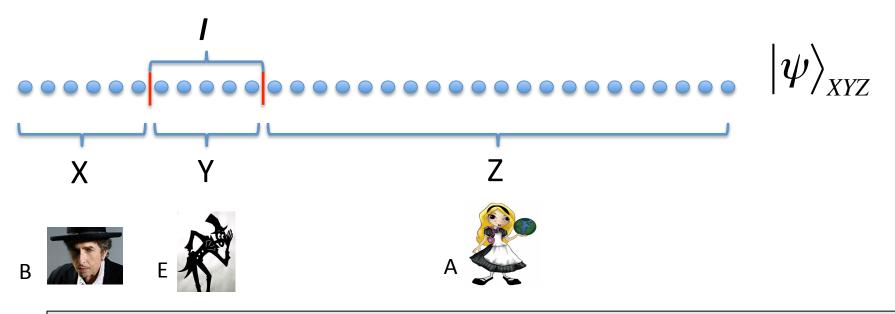
Entanglement distillation: Given $|\psi\rangle_{ABE}$ Alice can distill -S(A|B) = S(B) - S(AB) EPR pairs with Bob by making a measurement with N $\approx 2^{I(A:E)}$ elements, with I(A:E) := S(A) + S(E) - S(AE), and communicating the outcome to Bob. (Devetak, Winter '04)

Distillation Bound



$$S(X) > S(Y) \Rightarrow Cor(X : Z) \ge O(2^{-I(X : Y)})$$

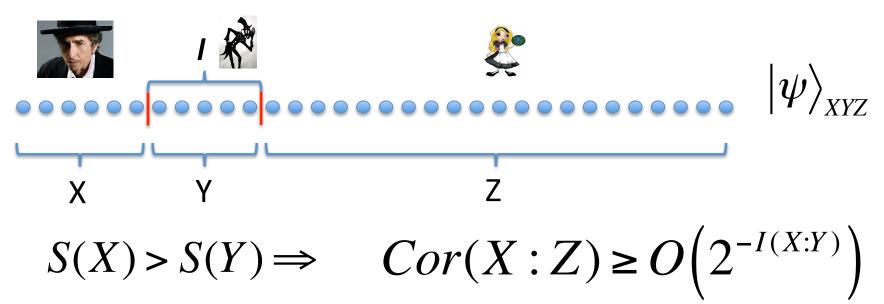
Distillation Bound

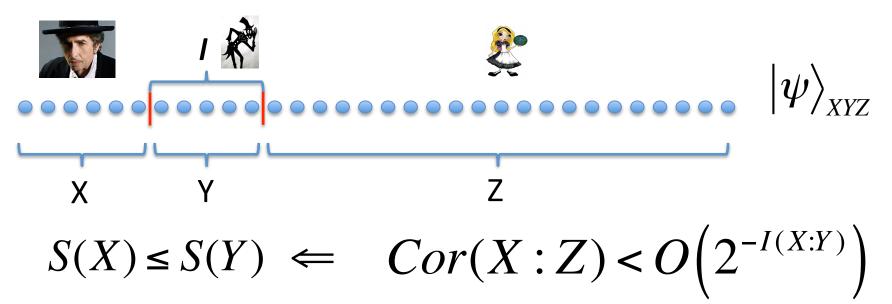


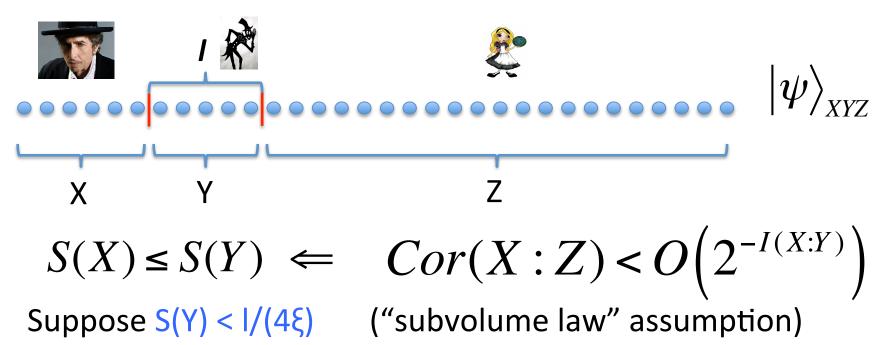
$$S(X) > S(Y) \Rightarrow Cor(X : Z) \ge O(2^{-I(X:Y)})$$

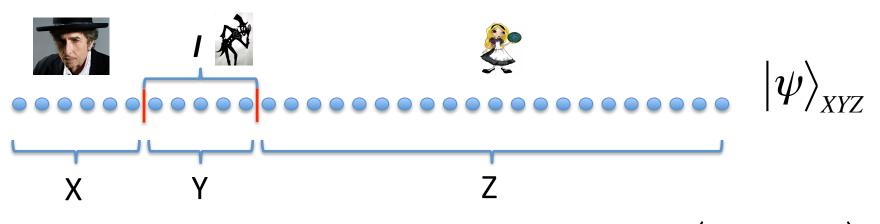
S(X) - S(XZ) > 0 (EPR pair distillation rate)

Prob. of getting one of the $2^{I(X:Y)}$ outcomes









$$S(X) \le S(Y) \iff Cor(X:Z) < O\left(2^{-I(X:Y)}\right)$$

Suppose $S(Y) < I/(4\xi)$ ("subvolume law" assumption)

Since $I(X:Y) < 2S(Y) < I/(2\xi)$, a correlation length ξ implies

 $Cor(X:Z) < 2^{-I/\xi} < 2^{-I(X:Y)}$

Thus: S(X) < S(Y)

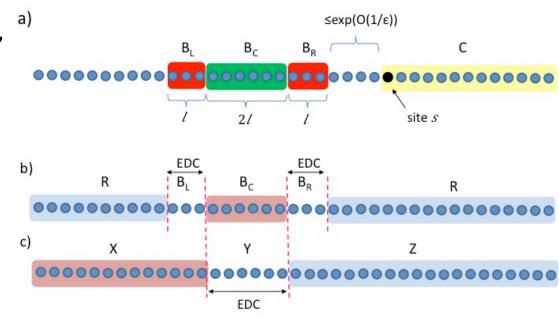
Actual Proof

We apply the bound from entanglement distillation to prove finite correlation length -> Area Law in 3 steps:

- c. Get area law from finite correlation length under assumption there is a region with "subvolume law"
- b. Get region with "subvolume law" from finite corr. length and assumption there is a region of "small mutual information"
- a. Show there is always a region of "small mutual info"

Each step uses the assumption of finite correlation length.

Obs: Must use single-shot info theory (Renner et al)



Area Law in Higher Dim?

Wide open...

Known proofs in 1D (for groundstates gapped models):

- Hastings '07. Analytical (Lieb-Robinson bounds, Fourier analysis,...)
- 2. Arad, Kitaev, Landau, Vazirani '13. Combinatorial (Chebyshev polynomial, ...)
- 3. B., Horodecki '12 (this talk). Information-Theoretical

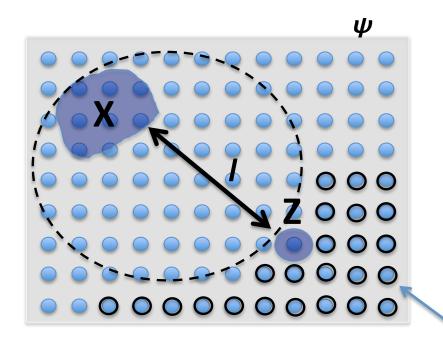
All fail in higher dimensions....

Area Law in Higher Dim?

New Approach:

"Conditional Correlation length":

$$\mathbb{E}_{i_1,...,i_k} \operatorname{cor}(X:Z)_{\psi_{i_1,...,i_k}} \le 2^{-l/\xi}$$



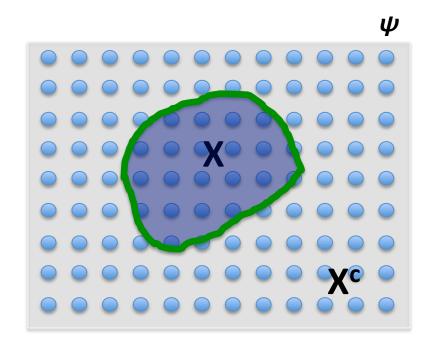
 $\psi_{i_1,...,i_k}$: post-measured state after measurement on sites (a₁,..., a_k) with outcomes (i₁, ..., i_k) in {0, 1}^k

$$|\psi_{i_1,\ldots,i_k}\rangle := \frac{1}{\|\ldots\|} (\mathrm{id} \otimes |i_k\rangle\langle i_k|_{a_k} \otimes \ldots \otimes |i_1\rangle\langle i_1|_{a_1}) |\psi\rangle$$

Measurement on site a_k

Area Law from Finite *Conditional*Correlation length

thm (B. '14) In any dim, if ψ has conditional correlation length ξ , then $S(X)_{\psi} \leq 4\xi \text{ Area}(X)$



Which states have a finite conditional correlation length?

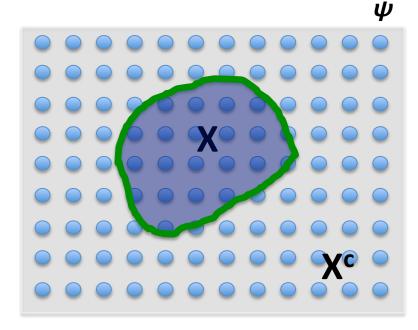
Conjecture1: Any groundstate of gapped local Hamiltonian.

Conjecture 2: Any state with a finite correlation length.

Obs: Can prove it for 1D models (finite CL -> area law -> MPS -> finite CCL)

Area Law from Finite *Conditional*Correlation length

thm (B. '14) In any dim, if ψ has conditional correlation length ξ , then $S(X)_{\cup} \le 4\xi \text{ Area}(X)$



Proof by quantum information theory:

$$2S(X)_{\psi} = I(X : X^{c})_{\psi}$$

$$\leq I(X : X^{c})_{\rho} + 2\xi \operatorname{Area}(X)$$

conditional corr. length

$$\rho = \mathrm{id}_X \otimes \Lambda^{\otimes \mathrm{size}(X^c)}(|\psi\rangle\langle\psi|)$$

$$\Lambda(X) = \langle 0|X|0\rangle|0\rangle\langle 0| + \langle 1|X|1\rangle|1\rangle\langle 1|$$

Application to Systems with Robust Gap

thm (B. '14) In any dim, if ψ has conditional correlation length ξ , then $S(X)_{\psi} \leq 4\xi \text{ Area}(X)$

(Verstraete '14) Groundstates of Hamiltonians with local topological order have finite conditional correlation length.

LTQO : Closely related to "robust gap", i.e. $H+\varepsilon\sum_k V_k$ is gapped for ε small enough and all $V_{\bf k}$.

cor Every groundstate of a system with local topological order fulfills area law

Application to Systems with Robust Gap

thm (B. '14) In any dim, if ψ has conditional correlation length ξ , then $S(X)_{\psi} \leq 4\xi \ Area(X)$

(Verstraete '14) Groundstates of Hamiltonians with local topological order have finite conditional correlation length.

LTQO : Closely related to "robust gap", i.e. $H+\varepsilon\sum_k V_k$ is gapped for ε small enough and all $V_{\bf k}$.

cor Every groundstate of a system with local topological order fulfills area law

Improves on (Michalakis, Pytel '11) who proved $S(X) \le Area(X)log(vol(X))$.

Obs: *Strict* area law is important, as it allows us to define the concept of topological entanglement entropy (Kitaev, Preskill '05, Levin, Wen '05)

Summary

- Finite correlation length gives an area law for entanglement in 1D. We don't know what happens in higher dimensions.
- More generally, thinking about entanglement from the perspective of quantum information theory is useful.
- Growing body of connections between concepts/techniques in quantum information science and other areas of physics.

Thanks!