Theory and Algorithms for Weighted Total Least-Squares Fitting of Lines, Planes, and Parallel Planes to Support Tolerancing Standards

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We present the theory and algorithms for fitting a line, a plane, two parallel planes (corresponding to a slot or a slab), or many parallel planes in a total (orthogonal) least-squares sense to coordinate data that is weighted. Each of these problems is reduced to a simple 3×3 matrix eigenvalue/eigenvector problem or an equivalent singular value decomposition problem, which can be solved using reliable and readily available commercial software. These methods were numerically verified by comparing them with brute-force minimization searches. We demonstrate the need for such weighted total least-squares fitting in coordinate metrology to support new and emerging tolerancing standards, for instance, ISO 14405-1:2010. The widespread practice of unweighted fitting works well enough when point sampling is controlled and can be made uniform (e.g., using a discrete point contact coordinate measuring machine). However, we show by example that nonuniformly sampled points (arising from many new measurement technologies) coupled with unweighted least-squares fitting can lead to erroneous results. When needed, the algorithms presented also solve the unweighted cases simply by assigning the value one to each weight. We additionally prove convergence from the discrete to continuous cases of least-squares fitting as the point sampling becomes dense. [DOI: 10.1115/1.4024854]

1 Introduction

The need for weighted total (here, orthogonal distance) leastsquares fitting of lines, planes, and parallel planes comes from at least two fronts: new tolerancing standards and new coordinate measuring instrumentation [1]. First, ISO tolerancing standards such as ISO 14405-1:2010 [2] and the future revisions of ISO 1101 [3], and ISO 14405-3 [4] will allow specifications to be called out based on a total least-squares criterion (see Table 1). As will be shown in Sec. 2, these must be interpreted as weighted total least-squares², where the weights are the discretely partitioned lengths/areas corresponding to the discretely sampled and measured points. This will be shown to arise from the definition of the tolerance specification, which applies to the continuous surface and which, according to the duality principle [5], defines the measurand in the verification process. The verification process, generally performed using discrete points, must correspond to the specification or suffer the appropriate additional uncertainty for any departure. Thus, there is an immediate need for weighted total least-squares fitting of lines, planes, and parallel planes.

Verification of slot and slab size specifications requires fitting two parallel planes. One can also see the need for fitting several parallel planes, as in the case a door hinge [6], where there are

¹The case of fitting parallel lines in two-dimensions is also solved in this paper as a simple restriction of the parallel planes fit to 2D. The case of fitting parallel lines in three-dimensions is solved by going from the single line solution to multiple parallel lines with exactly analogous methods as going from single plane solution to multiple parallel planes as shown in this paper. The reason this case is not explicitly dealt with in this paper is due to the lack of immediate application to tolerancing standards.

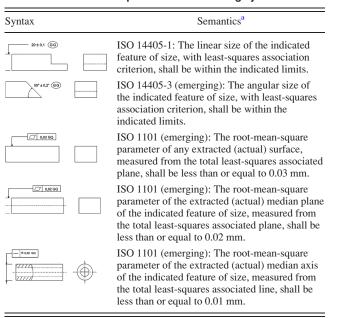
²Including the word "weighted" before total least-squares might be thought of as redundant, but total least-squares is often employed to weigh the coordinates of each point differently from each other. Here, the points themselves are weighted differently from each other, so we explicitly say weighted, even though the most general application of total least-squares allows for all such weights.

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several nominally parallel planes (perpendicular to multiple cylindrical surfaces representing holes that are nominally coaxial). Under such fitting, tolerancing a hinge can be treated as a one-dimensional tolerancing problem that is by no means trivial.

The second reason for weighted (as opposed to unweighted) total least-squares fitting arises from newer measuring instrumentation. Discrete point coordinate measuring machines (CMMs) can be programmed to sample a surface predictably and

Table 1 Examples of ISO tolerancing syntax



^aThese statements of semantics are composed from different statements in ISO 14405-1:2010 and the emerging ISO 1101. These are not the formal statements of explanation associated with the drawings in the official ISO standards.

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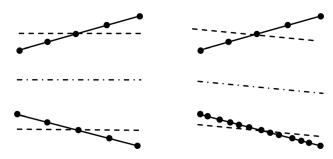


Fig. 1 Point sampling density adversely affecting the leastsquares fit

uniformly. That is, before any CMM probing occurs, the number of points and their approximate sampling pattern can already be known, and the operator generally chooses the points. This luxury allows for more or less even sampling of surfaces that are then associated with ideal form geometries based on an unweighted total least-squares fitting criterion.

But the introduction of newer measurement technologies takes away such a priori knowledge. An operator using—for examplean articulating arm CMM with a handheld laser scanner does not know how many points will be collected or, to a certain extent, how uniform the sampling will be (even considering deliberate attempts by the operator). The ramifications of this are important. Consider the problem of fitting two parallel planes to a slot—a task that can arise from a tolerance specification according to ISO 14405-1, which can specifically indicate a least-squares criterion. In this example, we suppose that the actual planar surfaces are not quite parallel. Figure 1 shows the effect that variations in sampling can have on the fit. In the figure, the unweighted leastsquares parallel planes are shown by the dashed lines when fit to points having uniform sampling (left) and nonuniform sampling (right) (in this figure, the sampled points happen to lie exactly on the surface, but this does not affect the idea conveyed).

As Fig. 1 shows, when the sampling is uniform all over, the orientation of the fit planes matches what is expected and desired, namely the orientation matches what one would get in the continuous case. But when the sampling is not uniform all over, and the number of points on one planar surface is much higher than on the other, the orientation of the fit planes is skewed, an effect that deviates from the continuous case. The use of weighted least-squares (where the weights correspond to the length or area around each point) avoids this effect of sampling.

Note that this problem still exists when greater numbers of points are taken, provided the relative disparity remains between the numbers of points sampled on the two surfaces. And this case is realistic, as optically based scanners can gather much more data on one surface than another based on various things such as distance, sampling time, lighting, and surface reflectivity.

An easily grasped example is the use of a 3D laser scanner that collects points on two nominally parallel planar surfaces of the same nominal size. If one surface is a distance of 2 m from the instrument while the other is 10 m from the instrument, the inverse square law would have us expect there to be 25 times as many points collected on the closer surface than the farther one (apart from special handling). An unweighted least-squares parallel planes fit would be almost entirely determined by the orientation of the closer plane. This could be remedied by the algorithms presented in this paper, if the points on the closer surface were given weights 1/25 (or whatever the actual ratio of points turned out to be) of the weights given to the points on the farther surface.

The problem can exist even when fitting a single plane to points taken on one surface, if one patch is sampled more densely than the rest of the surface. The characteristics of that patch would have undue influence on the fit when using an unweighted least-squares algorithm.

While in this paper, we will often associate the weights with the length/area of the arc/patch of the curve/surface corresponding

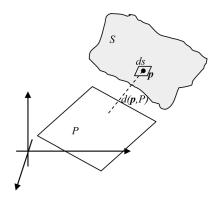


Fig. 2 Fitting a plane to a surface patch

to each point, the scope of this document does not include specific algorithms for calculating the weights, which is an interesting topic in itself. Further, making use of weights may be desired for other reasons such as sampled points with differing uncertainties or even specific coordinates with differing uncertainties. We also note that the case of unweighted least-squares fitting of a line or single plane has been well documented, for example in Refs. [7–11]. While not in the context of metrology, a general presentation of total least-squares is given in Refs. [12,13]. An iterative search algorithm for total least-squares line fitting is given in Ref. [14], where in their 2D problem, different weights are given to the individual coordinates of each point, based on the uncertainties of those coordinates.

The remainder of this paper is organized as follows. Section 2 poses the problem in the continuous cases and shows why this lends itself to weighted fitting in the discrete cases. Section 3 gives solutions to the continuous cases. Section 4 gives the algorithms for discrete cases separate from any proofs for convenient access of the reader. Section 5 contains proofs the discrete-case algorithms presented. Section 6 gives MATLAB³ code of the algorithms along with results of numerical testing. Section 7 contains the proof that the results from the discrete algorithms converge to their continuous counterparts as the points become dense, and Sec. 8 gives some concluding remarks.

2 Definitions in Continuous Cases and Their Discrete Approximations

As described in Ref. [15], to fit a total least-squares plane to a surface patch in space, we pose the following optimization problem (with reference to Fig. 2):

2.1 TlsqPlane. Given a bounded surface *S*, find the plane *P* that minimizes $\int_S d^2(\mathbf{p}, P) ds$

Here d(p, P) denotes the signed perpendicular distance (hence, the qualification "total" for the least-squares fitting) of a point p on surface patch S from the plane P that will be fitted. Once such a plane P has been found, the root-mean-square (rms) parameter for the bounded surface S is given by

$$\sqrt{\frac{\int_{S} d^{2}(\boldsymbol{p}, P)ds}{\int_{S} ds}} \tag{1}$$

We note that $\int_S ds$ is the area of the surface patch. If the surface consists of several patches, then the integrals can be evaluated over each patch and then summed.

031008-2 / Vol. 13, SEPTEMBER 2013

³Certain commercial software packages are identified in this paper in order to specify the experimental procedures and code adequately. Such identification is not intended to imply recommendation or endorsement by the National Institute of Standards and Technology, nor is it intended to imply that the software tools identified are necessarily the best available for the purpose.

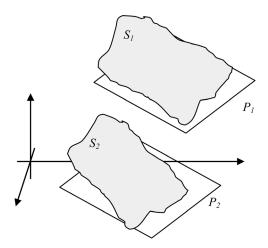


Fig. 3 Fitting two parallel planes

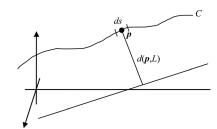


Fig. 4 Fitting a line to a curve

The objective function in TlsqPlane cannot, in general, be evaluated in closed form. So we resort to numerical integration over the surface S. We can sample points on a surface patch after dividing up the patch into discrete areas ΔA_i and approximate the objective function as

$$\int_{S} d^{2}(\boldsymbol{p}, P) ds \approx \sum_{i=1}^{N} \left\{ d^{2}(\boldsymbol{p}_{i}, P) \right\} \cdot \Delta A_{i}$$
 (2)

where p_i are the N sampled points, one in each subdivision. Thus, we are led to minimizing $\sum_{i=1}^{N} \left[\left\{ d^2(p_i, P) \right\} \cdot \Delta A_i \right]$ over the parameters of the plane P for TlsqPlane, where ΔA_i 's are treated as the weights.

When we need to fit two or more parallel planes, the problem can be formulated as follows. For simplicity, we will present the case of fitting two parallel planes (with corresponding illustration in Fig. 3):

2.2 TIsqParallelPlanes. Given two bounded surfaces S_1 and S_2 , find two parallel planes P_1 and P_2 that minimize $\int_{S_1} d^2(\mathbf{p}, P_1) ds + \int_{S_2} d^2(\mathbf{p}, P_2) ds$

Once such parallel planes P_1 and P_2 have been found, the distance between them is the linear size (that has been defined with the least-squares criterion) in ISO 14405-1:2010. The definition can be extended to any arbitrary number of parallel planes. A discrete approximation to the TlsqParallelPlanes problem can be defined as we did for the TlsqPlane problem.

For fitting a line, see Fig. 4.

TIsqLine. Given a bounded space curve C, find the line L that minimizes $\int_C d^2(\mathbf{p}, L) ds$

We can sample points on C after dividing up the curve into discrete segments ΔC_i and approximate the objective function as $\int_C d^2(\mathbf{p}, L) ds \approx \sum_{i=1}^N \left[d^2(\mathbf{p}_i, L) \right] \cdot \Delta C_i$.

3 Solutions for Continuous Cases

We begin by solving the continuous cases; the discrete cases follow from these in a straightforward manner. First, we give the relevant distance formulas. If a plane is defined by a point on the plane x = (x, y, z) and the direction cosines of the normal to the plane (i.e., the unit vector normal to the plane) a = (a, b, c), then the signed orthogonal distance from a point $x_i = (x_i, y_i, z_i)$ to the plane is given by:

$$d_i = \mathbf{a} \cdot (\mathbf{x_i} - \mathbf{x}) = a(x_i - \mathbf{x}) + b(y_i - \mathbf{y}) + c(z_i - z)$$

If a line is defined by a point on the line x = (x, y, z) and the direction cosines of the line (i.e., the unit vector in the direction of the line) a = (a, b, c), then the orthogonal distance from a point $x_i = (x_i, y_i, z_i)$ to the line is given by:

$$d_i = |\boldsymbol{a} \times (\boldsymbol{x_i} - \boldsymbol{x})|$$

Written in detail:

$$d_i = \sqrt{u^2 + v^2 + w^2}$$
, where
 $u = c(y_i - y) - b(z_i - z)$
 $v = a(z_i - z) - c(x_i - x)$
 $w = b(x_i - x) - a(y_i - y)$

However, it is also equal to the Pythagorean expression below, which is less numerically stable but more convenient for use in proofs:

$$d_i = \sqrt{|x_i - x|^2 - (\boldsymbol{a} \cdot (x_i - x))^2}$$

or written another way,

$$\left(d_i^{\text{line}}\right)^2 = \left|x_i - x\right|^2 - \left(d_i^{\text{plane}}\right)^2$$

Note: In this paper d_i can refer to the distance from the *i*th data point to either a plane or line, with the context making clear which. In cases when clarification is needed, d_i^{line} and d_i^{plane} are used to remove ambiguity, as in the equation above.

We now present two theorems solving the continuous case least-squares fitting problems, along with a few helpful lemmas.

Lemma 1. Assume that we are given a bounded, piecewise continuous surface, S of finite area, and an orientation, a. Then the total least-squares plane constrained to have normal a must pass through the centroid, $\bar{x} = (\bar{x}, \bar{y}, \bar{z})$, where $\bar{x} = \int_S x ds / \int_S ds$, $\bar{y} = \int_S y ds / \int_S ds$, and $\bar{z} = \int_S z ds / \int_S ds$. Any plane having normal a but not passing through the centroid has a sum-of-squares objective function value strictly greater than the plane passing through the centroid.

Proof. The equation of any plane having normal \boldsymbol{a} can be written as $\boldsymbol{a} \cdot \boldsymbol{x} - d = 0$, where d is the signed distance from the plane to the origin. The orthogonal distance from any arbitrary point \boldsymbol{p} to this plane is $\boldsymbol{a} \cdot \boldsymbol{p} - d$. Thus, the objective function for any plane of orientation \boldsymbol{a} is given by $F(d) = \int_{S} [\boldsymbol{a} \cdot \boldsymbol{p} - d]^{2} ds$.

Taking the first and second derivatives yields:

$$F'(d) = -2 \int_{S} [\mathbf{a} \cdot \mathbf{p} - d] ds, \text{ and}$$
$$F''(d) = 2 \int_{S} ds > 0$$

The fact that the second derivative is a positive constant implies that the function is strictly convex and has a unique minimum if and where its first derivative vanishes. This occurs when $\int_S [a \cdot p - d] ds = 0$. Distributing the integral and solving for d yields

$$d = \frac{\int_{S} [\boldsymbol{a} \cdot \boldsymbol{p}] ds}{\int_{S} ds} = \boldsymbol{a} \cdot \left(\frac{\int_{S} \boldsymbol{p} ds}{\int_{S} ds} \right) = \boldsymbol{a} \cdot \bar{\boldsymbol{x}}$$

But this means that that the distance from \bar{x} (the centroid) to the plane, which is $a \cdot \bar{x} - d$, is equal to zero, implying that the centroid must lie on the total least-squares plane constrained to have normal a. Furthermore, the objective function is strictly convex,

Journal of Computing and Information Science in Engineering

meaning any other plane not passing through the centroid must have a greater objective function value.

THEOREM 1. Assume that we are given $K \ge 1$ surfaces, S_1 , S_2 ,..., S_k , where each surface is bounded, piecewise continuous, of finite area, and each S_i having centroid $\bar{\mathbf{x}}_i = (\bar{x}_i, \bar{y}_i, \bar{z}_i)$. Then a set of K

parallel planes of total least-squares can be computed as the planes passing through the respective centroids of the surfaces and all sharing the same orientation defined by the eigenvector corresponding to the smallest eigenvalue of the 3×3 matrix (written as a sum of matrices):

$$\sum_{i=1}^{K} \left| \int_{S_{i}} (x - \bar{x}_{i})^{2} ds \int_{S_{i}} (x - \bar{x}_{i})(y - \bar{y}_{i}) ds \int_{S_{i}} (x - \bar{x}_{i})(z - \bar{z}_{i}) ds \right| \int_{S_{i}} (x - \bar{x}_{i})(y - \bar{y}_{i}) ds \int_{S_{i}} (y - \bar{y}_{i})^{2} ds \int_{S_{i}} (y - \bar{y}_{i})(z - \bar{z}_{i}) ds \int_{S_{i}} (z - \bar{z}_{i})^{2} ds$$

Proof. For the sake of simplicity of notation and ease for the reader, we demonstrate the proof for the case of K = 2. Using simpler notation then, we assume two surfaces, S_A and S_B .

If two parallel planes have normal orientation a and pass through the points x_A and x_B , then the objective function to be minimized is

$$F(\boldsymbol{a}, \boldsymbol{x}_A, \boldsymbol{x}_B) = \int_{S_A} [\boldsymbol{a} \cdot (\boldsymbol{p} - \boldsymbol{x}_A)]^2 ds + \int_{S_B} [\boldsymbol{a} \cdot (\boldsymbol{p} - \boldsymbol{x}_B)]^2 ds$$

Since x_A occurs in the first integral only, and x_B occurs in the second integral only, then for any fixed orientation a we see that the minimization problem can be separated into two minimization problems as:

$$\frac{\min}{X_A, X_B} [F(x_A, x_B)] = \frac{\min}{X_A} [F(x_A)] + \frac{\min}{X_B} [F(x_B)]$$
(3)

This means that Lemma 1 can be applied to each minimization on the right hand side of Eq. (3) indicating to us that when fitting parallel planes, each least-squares plane passes through the centroid of its respective surface. We denote these centroids as \bar{x}_A , and \bar{x}_B , respectively. Knowing this, the objective function to be minimized can be written as a function of a alone:

$$F(\boldsymbol{a}) = \int_{S_A} \left[\boldsymbol{a} \cdot (\boldsymbol{p} - \bar{\boldsymbol{x}}_A) \right]^2 ds + \int_{S_B} \left[\boldsymbol{a} \cdot (\boldsymbol{p} - \bar{\boldsymbol{x}}_B) \right]^2 ds$$

To solve the orientation problem, we use the method of Lagrange multipliers. The minimum of $F(\mathbf{a})$ subject to $G(\mathbf{a}) = 0$ (where $G(\mathbf{a}) = |\mathbf{a}|^2 - 1$) occurs when $\nabla F = \lambda \nabla G$. In this case,

$$\nabla F = \begin{vmatrix} \frac{\partial F}{\partial a} \\ \frac{\partial F}{\partial b} \\ \frac{\partial F}{\partial c} \end{vmatrix} = 2 \begin{vmatrix} \int_{S_A} [\mathbf{a} \cdot (\mathbf{p} - \bar{\mathbf{x}}_A)](x - \bar{\mathbf{x}}_A) ds + \int_{S_B} [\mathbf{a} \cdot (\mathbf{p} - \bar{\mathbf{x}}_B)](x - \bar{\mathbf{x}}_B) ds \\ \int_{S_A} [\mathbf{a} \cdot (\mathbf{p} - \bar{\mathbf{x}}_A)](y - \bar{y}_A) ds + \int_{S_B} [\mathbf{a} \cdot (\mathbf{p} - \bar{\mathbf{x}}_B)](y - \bar{y}_B) ds \\ \int_{S_A} [\mathbf{a} \cdot (\mathbf{p} - \bar{\mathbf{x}}_A)](z - \bar{z}_A) ds + \int_{S_B} [\mathbf{a} \cdot (\mathbf{p} - \bar{\mathbf{x}}_B)](z - \bar{z}_B) ds \end{vmatrix}$$

Let L denote the 3×3 matrix (written as a sum of two matrices for reasons of space):

$$\left| \int_{S_A} (x - \bar{x}_A)^2 ds \right|_{S_A} (x - \bar{x}_A)(y - \bar{y}_A) ds \int_{S_A} (x - \bar{x}_A)(z - \bar{z}_A) ds \\
\int_{S_A} (x - \bar{x}_A)(y - \bar{y}_A) ds \int_{S_A} (y - \bar{y}_A)^2 ds \int_{S_A} (y - \bar{y}_A)(z - \bar{z}_A) ds \\
\int_{S_A} (x - \bar{x}_A)(z - \bar{z}_A) ds \int_{S_A} (y - \bar{y}_A)(z - \bar{z}_A) ds \int_{S_A} (z - \bar{z}_A)^2 ds
\right| \int_{S_B} (x - \bar{x}_B)^2 ds \int_{S_B} (x - \bar{x}_B)(y - \bar{y}_B) ds \int_{S_B} (x - \bar{x}_B)(z - \bar{z}_B) ds \\
+ \int_{S_B} (x - \bar{x}_B)(y - \bar{y}_B) ds \int_{S_B} (y - \bar{y}_B)^2 ds \int_{S_B} (y - \bar{y}_B)(z - \bar{z}_B) ds \\
\int_{S_B} (x - \bar{x}_B)(z - \bar{z}_B) ds \int_{S_B} (y - \bar{y}_B)(z - \bar{z}_B) ds \int_{S_B} (z - \bar{z}_B)^2 ds$$

Then the orientation of the least-squares planes arising from $\nabla F = \lambda \nabla G$ can be written as the elegant eigen-problem:

$$\boldsymbol{L} \begin{bmatrix} a \\ b \\ c \end{bmatrix} = \lambda \begin{bmatrix} a \\ b \\ c \end{bmatrix}$$

We now need to select the correct eigenvector. The normal equations can be written as follows:

$$\int_{S_A} [\boldsymbol{a} \cdot (\boldsymbol{p} - \bar{\boldsymbol{x}}_A)](x - \bar{x}_A) ds + \int_{S_B} [\boldsymbol{a} \cdot (\boldsymbol{p} - \bar{\boldsymbol{x}}_B)](x - \bar{x}_B) ds = \lambda \boldsymbol{a}$$

$$\int_{S_A} [\boldsymbol{a} \cdot (\boldsymbol{p} - \bar{\boldsymbol{x}}_A)](y - \bar{y}_A) ds + \int_{S_B} [\boldsymbol{a} \cdot (\boldsymbol{p} - \bar{\boldsymbol{x}}_B)](y - \bar{y}_B) ds = \lambda \boldsymbol{b}$$

$$\int_{S_A} [\boldsymbol{a} \cdot (\boldsymbol{p} - \bar{\boldsymbol{x}}_A)](z - \bar{z}_A) ds + \int_{S_B} [\boldsymbol{a} \cdot (\boldsymbol{p} - \bar{\boldsymbol{x}}_B)](z - \bar{z}_B) ds = \lambda \boldsymbol{c}$$

031008-4 / Vol. 13, SEPTEMBER 2013

Multiplying these equations by a, b, and c, respectively, then summing the equations gives

$$\int_{S_A} \left[\boldsymbol{a} \cdot (\boldsymbol{p} - \bar{\boldsymbol{x}}_{\boldsymbol{A}}) \right]^2 ds + \int_{S_B} \left[\boldsymbol{a} \cdot (\boldsymbol{p} - \bar{\boldsymbol{x}}_{\boldsymbol{B}}) \right]^2 ds = \lambda |\boldsymbol{a}|^2 = \lambda \tag{4}$$

But the sum on the left is just the objective function, F(a), hence the correct eigenvector for the solution corresponds to the smallest eigenvalue (since $F(a) = \lambda$, and since we seek to minimize F).

Lemma 2. Assume that we are given a bounded, piecewise continuous space curve, C, of finite length. Then the point that minimizes the sum-of-squares objective function, $F(\mathbf{x_0}) = \int_C |\mathbf{p} - \mathbf{x_0}|^2 ds$ is the centroid, $\bar{\mathbf{x}} = (\bar{x}, \bar{y}, \bar{z})$, where $\bar{x} = \int_C x ds / \int_C ds$, $\bar{y} = \int_C y ds / \int_C ds$, and $\bar{z} = \int_C z ds / \int_C ds$. Any point that is not the centroid has a sum-of-squares objective function value strictly greater than that of the centroid.

Proof. Setting $\nabla F = 0$ yields $\int_C (x - x_0) ds = \int_C (y - y_0) ds$ = $\int_C (z - z_0) ds = 0$. Solving for (x_0, y_0, z_0) gives $(\int_C x ds / \int_C ds, \int_C y ds / \int_C ds, \int_C z ds / \int_C ds)$, which is simply \bar{x} .

The Hessian matrix of
$$F$$
 is $2\begin{bmatrix} \int_C ds & 0 & 0 \\ 0 & \int_C ds & 0 \\ 0 & 0 & \int_C ds \end{bmatrix}$, which is immediately seen as positive definite, implying that the objective

immediately seen as positive definite, implying that the objective function is strictly convex and that the centroid is the unique point that minimizes F.

Lemma 3. Assume that we are given a bounded, piecewise continuous space curve, C, of finite length, and an orientation, a. Then the total least-squares line constrained to have direction a must pass through the centroid, $\bar{x} = (\bar{x}, \bar{y}, \bar{z})$, where $\bar{x} = \int_C x ds / \int_C ds$, $\bar{y} = \int_C y ds / \int_C ds$, and $\bar{z} = \int_C z ds / \int_C ds$. Any line having direction a but not passing through the centroid has a sum-of-squares objective function value strictly greater than the line passing through the centroid.

Proof. Suppose not. That is, suppose the least-squares line does not pass through the centroid, \bar{x} . We construct a plane perpendicular to the line and passing through \bar{x} . Call the point where the line intersects this plane, x^* (noting that $x^* \neq \bar{x}$). The remainder of the proof will show that the objective function of the line having direction a and passing through x^* is strictly greater than the objective function of the parallel line passing through \bar{x} , which contradicts the original supposition, since the line supposed to minimize the objective function indeed does not.

The objective function of the supposed least-squares solution is $F(x^*) = \int_c \left[|p-x^*|^2 - \left(d^{\text{plane}}(p,x^*) \right)^2 \right] ds$. The objective function of the parallel line passing through \bar{x} is $F(\bar{x}) = \int_c \left[|p-\bar{x}|^2 - \left(d^{\text{plane}}(p,\bar{x}) \right)^2 \right] ds$. However, since the same constructed plane with normal a contains both x^* and \bar{x} , we see that $d^{\text{plane}}(p,\bar{x}) = d^{\text{plane}}(p,x^*)$ for all points, p, so $F(x^*) - F(\bar{x})$ reduces to $\int_c |p-x^*|^2 ds - \int_c |p-\bar{x}|^2 ds$. But from Lemma 2, we know that $\int_c |p-x^*|^2 ds > \int_c |p-\bar{x}|^2 ds$, meaning $F(x^*) - F(\bar{x}) > 0$ and hence the line passing through x^* yields a greater objective function value than the line passing through \bar{x} . This contradicts the original supposition, proving the lemma.

Theorem 2. Assume that we are given a bounded, piecewise continuous space curve of finite length, C, having centroid $\bar{\mathbf{x}} = (\bar{x}, \bar{y}, \bar{z})$. Then a line of total least-squares can be computed as the line passing through the centroid and having the direction of the eigenvector corresponding to the largest eigenvalue of the 3×3 matrix:

$$\begin{bmatrix} \int_{c} (x-\bar{x})^{2} ds & \int_{c} (x-\bar{x})(y-\bar{y}) ds & \int_{c} (x-\bar{x})(z-\bar{z}) ds \\ \int_{c} (x-\bar{x})(y-\bar{y}) ds & \int_{c} (y-\bar{y})^{2} ds & \int_{c} (y-\bar{y})(z-\bar{z}) ds \\ \int_{c} (x-\bar{x})(z-\bar{z}) ds & \int_{c} (y-\bar{y})(z-\bar{z}) ds & \int_{c} (z-\bar{z})^{2} ds \end{bmatrix}$$

Proof. That the least-squares line passes through the centroid comes straightforwardly from Lemma 3. Knowing this, the objective function to be minimized can be written as a function of a alone:

$$F(\boldsymbol{a}) = \int_{C} |\boldsymbol{x}_{i} - \boldsymbol{x}|^{2} ds - \int_{C} (\boldsymbol{a} \cdot (\boldsymbol{x}_{i} - \boldsymbol{x}))^{2} ds$$

But the first term does not depend on a, so the function is minimized when the second term is maximized. But the second term is simply the objective function we had for the plane, which ended up equaling the eigenvalue as seen in Eq. (4) of the proof of Theorem 1. Thus, the a that maximizes the objective function here is found as the eigenvector corresponding to the *largest* eigenvalue of the matrix from Theorem 1 (with K=1) which is written out above in this theorem.

4 Fitting Algorithms

For the convenience of the reader, the algorithms themselves are presented in this section, unencumbered by their proofs, which appear later.

4.1 Fitting a Single Plane. Given:

- (1) Data points $x_1, x_2, x_3, ..., x_N$, where each $x_i = (x_i, y_i, z_i)$, and
- (2) The corresponding weights, $w_1, w_2, w_3, ..., w_N$, where all the weights are positive.

Then the weighted total least-squares plane is defined as the plane that minimizes $\sum_{i=1}^{N} w_i d_i^2$, where d_i is the orthogonal distance from the *i*th data point to the plane.

The weighted total least-squares plane can be found as follows:

- (1) A point on the plane is the weighted centroid, namely $\bar{x} = (\bar{x}, \bar{y}, \bar{z}) = \frac{\sum_{i=1}^{N} w_i x_i}{\sum_{i=1}^{N} w_i}$.
- (2) The unit vector normal to the plane is the (right) singular vector corresponding to the smallest singular value in the singular value decomposition (SVD) of the N × 3 matrix given by:

$$\mathbf{M} = \begin{bmatrix} \sqrt{w_1}(x_1 - \bar{x}) & \sqrt{w_1}(y_1 - \bar{y}) & \sqrt{w_1}(z_1 - \bar{z}) \\ \sqrt{w_2}(x_2 - \bar{x}) & \sqrt{w_2}(y_2 - \bar{y}) & \sqrt{w_2}(z_2 - \bar{z}) \\ \vdots & \vdots & \vdots \\ \sqrt{w_N}(x_N - \bar{x}) & \sqrt{w_N}(y_N - \bar{y}) & \sqrt{w_N}(z_N - \bar{z}) \end{bmatrix}$$

The unweighted case (i.e., the equally weighted case) can be found by making the value of all the weights equal to one; thus, removing their appearance from the matrix. Scaling all the weights by the any fixed, positive value does not affect the solution

There is a close connection between the smallest singular value and the rms value between the fitted plane and the surface (as approximated by the sampled points). This rms value is a quantity of interest in future ISO 1101 revisions. As we will later see in the proof of Theorem 3, the square of the smallest singular value equals the objective function—the weighted sum-of-squares of residuals. If we denote the smallest singular value as σ , and if the weights correspond to the discretely partitioned areas about the sampled points, totaling A, then the (discrete) rms value can be simply obtained as:

$$\sqrt{\frac{\sigma^2}{\sum_{i=1}^N w_i}} = \frac{\sigma}{\sqrt{A}}$$

This result can be seen to follow from the fact that the smallest eigenvalue of M^TM (or equivalently the square of the smallest

Journal of Computing and Information Science in Engineering

singular vector of M) equals the weighted sum-of-squares of distances from the points to the weighted least-squares plane (which is seen in the proof of Theorem 3).

4.2 Two Parallel Planes Given:

- (1) Two sets of data points $x_1, x_2, ..., x_N$ and $x_{N+1}, x_{N+2}, ..., x_{N+M}$, where each $x_i = (x_i, y_i, z_i)$, where it is known a priori which points belong to each plane, and
- (2) The corresponding weights, $w_1, w_2, w_3, ..., w_{N+M}$, where all the weights are positive.

Then the weighted total least-squares fitting of two parallel planes (e.g., corresponding to a slot or slab) is defined as the pair of parallel planes that minimizes $\sum_{i=1}^{N+M} w_i d_i^2$, where d_i is the orthogonal distance from the ith data point to the first plane when $1 \le i \le N$ or to the second plane when $N+1 \le i \le N+M$.

If two parallel planes are defined by a point on the first plane $x_A = (x_A, y_A, z_A)$, a point on the second plane $x_B = (x_B, y_B, z_B)$, and the unit vector normal to the parallel planes, a = (a, b, c), then the weighted total least-squares parallel planes can be found as follows:

- (1) A point on the first plane is the weighted centroid of the first set of data points, namely, $\bar{x}_A = (\bar{x}_A, \bar{y}_A, \bar{z}_A)$ $= \frac{\sum_{i=1}^N w_i x_i}{\sum_{i=1}^N w_i}.$
- (2) A point on the second plane is the weighted centroid of the second set of data points, namely, $\bar{x}_B = (\bar{x}_B, \bar{y}_B, \bar{z}_B)$ $= \frac{\sum_{i=N+1}^{N+M} w_i x_i}{\sum_{i=N+1}^{N+M} w_i}.$
- (3) The unit vector normal to the plane is the singular vector corresponding to the smallest singular value of the $(N+M) \times 3$ matrix given by:

$$\begin{bmatrix} \sqrt{w_1}(x_1 - \bar{x}_A) & \sqrt{w_1}(y_1 - \bar{y}_A) & \sqrt{w_1}(z_1 - \bar{z}_A) \\ \sqrt{w_2}(x_2 - \bar{x}_A) & \sqrt{w_2}(y_2 - \bar{y}_A) & \sqrt{w_2}(z_2 - \bar{z}_A) \\ \vdots & \vdots & \vdots \\ \sqrt{w_N}(x_N - \bar{x}_A) & \sqrt{w_N}(y_N - \bar{y}_A) & \sqrt{w_N}(z_N - \bar{z}_A) \\ \sqrt{w_{N+1}}(x_{N+1} - \bar{x}_B) & \sqrt{w_{N+1}}(y_{N+1} - \bar{y}_B) & \sqrt{w_{N+1}}(z_{N+1} - \bar{z}_B) \\ \sqrt{w_{N+2}}(x_{N+2} - \bar{x}_B) & \sqrt{w_{N+2}}(y_{N+2} - \bar{y}_B) & \sqrt{w_{N+2}}(z_{N+2} - \bar{z}_B) \\ \vdots & \vdots & \vdots \\ \sqrt{w_{N+M}}(x_{N+M} - \bar{x}_B) \sqrt{w_{N+M}}(y_{N+M} - \bar{y}_B) \sqrt{w_{N+M}}(z_{N+M} - \bar{z}_B) \end{bmatrix}$$

The distance between the two planes can be easily calculated as $|a \cdot (\bar{x}_A - \bar{x}_B)|$.

The unweighted case (i.e., the equally weighted case) can be found by making the value of all the weights equal to one; thus, removing their appearance from the matrix. Scaling all the weights (not just the weights for one plane) by the same factor does not affect the solution.

4.3 Arbitrarily Many Parallel Planes. The solution for two planes extends to arbitrarily many planes, where every plane passes through its weighted centroid, and the SVD is performed on the matrix written out in Theorem 3 in Sec. 5. We find it a pleasing result that the algorithm for parallel planes is found as such a straightforward extension of the single plane case.

4.4 Fitting a Line to 3D Coordinates. Given:

- (1) Data points $x_1, x_2, x_3, ..., x_N$, where each $x_i = (x_i, y_i, z_i)$, and
- (2) The corresponding weights, $w_1, w_2, w_3, ..., w_N$, where all the weights are positive.

The weighted total least-squares line can be found as follows:

- (1) A point on the line is the weighted centroid, namely $\bar{x} = (\bar{x}, \bar{y}, \bar{z}) = \frac{\sum_{i=1}^{N} w_i x_i}{\sum_{i=1}^{N} w_i}$.
- (2) The unit vector direction of the line is the (right) singular vector corresponding to the largest singular value in the singular-value decomposition (SVD) of the $N \times 3$ matrix denoted M in Sec. 4.1.

Note that this algorithm is identical to that for fitting a plane, except the singular vector corresponding to the largest (rather than smallest) singular vector is chosen. Also, the fitting of parallel lines in space is not considered in this paper, but could be performed in a manner that is likewise analogous to Secs. 4.2 and 4.3 (the fitting of parallel lines in a plane is considered in this paper, accomplished by implementing the parallel plane algorithm in two-dimensions).

There is a close connection between the largest singular value and the rms value between the fitted line and the curve (as approximated by the sampled points). This rms value is a quantity of interest in future ISO 1101 revisions. As we will later see in the proof of Theorem 4, the square of the smallest singular value is related to the objective function—the weighted sum-of-squares of residuals. If we denote the singular values as σ_1 , σ_2 , and σ_3 , with σ_3 being the largest and if the weights correspond to the discretely partitioned lengths about the sampled points, totaling the curve length L, then the (discrete) rms value can be simply obtained as:

$$\sqrt{\frac{\sigma_{1}^{2} + \sigma_{2}^{2}}{\sum_{i=1}^{N} w_{i}}} = \frac{\sqrt{\sigma_{1}^{2} + \sigma_{2}^{2}}}{\sqrt{L}}$$

This (perhaps surprising) result can be understood as follows: The planes orthogonal to the eigenvectors of M^TM and all passing through the centroid form three orthogonal planes (since M^TM is symmetric). The weighted sum-of-squares of distances from the points to any one of the three planes equals the eigenvalue of M^TM associated with that plane; this will be seen in the proof of Theorem 3. But the two planes corresponding to the smallest two eigenvalues both contain the weighted least-squares line. So by the Pythagorean Theorem, the distance from each point to the line is simply the root-sum-of-squares of the distances to each of the two planes. Summing the weighted squares of all these distances gives the sum of the eigenvalues of the two smallest eigenvalues or, equivalently, the sum of the squares of the smallest two singular values of M.

We note that when the points are projected into the plane perpendicular to the line, this result along with its proof bears strong resemblance to the perpendicular axis theorem (applicable for calculating moments of inertia).

5 Proofs of Algorithms in Discrete Cases

We prove the correctness of the algorithms given in Sec. 4 above. Before giving lemmas and theorems, we note that the least-squares solutions might not be unique (this is also true for the continuous cases already considered). While we do not rely on uniqueness in the proofs here, nonuniqueness arises only in pathological cases and is not a problem when we deal with planes in practical measurements on realistic parts [16]. Hence, we proceed, simply speaking of *the* least-squares plane or line. We start with a lemma related to the location of the weighted total least-squares plane.

LEMMA 4. (Discrete form of Lemma 1) Assume that we are given data points $x_1, x_2, x_3, ..., x_N$, where $x_i = (x_i, y_i, z_i)$, corresponding positive weights $w_1, w_2, w_3, \cdots, w_N$, and an arbitrary unit normal vector, $\mathbf{a} = (a, b, c)$. Then the weighted total least-squares plane constrained to have normal \mathbf{a} must pass through

031008-6 / Vol. 13, SEPTEMBER 2013

the weighted centroid, $\bar{\mathbf{x}} = (\bar{x}, \bar{y}, \bar{z}) = \frac{\sum_{i=1}^{N} w_i x_i}{\sum_{i=1}^{N} w_i}$. Furthermore, any

plane with normal a but not passing through the weighted centroid has a weighted sum-of-squares strictly greater than the plane passing through the centroid.

Proof. The equation of any plane having normal a can be written as $a \cdot x - d = 0$, where d is the signed distance from the plane to the origin. The signed orthogonal distance from any arbitrary point x_i to this plane is $a \cdot x_i - d$. Thus, the weighted sum-ofsquares objective function for any plane of orientation \boldsymbol{a} is given by $F(d) = \sum_{i=1}^{N} w_i (\boldsymbol{a} \cdot \boldsymbol{x}_i - d)^2$.

Taking the first and second derivatives yields

$$F'(d)=-2\sum_{i=1}^N w_i(\pmb a\cdot \pmb x_i-d),$$
 and
$$F''(d)=2\sum_{i=1}^N w_i>0,\quad \text{as all weights are positive}$$

The fact that the second derivative is a positive constant implies that the function is strictly convex, and has a unique minimum if and where its first derivative vanishes. This occurs when $\sum_{i=1}^{N} w_i(\boldsymbol{a} \cdot \boldsymbol{x}_i - d) = 0.$ Distributing the sum and solving for d

$$d = \frac{\sum_{i=1}^{N} w_i(\boldsymbol{a} \cdot \boldsymbol{x}_i)}{\sum_{i=1}^{N} w_i} = \boldsymbol{a} \cdot \left(\frac{\sum_{i=1}^{N} w_i \boldsymbol{x}_i}{\sum_{i=1}^{N} w_i}\right) = \boldsymbol{a} \cdot \bar{\boldsymbol{x}}$$

But this means that the distance from \bar{x} (the weighted centroid) to the plane is $\mathbf{a} \cdot \bar{\mathbf{x}} - d = 0$, implying that the weighted centroid must lie on the weighted total least-squares plane constrained to have normal a. Furthermore, the objective function is strictly convex, meaning any other plane not passing through the weighted centroid must have a greater weighted sum-of-squares.

Theorem 3. Assume that we are given $K \ge 1$ sets of data points $\{x_{1,1},x_{2,1},...,x_{N1,1}\}, \{x_{1,2},x_{2,2},...,x_{N2,2}\},...,\{x_{1,K},x_{2,K},...,x_{NK,K}\},$ where $x_{ij} = (x_{ij}, y_{ij}, z_{ij})$, and the corresponding positive weights: $w_{1,1}, w_{2,1}, ..., w_{N1,1}, w_{1,2}, w_{2,2}, ..., w_{NK,K}$, where all the weights are positive. Then a set of K parallel planes of weighted total leastsquares can be computed as the planes passing through the respective weighted centroids of the data sets and all sharing the same orientation defined by the singular vector corresponding to the smallest singular value of the matrix

$$\begin{bmatrix} \sqrt{w_{1,1}}(x_{1,1} - \bar{x}_{A1}) & \sqrt{w_{1,1}}(y_{1,1} - \bar{y}_{A1}) & \sqrt{w_{1,1}}(z_{1,1} - \bar{z}_{A1}) \\ \sqrt{w_{2,1}}(x_{2,1} - \bar{x}_{A1}) & \sqrt{w_{2,1}}(y_{2,1} - \bar{y}_{A1}) & \sqrt{w_{2,1}}(z_{2,1} - \bar{z}_{A1}) \\ \vdots & \vdots & \vdots \\ \sqrt{w_{N1,1}}(x_{N1,1} - \bar{x}_{A1}) & \sqrt{w_{N,1}}(y_{N1,1} - \bar{y}_{A1}) & \sqrt{w_{N1,1}}(z_{N1,1} - \bar{z}_{A1}) \\ \sqrt{w_{1,2}}(x_{1,2} - \bar{x}_{A2}) & \sqrt{w_{1,2}}(y_{1,2} - \bar{y}_{A2}) & \sqrt{w_{1,2}}(z_{1,2} - \bar{z}_{A2}) \\ \sqrt{w_{2,2}}(x_{2,2} - \bar{x}_{A2}) & \sqrt{w_{2,2}}(y_{2,2} - \bar{y}_{A2}) & \sqrt{w_{2,2}}(z_{2,2} - \bar{z}_{A2}) \\ \vdots & \vdots & \vdots \\ \sqrt{w_{NK,K}}(x_{NK,K} - \bar{x}_{AK}) \sqrt{w_{NK,K}}(y_{NK,K} - \bar{y}_{AK}) \sqrt{w_{NK,K}}(z_{NK,K} - \bar{z}_{AK}) \end{bmatrix} \end{bmatrix}$$
which, when expanded becomes
$$\begin{bmatrix} \sum_{i=1}^{N} w_{i} [\boldsymbol{a} \cdot (\boldsymbol{x}_{i} - \bar{\boldsymbol{x}}_{A})](x_{i} - \bar{\boldsymbol{x}}_{A}) + \sum_{i=N+1}^{N+M} w_{i} [\boldsymbol{a} \cdot (\boldsymbol{x}_{i} - \bar{\boldsymbol{x}}_{B})](x_{i} - \bar{\boldsymbol{x}}_{B}) \\ \sum_{i=1}^{N} w_{i} [\boldsymbol{a} \cdot (\boldsymbol{x}_{i} - \bar{\boldsymbol{x}}_{A})](y_{i} - \bar{y}_{A}) + \sum_{i=N+1}^{N+M} w_{i} [\boldsymbol{a} \cdot (\boldsymbol{x}_{i} - \bar{\boldsymbol{x}}_{B})](y_{i} - \bar{y}_{B}) \\ \sum_{i=1}^{N} w_{i} [\boldsymbol{a} \cdot (\boldsymbol{x}_{i} - \bar{\boldsymbol{x}}_{A})](z_{i} - \bar{z}_{A}) + \sum_{i=N+1}^{N+M} w_{i} [\boldsymbol{a} \cdot (\boldsymbol{x}_{i} - \bar{\boldsymbol{x}}_{B})](z_{i} - \bar{z}_{B}) \end{bmatrix}$$

Proof. For the sake of simplicity of notation and ease for the reader, we demonstrate the proof for the case of K=2, as we did in Theorem 1. Using simpler notation then, we assume two sets of data points, $x_1, x_2, ..., x_N$, and $x_{N+1}, x_{N+2}, ..., x_{N+M}$, and corresponding weights, $w_1, w_2, w_3, ..., w_{N+M}$. The weighted sum-ofsquares objective function is given by

$$F(\boldsymbol{a}, d_A, d_B) = \sum_{i=1}^{N} w_i [\boldsymbol{a} \cdot \boldsymbol{x}_i - d_A]^2 + \sum_{i=N+1}^{N+M} w_i [\boldsymbol{a} \cdot \boldsymbol{x}_i - d_B]^2$$

(Here, d_A and d_B are the distances from the two planes to the origin). For a fixed orientation, a^* , the objective function becomes

$$F(d_A, d_B) = \sum_{i=1}^{N} w_i [\mathbf{a}^* \cdot \mathbf{x}_i - d_A]^2 + \sum_{i=N+1}^{N+M} w_i [\mathbf{a}^* \cdot \mathbf{x}_i - d_B]^2$$

But we note that d_A occurs in the first sum only and d_B occurs in the second sum only, allowing us to express the objective function as $F(d_A,d_B) = F(d_A) + F(d_B)$ where $F(d_A)$ and $F(d_B)$ are defined as in the proof of Lemma 4 and each corresponding to its own set of data. This decomposition allows us to see that $F(d_A, d_B)$ is minimized when each term in its sum is minimized. That is,

$$\frac{\min}{d_A, d_B} [F(d_A, d_B)] = \frac{\min}{d_A} [F(d_A)] + \frac{\min}{d_B} [F(d_B)]$$
(5)

Thus, for the fixed orientation, the parallel planes minimizing the objective function are in fact the planes that each individually minimize the weighted sum-of-squares for its own set of data.

Lemma 4 can be applied then to the two minimization problems on the right hand side of Eq. (5), yielding that each plane must pass through the weighted centroid of its individual data set. Since this result is true for any fixed direction a^* , it follows that it holds for the least-squares direction in particular.

Knowing this, the objective function to be minimized can be written as a function of a alone

$$F(\boldsymbol{a}) = \sum_{i=1}^{N} w_i [\boldsymbol{a} \cdot (\boldsymbol{x}_i - \bar{\boldsymbol{x}}_A)]^2 + \sum_{i=N+1}^{N+M} w_i [\boldsymbol{a} \cdot (\boldsymbol{x}_i - \bar{\boldsymbol{x}}_B)]^2$$

To solve the orientation problem, we again use the method of Lagrange multipliers. The minimum of F(a) subject to G(a) = 0(where $G(\mathbf{a}) = |\mathbf{a}|^2 - 1$) occurs when $\nabla F = \lambda \nabla G$. In this case we

$$\nabla F = \begin{bmatrix} \frac{\partial F}{\partial a} \\ \frac{\partial F}{\partial b} \\ \frac{\partial F}{\partial c} \end{bmatrix}$$

$$\sum_{i=1}^{N} w_{i} [\boldsymbol{a} \cdot (\boldsymbol{x}_{i} - \bar{\boldsymbol{x}}_{A})] (x_{i} - \bar{\boldsymbol{x}}_{A}) + \sum_{i=N+1}^{N+M} w_{i} [\boldsymbol{a} \cdot (\boldsymbol{x}_{i} - \bar{\boldsymbol{x}}_{B})] (x_{i} - \bar{\boldsymbol{x}}_{B}) \\
2 \left[\sum_{i=1}^{N} w_{i} [\boldsymbol{a} \cdot (\boldsymbol{x}_{i} - \bar{\boldsymbol{x}}_{A})] (y_{i} - \bar{y}_{A}) + \sum_{i=N+1}^{N+M} w_{i} [\boldsymbol{a} \cdot (\boldsymbol{x}_{i} - \bar{\boldsymbol{x}}_{B})] (y_{i} - \bar{y}_{B}) \\
\sum_{i=1}^{N} w_{i} [\boldsymbol{a} \cdot (\boldsymbol{x}_{i} - \bar{\boldsymbol{x}}_{A})] (z_{i} - \bar{z}_{A}) + \sum_{i=N+1}^{N+M} w_{i} [\boldsymbol{a} \cdot (\boldsymbol{x}_{i} - \bar{\boldsymbol{x}}_{B})] (z_{i} - \bar{z}_{B})$$

Similar to the single plane case, computing the gradients yields an eigenvector problem, but the sum in each matrix entry is replaced by the two corresponding sums for data sets A and B.

As in the proof of Theorem 1, we have $\nabla G = 2a$ and $\nabla F = 2La = 2(M^{T}M)a$, where L is defined to be $M^{T}M$, and where *M* is now defined as the $(N+M) \times 3$ matrix given as

Journal of Computing and Information Science in Engineering

$$\begin{bmatrix} \sqrt{w_1}(x_1 - \bar{x}_A) & \sqrt{w_1}(y_1 - \bar{y}_A) & \sqrt{w_1}(z_1 - \bar{z}_A) \\ \sqrt{w_2}(x_2 - \bar{x}_A) & \sqrt{w_2}(y_2 - \bar{y}_A) & \sqrt{w_2}(z_2 - \bar{z}_A) \\ \vdots & \vdots & \vdots \\ \sqrt{w_N}(x_N - \bar{x}_A) & \sqrt{w_N}(y_N - \bar{y}_A) & \sqrt{w_N}(z_N - \bar{z}_A) \\ \sqrt{w_{N+1}}(x_{N+1} - \bar{x}_B) & \sqrt{w_{N+1}}(y_{N+1} - \bar{y}_B) & \sqrt{w_{N+1}}(z_{N+1} - \bar{z}_B) \\ \sqrt{w_{N+2}}(x_{N+2} - \bar{x}_B) & \sqrt{w_{N+2}}(y_{N+2} - \bar{y}_B) & \sqrt{w_{N+2}}(z_{N+2} - \bar{z}_B) \\ \vdots & \vdots & \vdots \\ \sqrt{w_{N+M}}(x_{N+M} - \bar{x}_B) & \sqrt{w_{N+M}}(y_{N+M} - \bar{y}_B) & \sqrt{w_{N+M}}(z_{N+M} - \bar{z}_B) \end{bmatrix}$$

The orientation, (a, b, c), can be found by solving the 3×3 eigenproblem given by $L\begin{bmatrix} a \\ b \\ c \end{bmatrix} = \lambda \begin{bmatrix} a \\ b \\ c \end{bmatrix}$, by using well known meth-

ods. As in the proof of Theorem 1, the eigenvectors of M^TM are also the singular vectors of M. This allows us to gain better numerical results by applying the SVD to M without ever computing M^TM

We further mimic the proof of Theorem 1 to select the proper eigenvector (singular vector). The normal equations can be written as follows:

$$\sum_{i=1}^{N} w_i(x_i - \bar{x})[\boldsymbol{a} \cdot (\boldsymbol{x}_i - \bar{\boldsymbol{x}}_A)] + \sum_{i=N+1}^{N+M} w_i(x_i - \bar{x})[\boldsymbol{a} \cdot (\boldsymbol{x}_i - \bar{\boldsymbol{x}}_B)] = \lambda a$$

$$\sum_{i=1}^{N} w_i(y_i - \bar{y})[\boldsymbol{a} \cdot (\boldsymbol{x}_i - \bar{\boldsymbol{x}}_A)] + \sum_{i=N+1}^{N+M} w_i(y_i - \bar{y})[\boldsymbol{a} \cdot (\boldsymbol{x}_i - \bar{\boldsymbol{x}}_B)] = \lambda b$$

$$\sum_{i=1}^{N} w_i(z_i - \bar{z})[\boldsymbol{a} \cdot (\boldsymbol{x}_i - \bar{\boldsymbol{x}}_A)] + \sum_{i=N+1}^{N+M} w_i(z_i - \bar{z})[\boldsymbol{a} \cdot (\boldsymbol{x}_i - \bar{\boldsymbol{x}}_B)] = \lambda c$$

Multiplying these equations by a, b, and c, respectively, then summing the equations gives

$$\sum_{i=1}^{N} w_{i} [\boldsymbol{a} \cdot (\boldsymbol{x}_{i} - \bar{\boldsymbol{x}}_{A})]^{2} + \sum_{i=N+1}^{N+M} w_{i} [\boldsymbol{a} \cdot (\boldsymbol{x}_{i} - \bar{\boldsymbol{x}}_{B})]^{2} = \lambda |\boldsymbol{a}|^{2} = \lambda \quad (6)$$

But the sum on the left is just the objective function, F(a), hence the correct eigenvector for the solution corresponds to the smallest eigenvalue (since $F(a) = \lambda$, and since we seek to minimize F). When using the SVD, we choose the singular vector corresponding to the smallest singular value, since under these conditions the singular values are the square roots of the eigenvalues.

LEMMA 5. (Discrete form of Lemma 2) Assume that we are given data points $\mathbf{x}_1, \mathbf{x}_2, \mathbf{x}_3, ..., \mathbf{x}_N$, where $\mathbf{x}_i = (x_i, y_i, z_i)$, corresponding positive weights $w_1, w_2, w_3, ..., w_N$. Then the point that minimizes the sum-of-squares objective function, $F(\mathbf{x}_0) = \sum_{i=1}^N w_i |\mathbf{x}_i - \mathbf{x}_0|^2$ is the weighted centroid, $\bar{\mathbf{x}} = (\bar{x}, \bar{y}, \bar{z})$ = $\frac{\sum_{i=1}^N w_i x_i}{\sum_{i=1}^N w_i}$. Any point that is not the centroid has a sum-of-squares objective function value strictly greater than that of the centroid.

Proof. Setting $\nabla F = 0$ yields $\sum_{i=1}^{N} w_i(x_i - x_0) = \sum_{i=1}^{N} w_i(y_i - y_0) = \sum_{i=1}^{N} w_i(z_i - z_0) = 0$. Solving for (x_0, y_0, z_0) gives $\left(\frac{\sum_{i=1}^{N} w_i x_i}{\sum_{i=1}^{N} w_i}, \frac{\sum_{i=1}^{N} w_i y_i}{\sum_{i=1}^{N} w_i}, \frac{\sum_{i=1}^{N} w_i z_i}{\sum_{i=1}^{N} w_i}\right)$, which is simply \bar{x} .

The Hessian matrix of
$$F$$
 is $2\begin{bmatrix} \sum_{i=1}^{N} w_i & 0 & 0\\ 0 & \sum_{i=1}^{N} w_i & 0\\ 0 & 0 & \sum_{i=1}^{N} w_i \end{bmatrix}$,

which is immediately seen as positive definite, implying that the

objective function is strictly convex and that the centroid is the unique point that minimizes F

LEMMA 6. (Discrete form of Lemma 3) Assume that we are given data points $x_1, x_2, x_3, ..., x_N$, where $x_i = (x_i, y_i, z_i)$, corresponding positive weights $w_1, w_2, w_3, ..., w_N$, and an arbitrary unit vector, $\mathbf{a} = (a, b, c)$. Then the weighted total least-squares line constrained to have direction \mathbf{a} must pass through the

weighted centroid, $\bar{x} = (\bar{x}, \bar{y}, \bar{z}) = \frac{\sum_{i=1}^{N} w_i x_i}{\sum_{i=1}^{N} w_i}$. Furthermore, any line with direction **a** but not passing through the weighted centroid

with direction **a** but not passing through the weighted centroid has a weighted sum-of-squares strictly greater than the line passing through the centroid.

Proof. Suppose not. That is, suppose the least-squares line does not pass through the weighted centroid, \bar{x} . We construct a plane perpendicular to the line and passing through \bar{x} . Call the point where the line intersects this plane, x^* (noting that $x^* \neq \bar{x}$). The remainder of the proof will show that the objective function of the line having direction a and passing through x^* is strictly greater than the objective function of the parallel line passing through \bar{x} , which contradicts the original supposition, since the line supposed to minimize the objective function indeed does not.

The objective function of the supposed least-squares solution is $F(\mathbf{x}^*) = \sum_{i=1}^N w_i \left[|\mathbf{x}_i - \mathbf{x}^*|^2 - \left(d^{\text{plane}}(\mathbf{x}_i, \mathbf{x}^*)\right)^2 \right]$. The objective function of the parallel line passing through $\bar{\mathbf{x}}$ is $F(\bar{\mathbf{x}}) = \sum_{i=1}^N w_i \left[|\mathbf{x}_i - \bar{\mathbf{x}}|^2 - \left(d^{\text{plane}}(\mathbf{x}_i, \bar{\mathbf{x}})\right)^2 \right]$. However, since the same constructed plane with normal \mathbf{a} contains both \mathbf{x}^* and $\bar{\mathbf{x}}$, we see that $d^{\text{plane}}(\mathbf{x}_i, \bar{\mathbf{x}}) = d^{\text{plane}}(\mathbf{x}_i, \mathbf{x}^*)$ for all points, \mathbf{x}_i , so $F(\mathbf{x}^*) - F(\bar{\mathbf{x}})$ reduces to $\sum_{i=1}^N w_i |\mathbf{x}_i - \mathbf{x}^*|^2 - \sum_{i=1}^N w_i |\mathbf{x}_i - \bar{\mathbf{x}}|^2$. But from Lemma 5, we know that $\sum_{i=1}^N w_i |\mathbf{x}_i - \mathbf{x}^*|^2 > \sum_{i=1}^N w_i |\mathbf{x}_i - \bar{\mathbf{x}}|^2$, meaning $F(\mathbf{x}^*) - F(\bar{\mathbf{x}}) > 0$; and hence, the line passing through \mathbf{x}^* yields a greater objective function value than the line passing through $\bar{\mathbf{x}}$. This contradicts the original supposition, proving the lemma.

THEOREM 4. Assume that we are given data points $x_1, x_2, x_3, ..., x_N$, where each $x_i = (x_i, y_i, z_i)$, and corresponding positive weights $w_1, w_2, w_3, \cdots, w_N$, then a weighted total least-squares line can be computed as the line passing through its weighted centroid, \bar{x} , and having its orientation defined by the singular vector corresponding to the largest singular value of the $N \times 3$ matrix denoted M in Sec. 4.1.

Proof. That the least-squares line passes through the centroid comes straightforwardly from Lemma 6. Knowing this, the objective function to be minimized can be written as a function of a alone

$$F(\mathbf{a}) = \sum_{i=1}^{N} w_i |\mathbf{x}_i - \bar{\mathbf{x}}|^2 - \sum_{i=1}^{N} w_i (d^{\text{plane}}(\mathbf{x}_i, \bar{\mathbf{x}}))^2$$

But the first term does not depend on a, so the function is minimized when the second term is maximized. But the second term is simply the objective function we had for the case of plane fitting, which ended up equaling the eigenvalue as seen in Eq. (6) of Theorem 3 (when K=1). Thus, the direction a that maximizes the second term in the objective function here is found as the singular vector corresponding to the *largest* singular value of the matrix from Theorem 1 (with K=1) which is denoted M in Sec. 4.1.

6 MATLAB Code and Numerical Testing

MATLAB code for the cases of fitting of one or two planes or a line is included here. The function names should be understood as wtlsqPlane="weighted total least-squares plane," wtlsq2pp="weighted total least-squares two parallel planes," wtlsqLine="weighted total least-squares line," In this code, w1 and w2 are column vectors of weights. pts1 and pts2 are matrices three columns wide containing the coordinates of the points, one point per row. q is a point on

031008-8 / Vol. 13, SEPTEMBER 2013

the single least-squares plane/line; q1 and q2 are points on the two parallel least-squares planes. In all cases, v is the unit vector that is normal to the least-squares plane(s) or in the direction of the least-squares line. The distance between the least-squares parallel planes can be computed as abs((q1-q2)*v).

```
function [q, v] = wtlsqPlane(w1, pts1)
 q = sum(bsxfun(@times,w1,pts1))/sum(w1);
 A1 = bsxfun(@minus,pts1,q);
A = [bsxfun(@times, sqrt(w1), A1)];
 [U, S, V] = svd(A, 0); % Can use [V, S] = eig(A'*A);
 [s, i] = min(diag(S));
v = V(:, i);
return;
function [q1, q2, v] = wtlsq2pp(w1, pts1, w2, pts2)
 q1 = sum (bsxfun (@times, w1, pts1))/sum (w1);
 q2 = sum(bsxfun(@times, w2, pts2))/sum(w2);
A1 = bsxfun(@minus,pts1,q1);
A2 = bsxfun(@minus,pts2,q2);
A = [bsxfun(@times, sqrt(w1), A1); ...
     bsxfun(@times,sqrt(w2),A2)];
 [U, S, V] = svd(A, 0); % Can use [V, S] = eig(A'*A);
 [s, i] = min(diag(S));
 v = V(:, i);
return;
function [q, v] = wtlsqLine(w1, pts1)
 q = sum (bsxfun (@times, w1, pts1)) / sum (w1);
A1 = bsxfun(@minus,pts1,q);
A = [bsxfun(@times, sqrt(w1), A1)];
 [U,S,V] = svd(A,0); % Can use [V,S] = eig(A'*A);
 [s, i] = \max(\text{diag}(S));
v = V(:, i);
return;
```

The algorithms in this paper were also implemented in Mathematica for the cases of fitting one, two, and three parallel planes and for fitting a line. The MATLAB and Mathematica algorithms were compared with each other to assure they give the same results (up to computational precision limits). The Mathematica code was also used to test the algorithms presented in this paper against a brute-force minimization search algorithm.

Test data sets were simulated randomly. Planar data sets were generated having varying aspect ratios, varying distances between them, varying numbers of points per plane, varying weights assigned, varying perturbations of the points from an exact plane, and varying nominal orientations of the planes to one another. Data sets for testing line fitting were generated having various lengths, numbers of points, weights, and perturbations of data points from a straight line.

The brute-force iterative search was performed using Mathematica's FindMinimum function. The options were tweaked to improve the desired accuracy and increase the working precision beyond machine precision in seeking accurate answers. The required initial guess for the function was obtained using knowledge of how the test data sets were generated. In contrast to the iterative search method, the algorithms presented in this paper were coded using normal machine precision, purposely not taking advantage of Mathematica's ability to increase the precision.

Results are shown here for the cases of two parallel planes for 100 simulated data sets (all other results were similar). Figure 5 shows a histogram of the anglular differences between the normal directions computed by the two methods. Figure 6 shows a

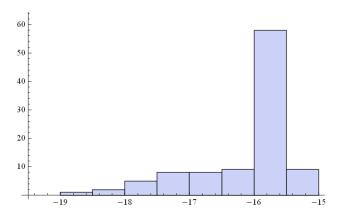


Fig. 5 Histogram of angular deviations in radians (log base 10 scaling)

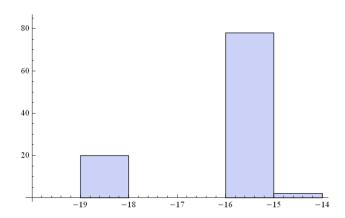


Fig. 6 Histogram of relative distance deviations (mm, log base 10 scaling)

histogram of the relative distances between the two computed parallel plane pairs. In both figures, the deviations are given on a log base 10 scale, meaning the greatest deviation shown in Fig. 5 is 10^{-15} radians and the greatest deviation shown in Fig. 6 is 10^{-14} mm. The magnitudes of these maximum deviations are more important than the histograms themselves, since they show that in all of the cases tested, the two methods agreed to within amounts attributable to machine precision.

7 Convergence of the Discrete to the Continuous Cases

Tolerance specifications according to standards are considered to apply to the continuous surface. Verification of such specifications is done using discrete points. Thus, we seek to prove the fundamental connection between the discrete and continuous cases. Specifically, we seek to show that the discrete cases converge to their corresponding continuous cases as the sampled points become dense (under idealized conditions of no measurement error and when the weights are the lengths/areas of the arcs/patches associated with each point). We show this first for the case of planes and then note that the case of lines is similar.

We seek to prove convergence for surfaces over which 2nd degree polynomials (as encountered in the continuous case in Theorem 1) are Riemann integrable—surfaces that are piecewise smooth, bounded, having finite area—conditions that are reasonable for real workpieces. In order to prove that the discrete solution converges to the continuous solution as the points become

⁴When considering convergence to the continuous case, we treat the surfaces as mathematical, ignoring the fact that, at very small scales, the molecular makeup of the material differs from our understanding of a continuous, mathematical surface.

dense (when the weights are assigned the values of the areas of the surface corresponding to the points, as shown in Sec. 2) we show the following three steps:

- The weighted centroid computed discretely converges to the centroid in the continuous case as the points become dense.
- (2) Every individual cell value of the 3 × 3 matrix in the discrete case (denoted L in the proof of Theorem 3) converges to its corresponding cell value in the matrix in the continuous case (also denoted L in the proof of Theorem 1) as the points become dense.
- (3) The fact that the individual cell values converge implies that the eigenvalues and eigenvectors from the discrete case also converge to their continuous case counterparts as the points become dense.

Step 1: Looking at the first coordinate in the weighted centroid (the others being similar), we have $\bar{x} = \sum_{i=1}^{N} w_i x_i / \sum_{i=1}^{N} w_i$. We want to show this value converges to its continuous counterpart, $\int_{S}^{x} ds / \int_{S}^{d} s$. Since we are assuming the weights are the areas of the patches associated with each point, the denominator in both the discrete and continuous cases is simply the area of the patch. Thus, we need to show convergence only of the numerator. The numerator converges as the points become dense because it forms a converging sequence of Riemann sums. The details of this statement are the same as will be given in step 2, below.

Step 2: That the sum converges to the surface integral for each of the nine cells is immediate from the definition of the Riemann surface integral (see, e.g., 1.5.1 in Ref. [17]). Specifically, if a partition of a surface called S is given by $S^{(n)} = \{S_1^{(n)}, S_2^{(n)}, \dots, S_n^{(n)}\}$ where these disjoint subsets have well defined areas, and if n points are chosen such that $x_i \in S_i^{(n)}$, for $i = 1, 2, \dots, n$, then the Riemann sum for a real valued function f over S is defined by

$$R_n = \sum_{i=1}^n \left(\Delta A_i^{(n)} \right) f(\mathbf{x}_i) \tag{7}$$

where $\Delta A_i^{(n)}$ represents the area of the subset $S_i^{(n)}$. If all sequences $\{R_n\}$ approach the same limit R with $\Delta_n = \max\{\operatorname{Area}(S_1^{(n)}), \ldots, \operatorname{Area}(S_1^{(n)})\} \to 0$ as $n \to \infty$, then f is defined as $Riemann\ integrable$ over S and the surface integral of f over S is defined by $\int_S^f (x) dS := R$.

In our discrete case, each cell in the matrix L contains one or more sums of the form $\sum_{i=1}^{N} w_i f(x_i)$. But, since we are assigning each w_i to have the value of its point's associated area, we have in fact a Riemann sum as in Eq. (7). Furthermore, as the points become dense, we then have a sequence of Riemann sums where the maximum area of a partition approaches zero. Then by definition, this must converge to its corresponding integral when the function is integrable over S. However, since the function f is simply a second degree polynomial, it is uniformly continuous over the piecewise smooth S and thus is integrable. Hence, Step 2 is shown.

Step 3: We now show that, since the individual cell values converge, the smallest eigenvalue and its corresponding eigenvector converge as well (when the smallest eigenvalue is unique, i.e., simple).

We first note that since there are a finite number of cells (namely nine) the convergence is uniform. That is, for any $\epsilon>0$ the points will become dense enough that every cell in the discrete case differs from the continuous case by less than ϵ .

We look at the sensitivities of the eigenvalues and eigenvectors in this case. Since the matrix is symmetric, we gain the advantage of several theorems that bound the changes in eigenvalues and eigenvectors under small changes to the matrix.

First, in this symmetric case, Stewart shows (p. 309 of Ref. [18]) that the eigenvalues are perfectly conditioned—that suffi-

ciently small ϵ -sized perturbations in the cells of the matrix yield differences in the eigenvalues essentially no greater than ϵ . Thus, convergence of eigenvalues is assured. Furthermore, this ensures that if the smallest eigenvalue in the continuous case is unique, then the smallest eigenvalue in the discrete case will eventually be as well.

It is also shown (p. 310 of Ref. [18]; see also Refs. [19,20]) that the eigenvector associated with the any unique eigenvalue λ is also well conditioned. Specifically, if the unit eigenvector associated with the smallest eigenvalue of the matrix L is denoted v, and the corresponding unit eigenvector of matrix $L + \Delta L$ is $v + \Delta v$, then the size of Δv is bounded as follows:

$$\|\Delta v\|_2 \le \frac{\|\Delta L\|_2}{\min|\lambda - \lambda_i|} + O\left(\|\Delta L\|_2^2\right)$$

where λ_j (j=1,2) represents the other two eigenvalues of the matrix. Thus (eventually), as the points become dense, we can find a constant κ such that if every element of ΔL is less than ϵ , then $\|\Delta v\|_2 \le \kappa \epsilon$. This means that—provided the smallest eigenvalue is unique—its eigenvector must also converge from the discrete to continuous case as the points become dense. Convergence for the case of lines is similar, only the largest eigenvalue must be unique instead of the smallest.

8 Concluding Remarks

We have presented and proved solutions to the problems of weighted total least-squares fitting of lines, planes and parallel planes. Furthermore, the solutions are conducive to implementation in computer algorithms using reliable and readily available linear algebra functions. The weighted fitting cases can be easily simplified to equally weighted fitting if desired. The need for such algorithms has been demonstrated and is relevant to newer tolerancing standards and instrumentation.

Furthermore, these fits in the discrete cases converge to their corresponding continuous cases as the points become dense (independent of the sampling strategy). This result is not generally true in the case of unweighted fitting.

Future work includes creating an appropriate practical means to assign weights to points as well as extending the work to nonlinear geometries.

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031008-10 / Vol. 13, SEPTEMBER 2013

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